

Resilience in Complex Linked Systems

Fred Roberts, DIMACS



Image credits:

Hurricane damage: FEMA Photo by
photographer Leif Skoogfors

Forest fire: USFS Region 5

Ebola treatment unit: CDC Global

No changes made in any image

DIMACS

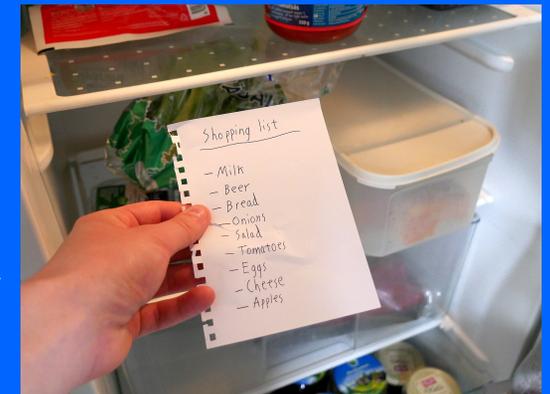
Center for Discrete Mathematics & Theoretical Computer Science
Founded as a National Science Foundation Science and
Technology Center



Resilience

• Today's society has become dependent on complex systems, enabled by increased digitization of our world, that have had a great impact on virtually all facets of our lives:

- Instant communication
- Ability to move money anywhere and quickly
- Ability to ask a machine to make our shopping list or turn on our favorite music.



- Yet these changes have made us vulnerable.
- To natural disasters, deliberate attacks, just plain errors.
- In recent years, *“resilience”* of complex natural and social systems has become a major area of emphasis.

Credit: Santeri Viinamäki via Wikimedia commons no changes made

Resilience



Hospital in Kansas during 1918 influenza pandemic

•Resilience in response to hurricanes, disease events, floods, earthquakes, cyber attacks, ...

Image credits:

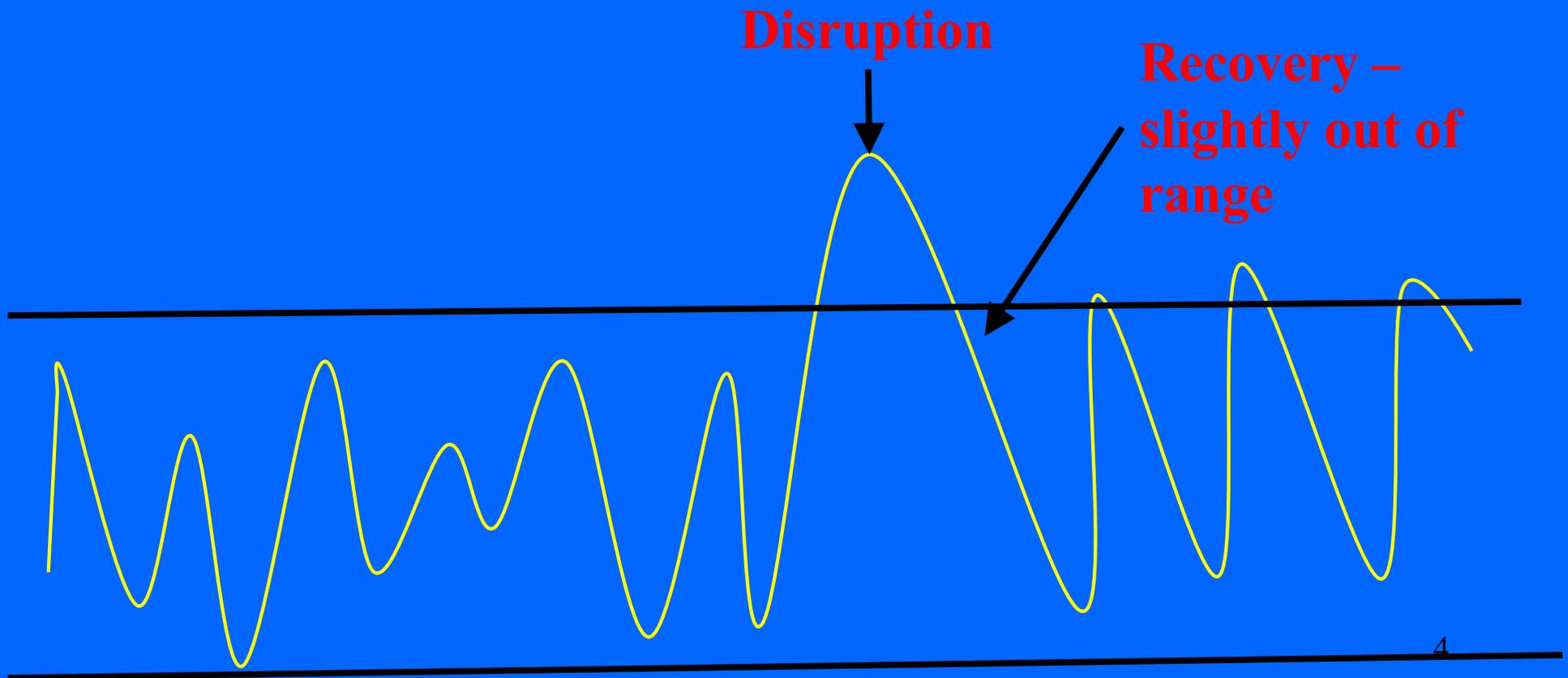
Earthquake: U.S. Air Force photo by Master Sgt. Jeremy Lock via Wikimedia commons

Flood: Voice of America Indonesian Service via Wikimedia commons

1918 influenza outbreak: Otis Historical Archives, National Museum of Health and Medicine via Wikimedia.com No changes made in any image

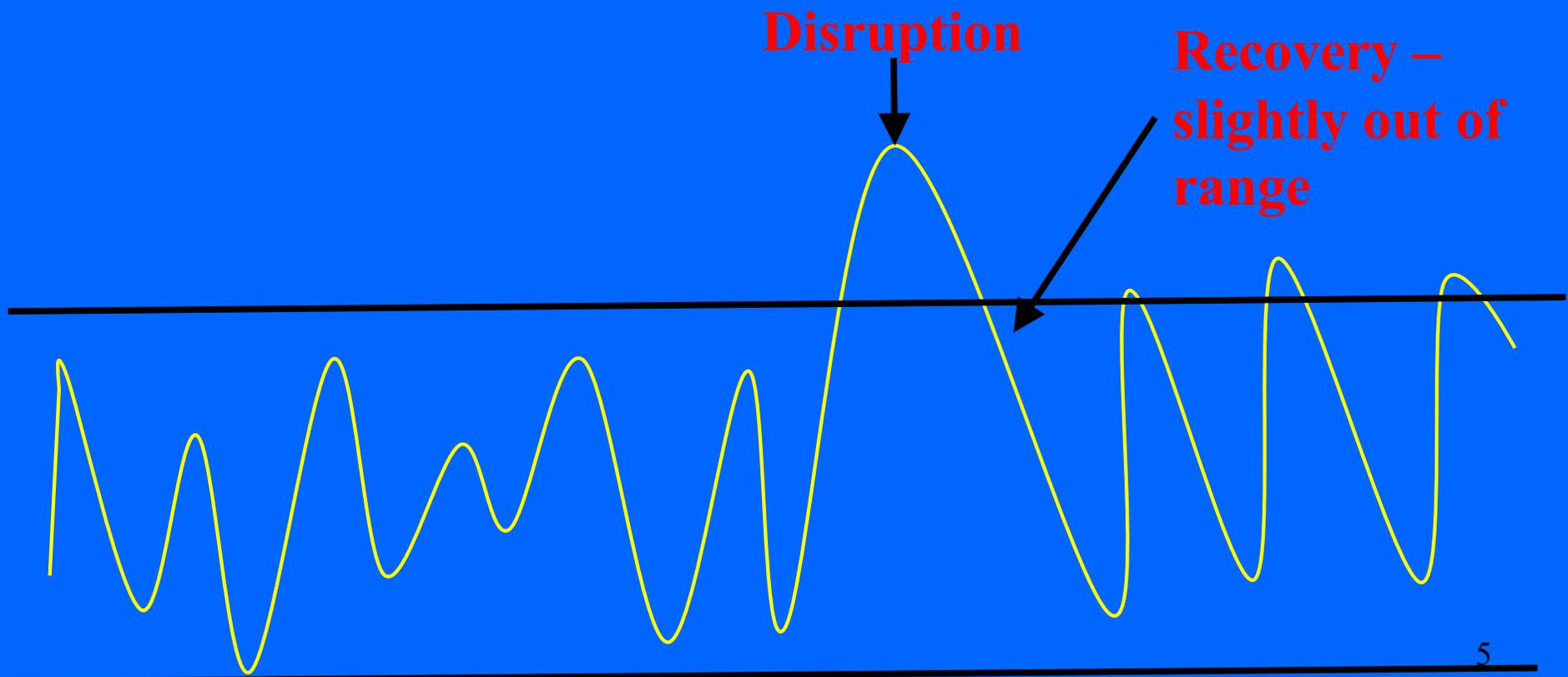
Resilience

- General concept of resilience: ability of a system to recover from disasters or attacks and avoid catastrophic collapse.



Resilience

- In a resilient system, values will return to the normal healthy range.
- Or they might establish a new healthy range – one that is not that far from the previous one



Resilience

- There are many parameters that measure a “healthy” system.
- Some will get back into their normal healthy range faster than others.
- Do we ask that the longest time to return to this range be small?
- Or that the average time to return to this range be small?

Approaches to Achieving Resilience

- One approach to resilience is to develop algorithms for responding to a disruption that will minimize the departure from the previous state when things settle down.
- Another is to design systems that can bounce back from disruptions quickly.
- I will emphasize the former.
- Will illustrate with four examples built around models using graphs and networks.

Example I: Spread and Control of Disease



- The spread of the new Coronavirus COVID-19 is just the latest worrisome example of a newly emerging disease that threatens not only lives but our economy and our social systems.

Image credit: Wikimedia commons

<https://www.youtube.com/watch?v=SBboFVjLQak> , 1:10

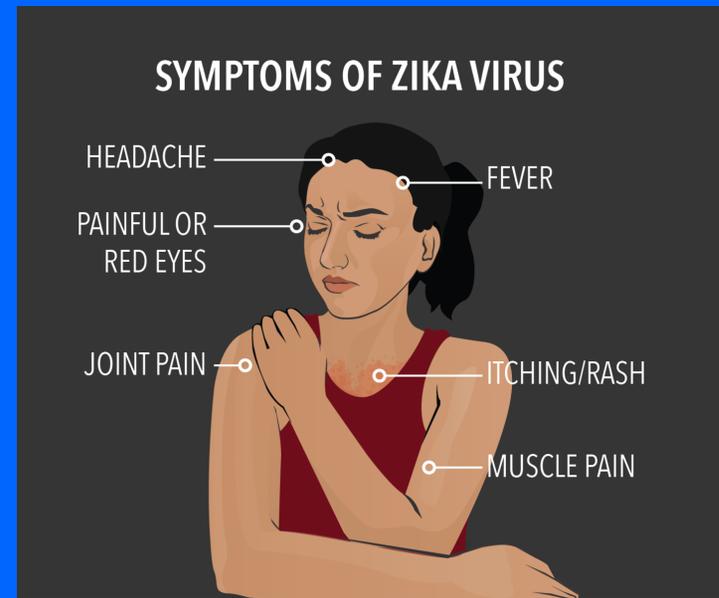
Chinanews.com/China News Service

Unchanged

Example I: Spread and Control of Disease



Ebola



Zika

- Ebola, Zika are other recent examples

Image credits: Wikimedia commons

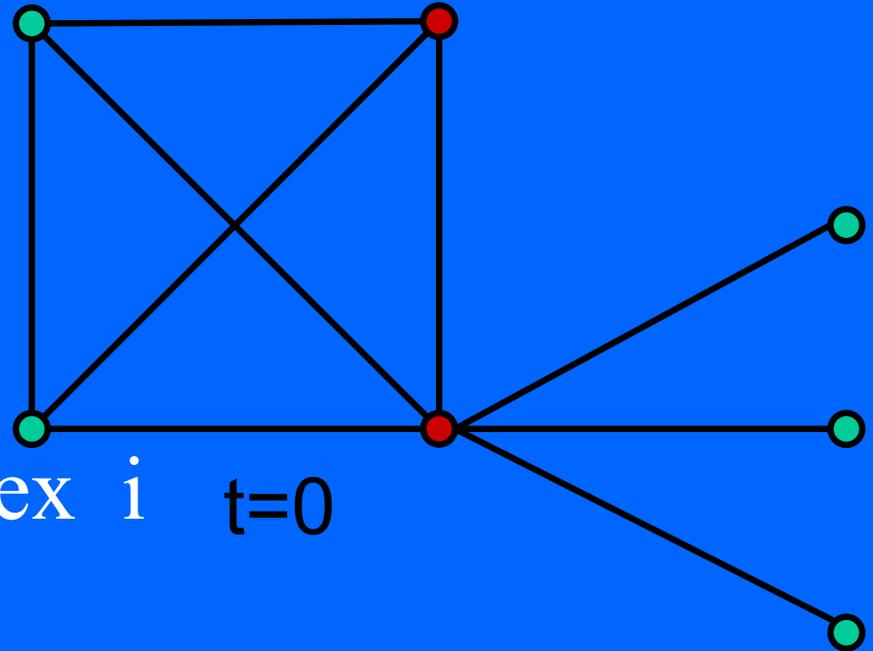
Ebola: Army Medicine; Zika: Beth.herlin no changes made

Example I: Spread and Control of Disease

- Modern transportation systems allow for rapid spread of diseases.
- Diseases are spread through social networks.
- “*Contact tracing*” is an important part of any strategy to combat outbreaks of infectious diseases, whether naturally occurring or resulting from bioterrorist attacks.
- I will illustrate the ideas with some fairly simple “toy” models that will illustrate concepts of resilience.

Simple Model: Moving From State to State

Social Network = Graph
Vertices = People
Edges = contact



Let $s_i(t)$ give the state of vertex i at time t .

Very simplified “toy” model: Two states: ● ●
● = susceptible, ● = infected (SI Model)

Times are discrete: $t = 0, 1, 2, \dots$

The Model: Moving From State to State

More complex models: SI, SEI, SEIR, etc.

S = susceptible, E = exposed, I = infected, R = recovered (or removed)



measles



SARS

Credit: measles: Wikimedia.org
SARS: Medical News Today

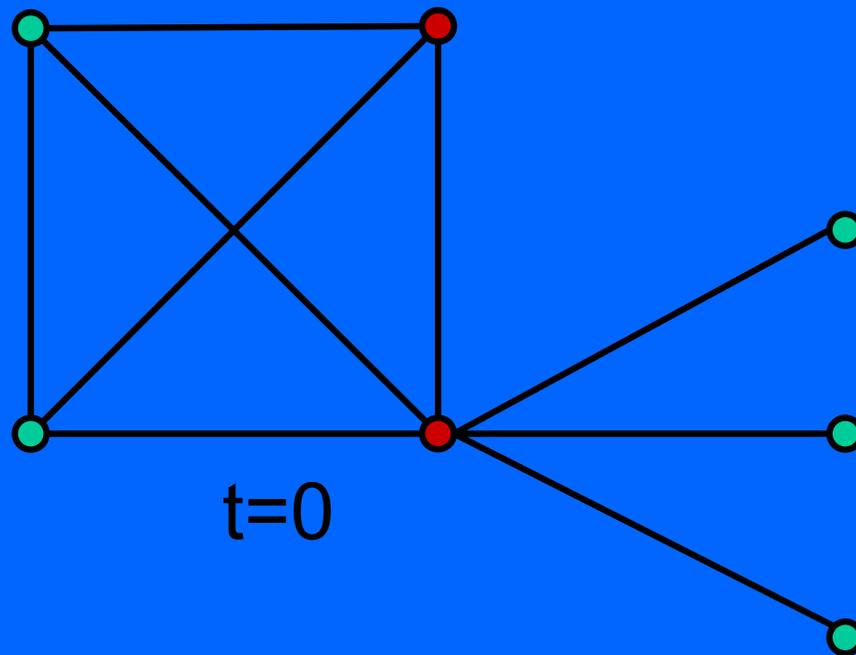
Threshold Processes

Irreversible k -Threshold Process: You change your state from \bullet to \bullet at time $t+1$ if at least k of your neighbors have state \bullet at time t . You never leave state \bullet .

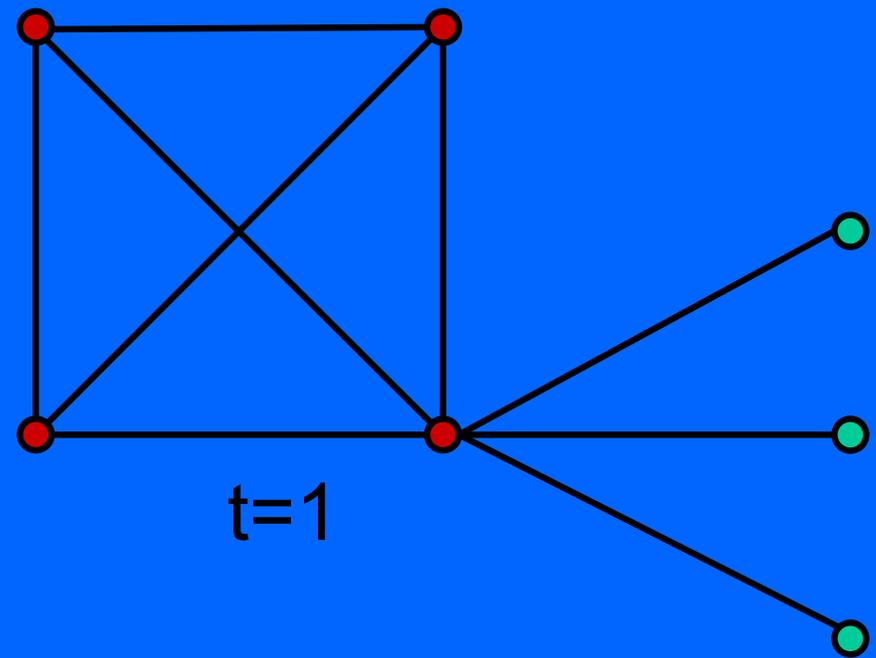
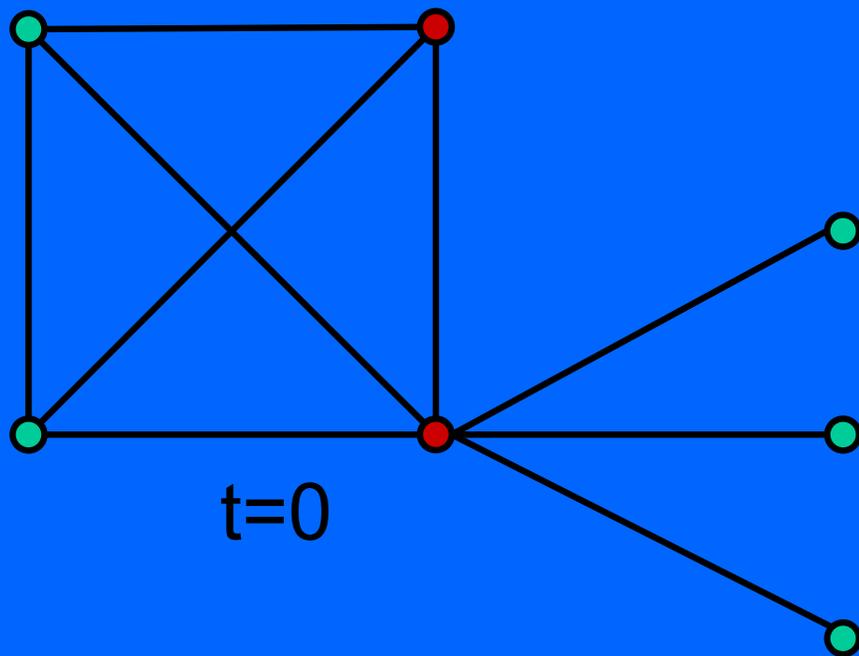
Disease interpretation? Infected if sufficiently many of your neighbors are infected.

Special Case $k = 1$: Infected if any of your neighbors is infected.

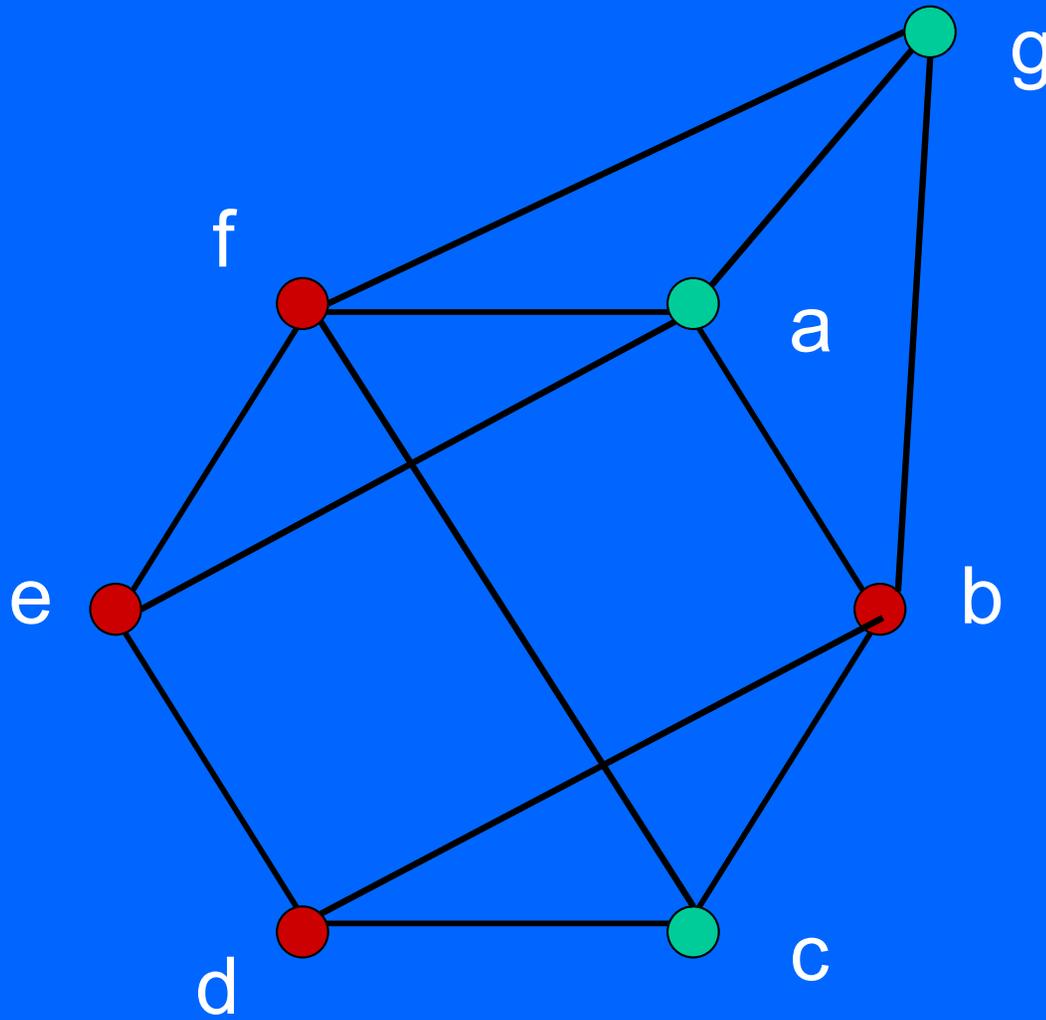
Irreversible 2-Threshold Process



Irreversible 2-Threshold Process

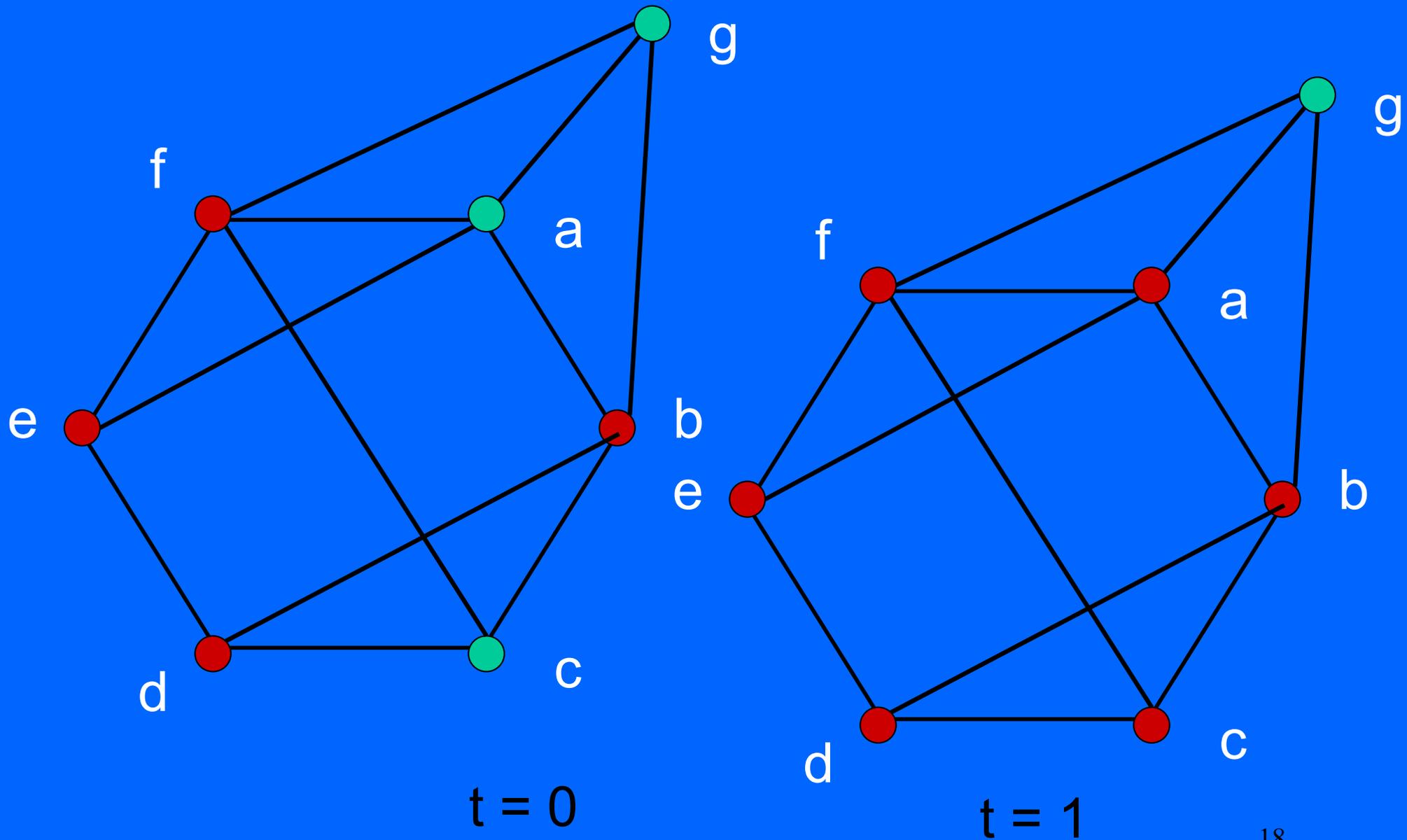


Irreversible 3-Threshold Process

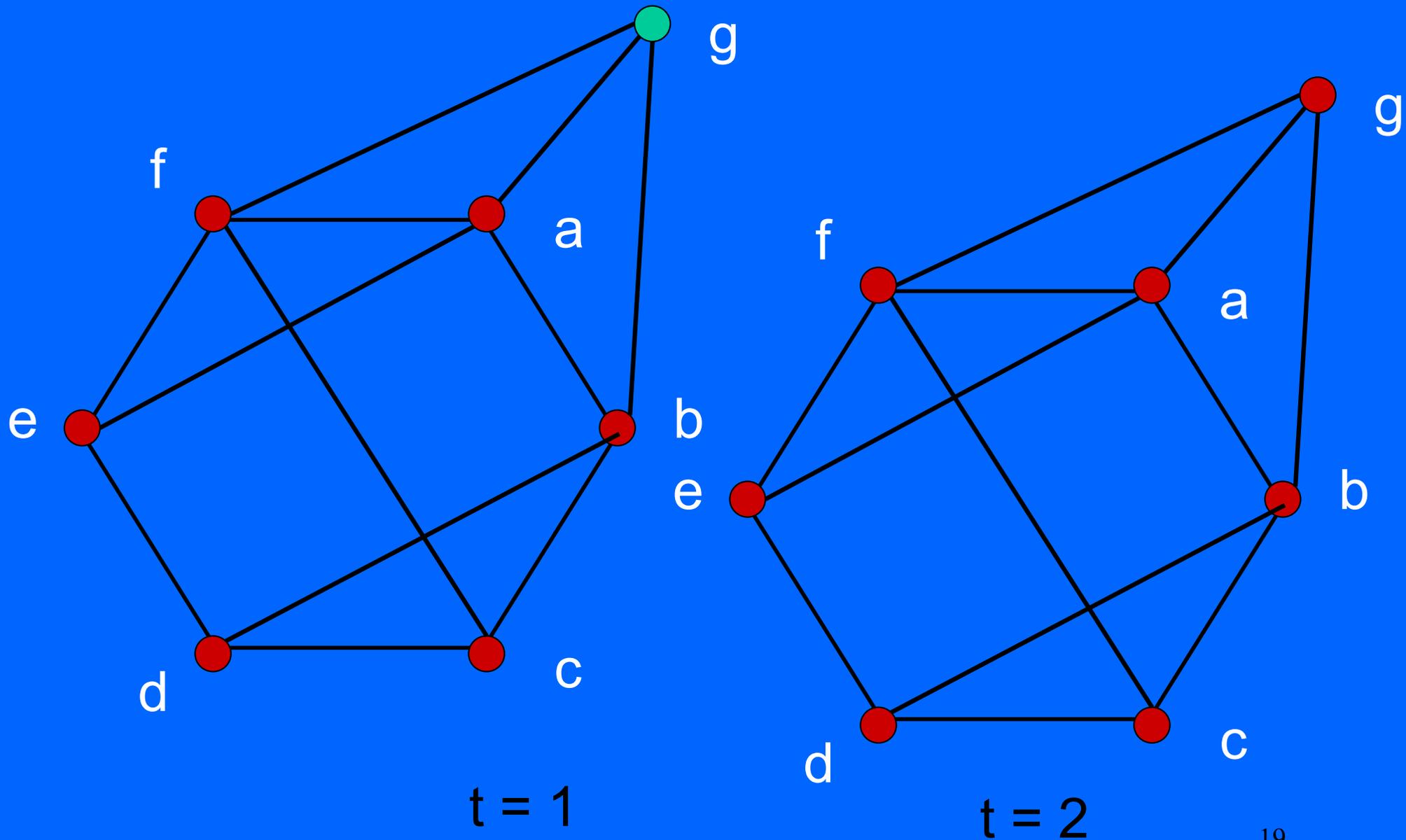


$t = 0$

Irreversible 3-Threshold Process



Irreversible 3-Threshold Process



The Saturation Problem

A great deal of attention has been paid to:

Attacker's Problem: Given a graph, what subsets S of the vertices should we plant a disease with so that ultimately the maximum number of people will get it?

Economic interpretation: What set of people do we place a new product with to guarantee “saturation” of the product in the population?

These Problems are "Hard"

Problem IRREVERSIBLE k-CONVERSION

SET: Given a positive integer p and a graph G , does G have a set S of size at most p so that if all vertices of S are infected at the beginning, then all vertices will ultimately be infected?

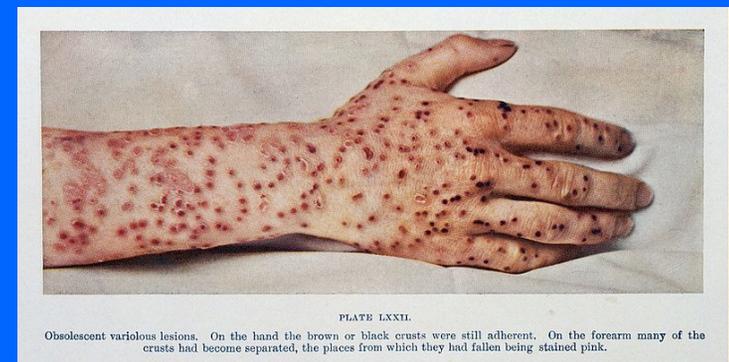
Theorem (Dreyer and Roberts): IRREVERSIBLE k-CONVERSION SET is NP-complete for fixed $k > 2$.

Complications to Add to Model

- $k = 1$, but you only get infected with a certain probability.
- You are automatically cured after you are in the infected state for d time periods.
- *A public health authority has the ability to “vaccinate” a certain number of vertices, making them immune from infection.*
- It's the vaccination strategy that relates to the resilience question.

Image credit:

https://wellcomeimages.org/indexplus/obf_images/b2/a8/9ca500938fc44f77d4c4e49a4d90.jpg



Smallpox

Vaccination Strategies



Credit: wikimedia
commons.org

Mathematical models are very helpful in comparing alternative vaccination strategies. The problem is especially interesting if we think of protecting against deliberate infection by a bioterrorist attacker but applies if we think of "nature" as the attacker.

Example II: Vaccinations and Fighting Fires

Stephen Hartke and others worked on a vaccination problem:

Defender: can vaccinate v people *per time period*.

Attacker: can only infect people at the beginning.

Irreversible k -threshold model.

What vaccination strategy minimizes number of people infected?

Variation: The vaccinator and infector alternate turns, having v vaccinations per period and i doses of pathogen per period.

What is a good strategy for the vaccinator?

Example II: Vaccinations and Fighting Fires

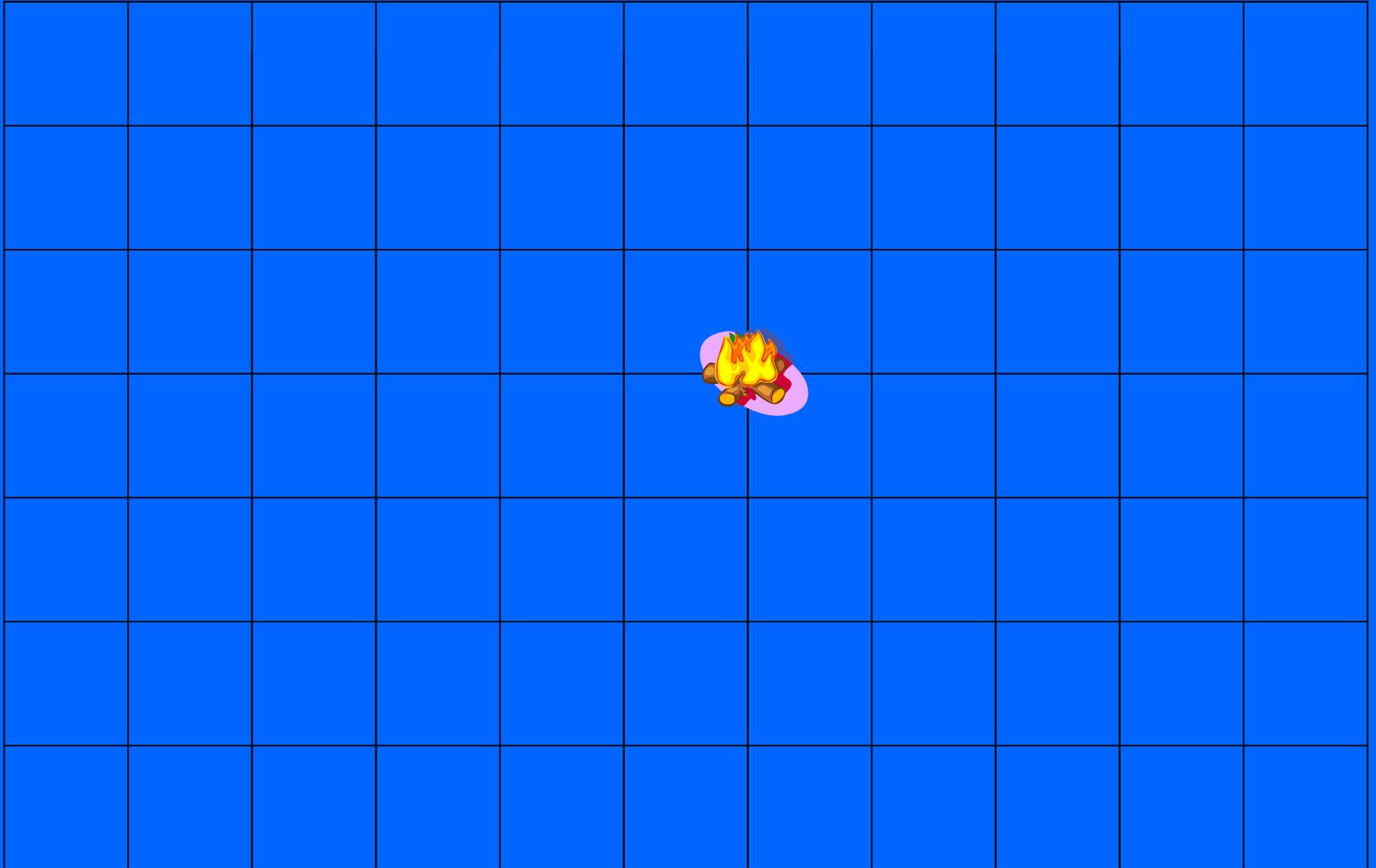
Sometimes called the *firefighter problem*:
alternate fire spread and firefighter placement.
Usual assumption: $k = 1$. (We will assume this.)

Problem goes back to Bert Hartnell 1995

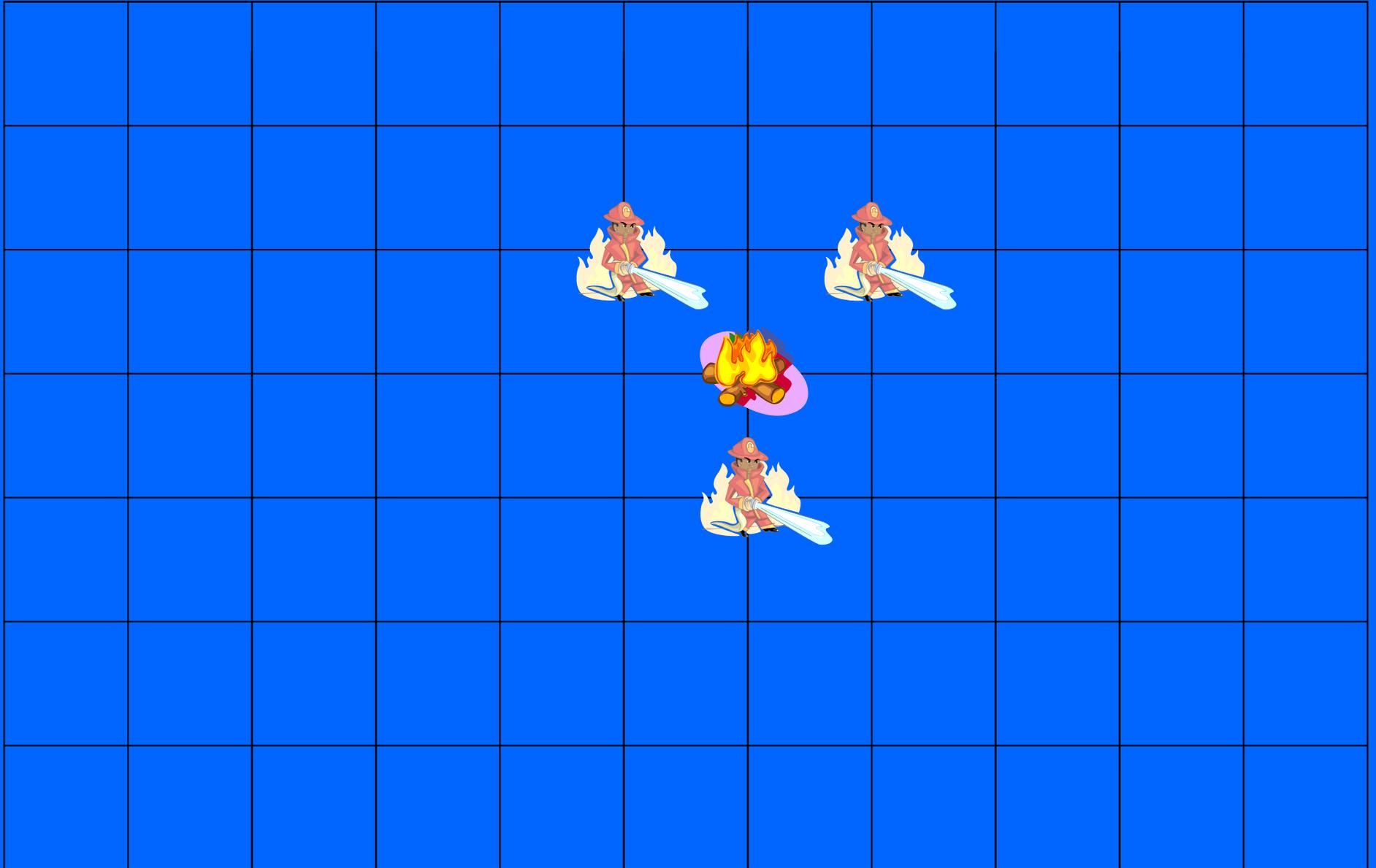
Image credit:
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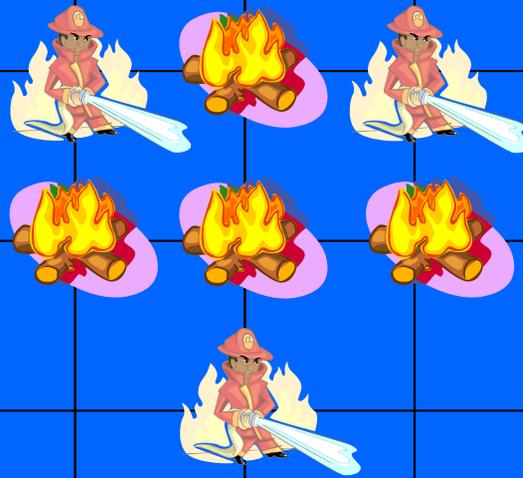
A Simple Model ($k = 1$) ($v = 3$)



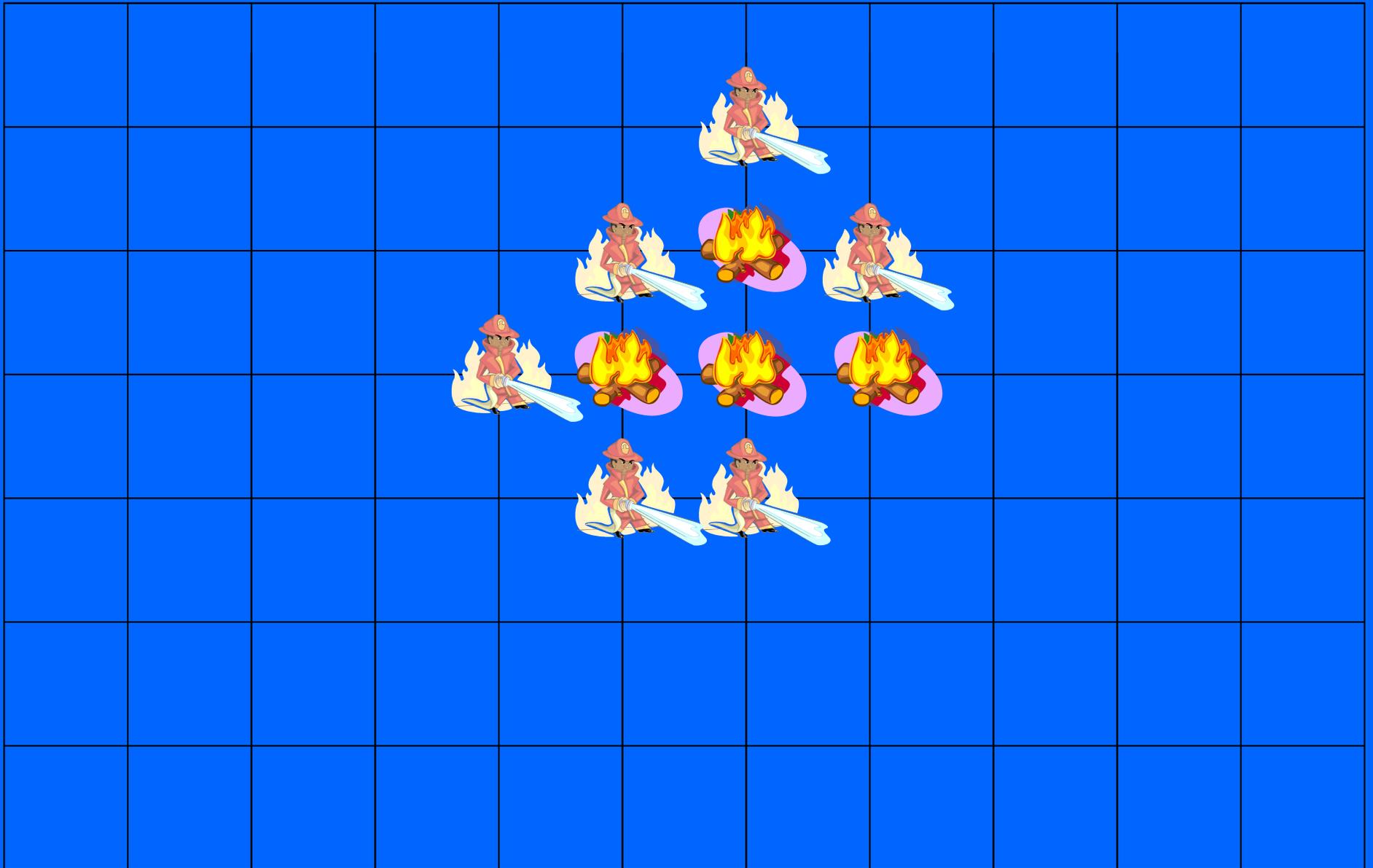
A Simple Model



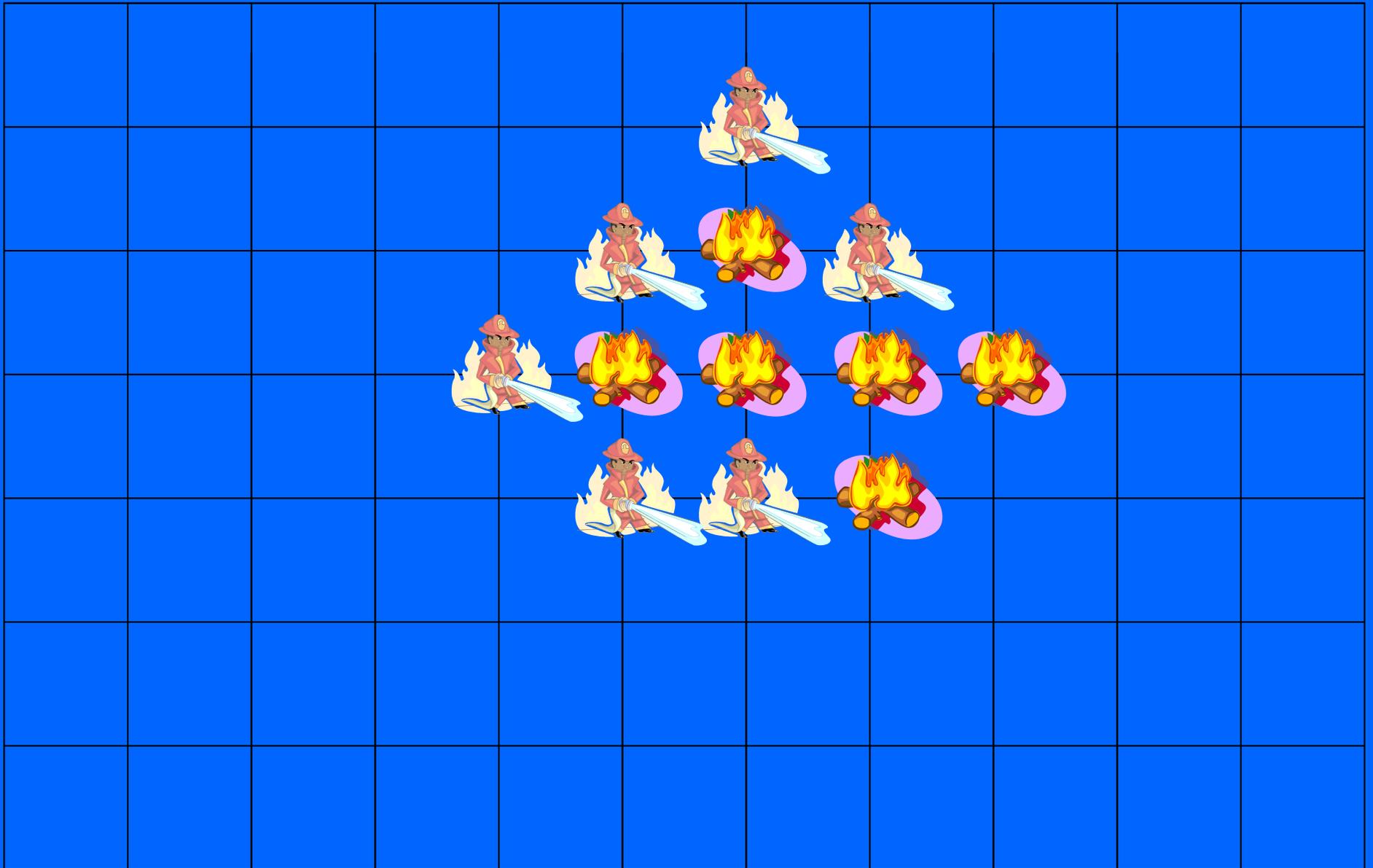
A Simple Model



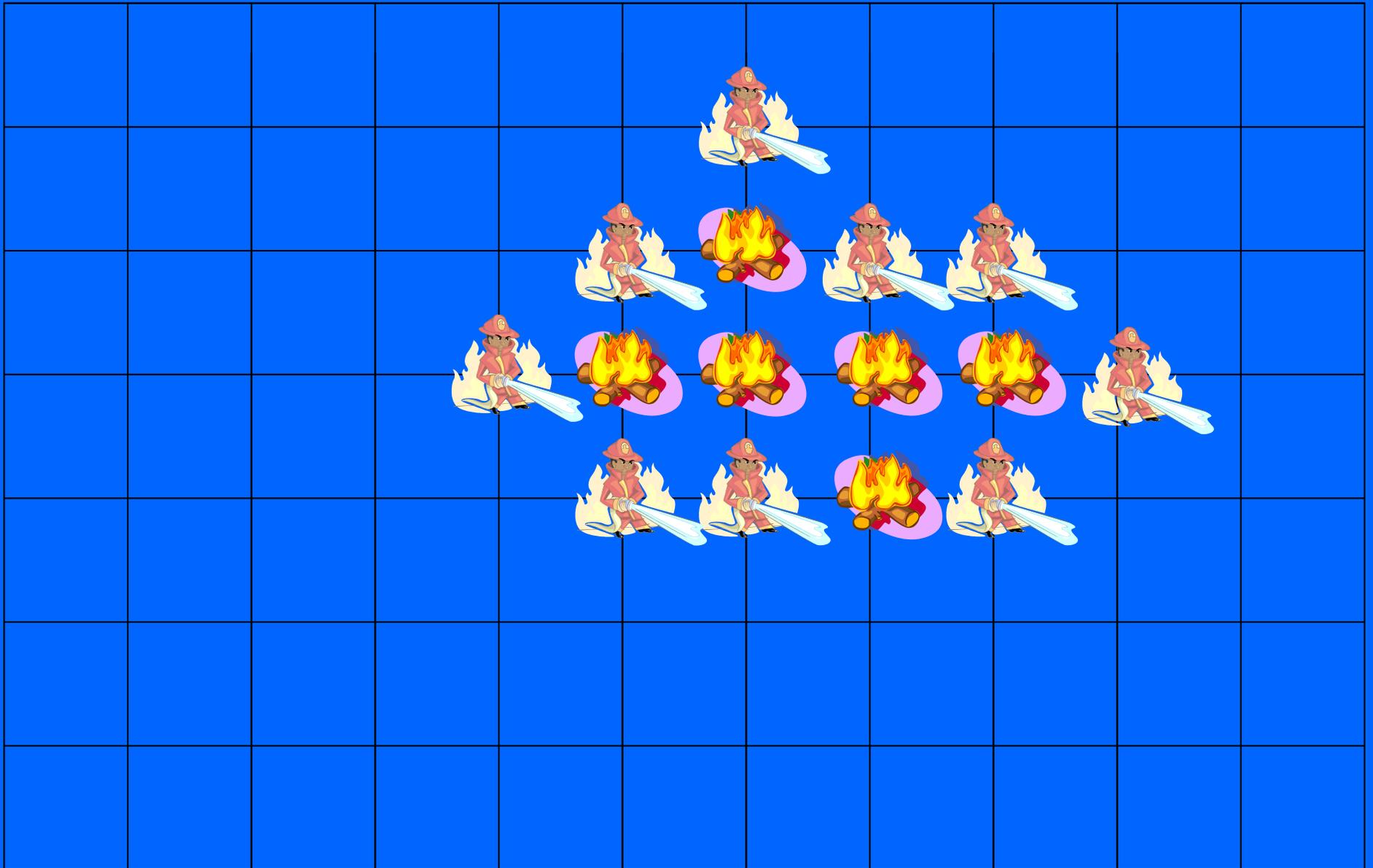
A Simple Model



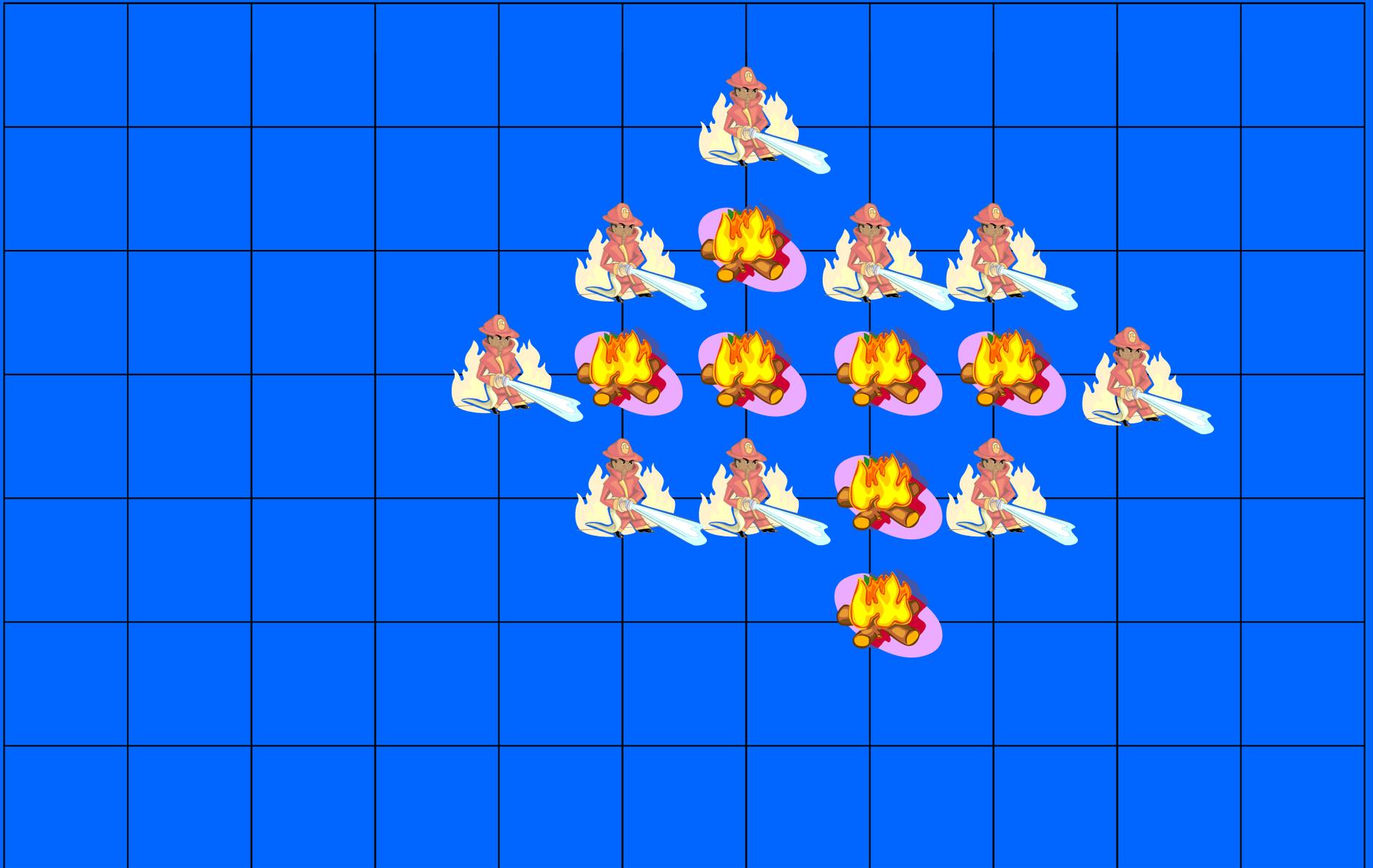
A Simple Model



A Simple Model



A Simple Model



A Simple Model





Some resilience questions that can be asked



- Can the fire (epidemic) be contained?
- How many time steps are required before fire is contained?
- How many firefighters per time step are necessary?
- What fraction of all vertices will be saved (burnt)?
- Does where the fire breaks out matter?
- Fire starting at more than 1 vertex?



Containing Fires in Infinite Grids

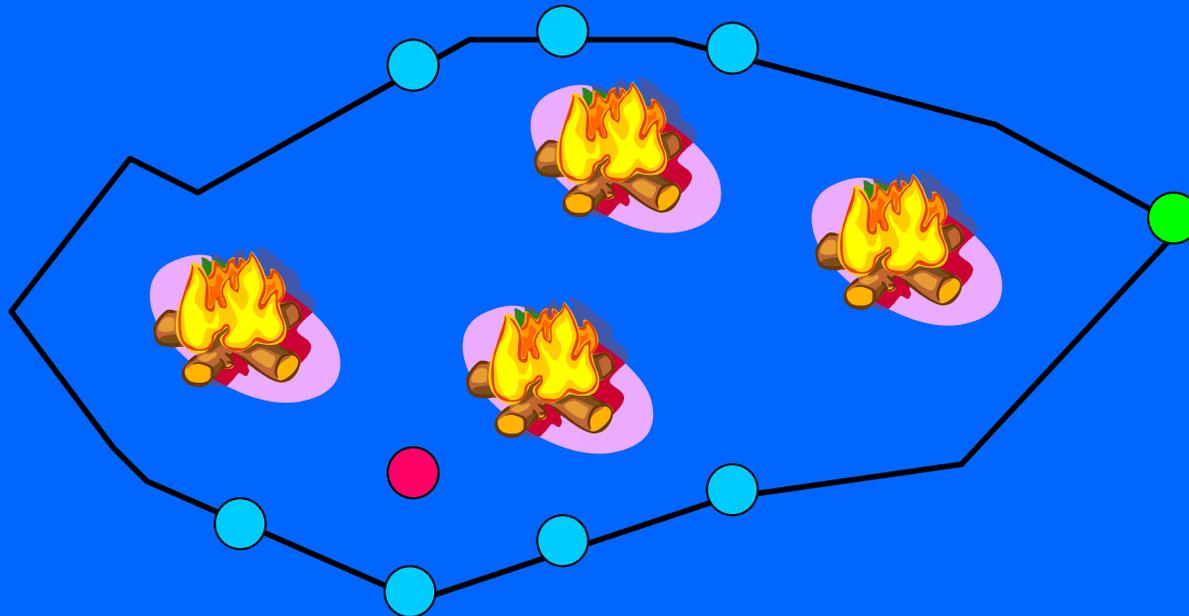


$$L_d$$

Fire starts at only one vertex:

$d = 1$: Trivial.

$d = 2$: Impossible to contain the fire with 1 firefighter per time step

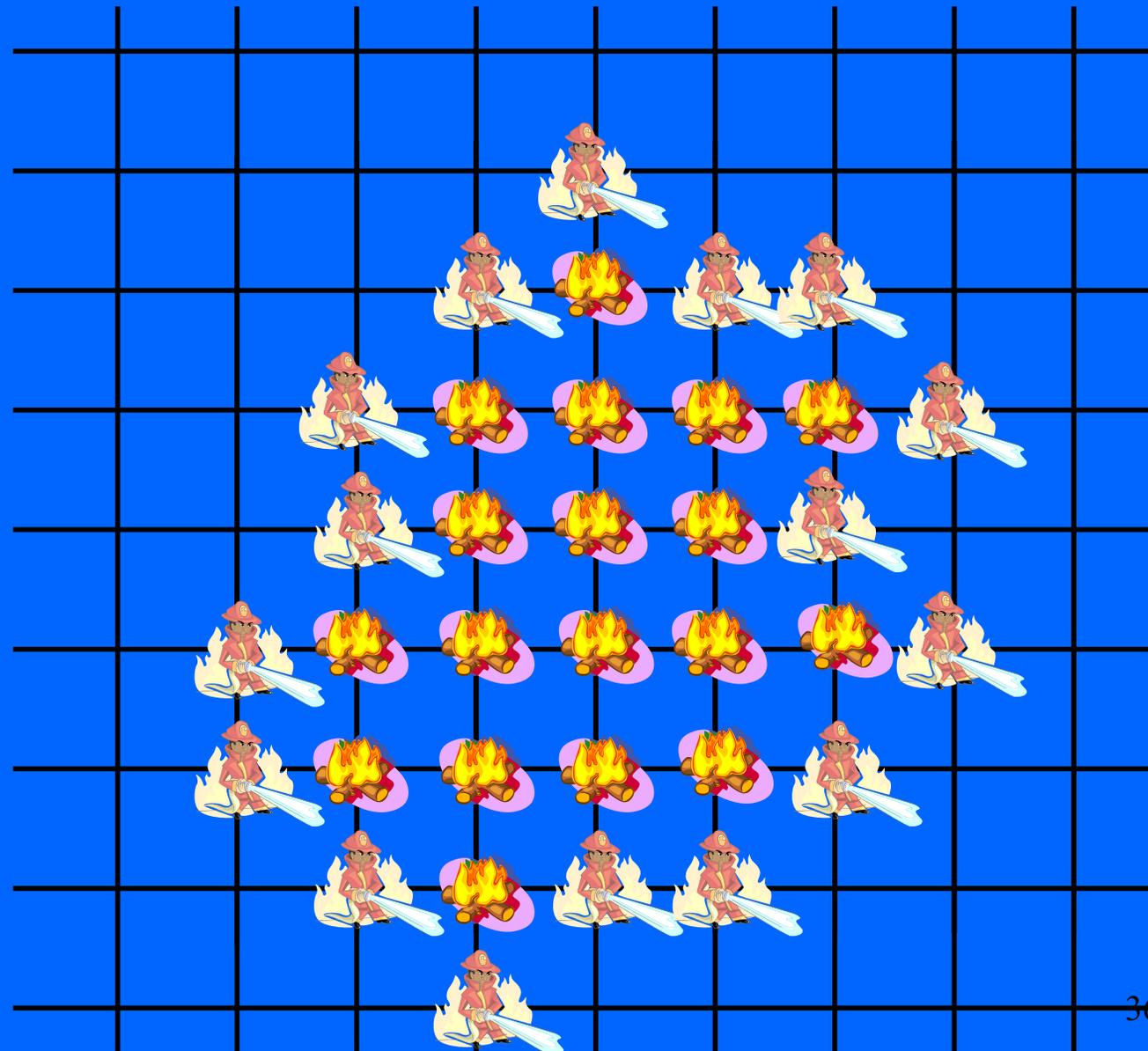


Containing Fires in Infinite Grids L_d

$d = 2$: Two firefighters per time step needed to contain the fire.

8 time steps

18 burnt vertices



Containing Fires in Infinite Grids L_d

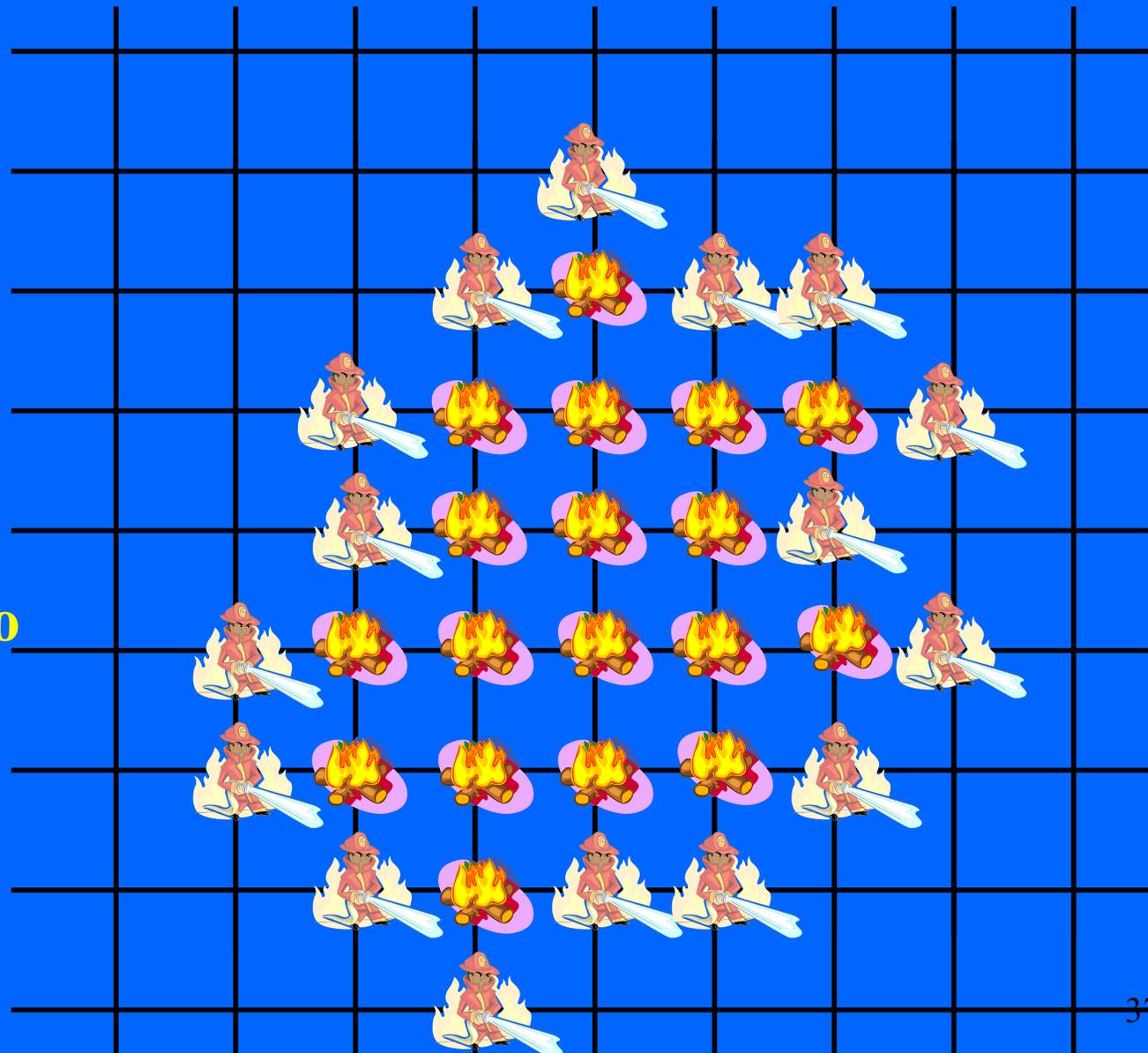
$d = 2$: Two firefighters per time step needed to contain the fire.

8 time steps

18 burnt vertices

**Develin &
Hartke: cannot do
better than 18**

**Wang & Moeller:
Cannot contain
fire in < 8 steps**



Containing Fires in Infinite Grids L_d

Sample Result:

$d \geq 3$: In L_d , every vertex has $2d$ neighbors.

Thus: $2d-1$ firefighters per time step are sufficient to contain any outbreak starting at a single vertex.

Theorem (Develin and Hartke): If $d \geq 3$, $2d - 2$ firefighters per time step are not enough to contain an outbreak in L_d .

Thus, $2d - 1$ firefighters per time step is the minimum number required to contain an outbreak in L_d and containment can be attained in 2 time steps.

More Realistic Models

- You stay in the infected state (state \bullet) for d time periods after entering it and then go back to the uninfected state (state \circ).
- We vaccinate a person in state \circ once $k-1$ neighbors are infected (in state \bullet).
- What if you only get infected with a certain probability if you meet an infected person?
- What if vaccines only work with a certain probability?
- What if the amount of time you remain infective exhibits a probability distribution?

Example III: Cascading Outages in the Power Grid

- Today's electric power systems operate under considerable uncertainty.
- Cascading failures can have dramatic consequences.

Blackout

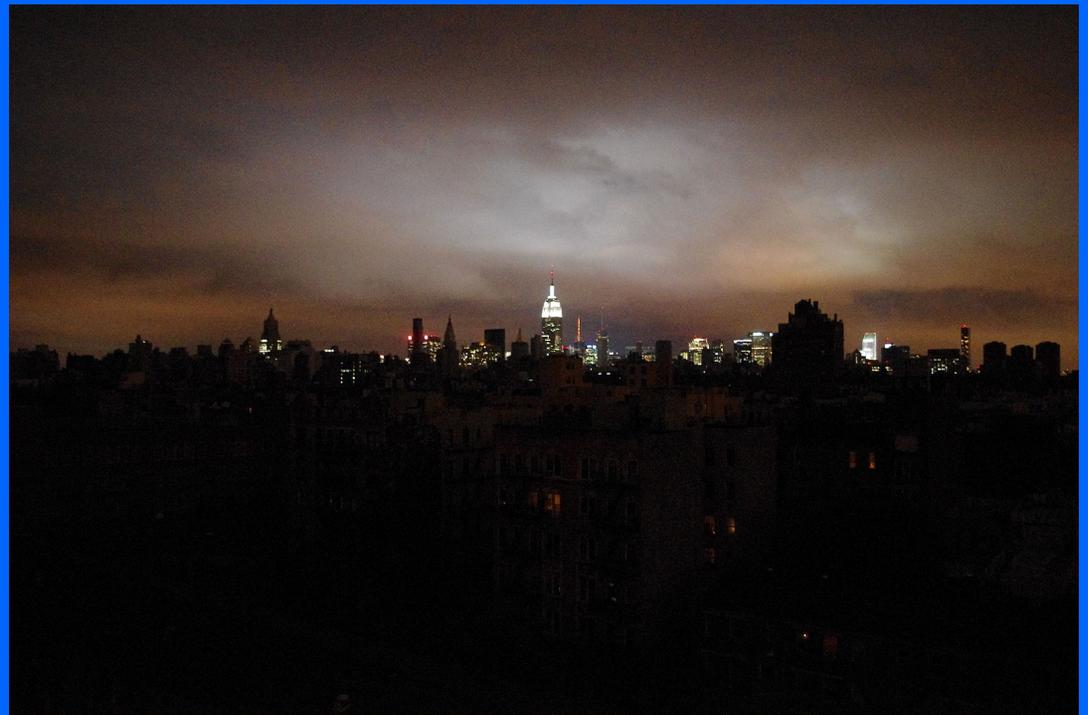


Image credit:

Wikimedia commons; David Shankbone no changes

Cascading Outages in the Power Grid

Grid Resilience:

- How can we design “control” procedures so that the power grid can quickly and efficiently respond to disturbances and quickly be restored to its healthy state?
- Grid disruptions can cascade so fast that a human being may not be able to react fast enough to prevent the cascading disaster leading to a major blackout.
- We are dependent on rapid response through algorithms.

Cascading Outages in the Power Grid

Grid Resilience:

- We are dependent on rapid response through algorithms.
- Need fast, reliable algorithm to respond to a detected problem.
 - Should not necessarily require human input
 - Has to be able to handle multiple possible “solutions”
 - Has to be able to understand what to do if all possible solutions are “bad”

Cascading Outages in Power Grid

Grid Resilience:

- Tool of interest: cascade model of Dobson, et al.
 - An initial “event” takes place
 - Reconfigure demands and generator output levels
 - New power flows are instantiated
 - The next set of faults takes place according to some stochastic model

Cascading Outages in Power Grid

Grid Resilience:

- The power grid model is not the same as a disease-spread model.
- Energy flows from generators through power lines (edges in the power grid graph).
- Each edge has a maximum capacity.
- When a vertex (substation) or edge (transmission line) outage occurs, power reroutes according to physical laws (Kirchhoff's Law, Ohm's Law).

Cascading Outages in Power Grid

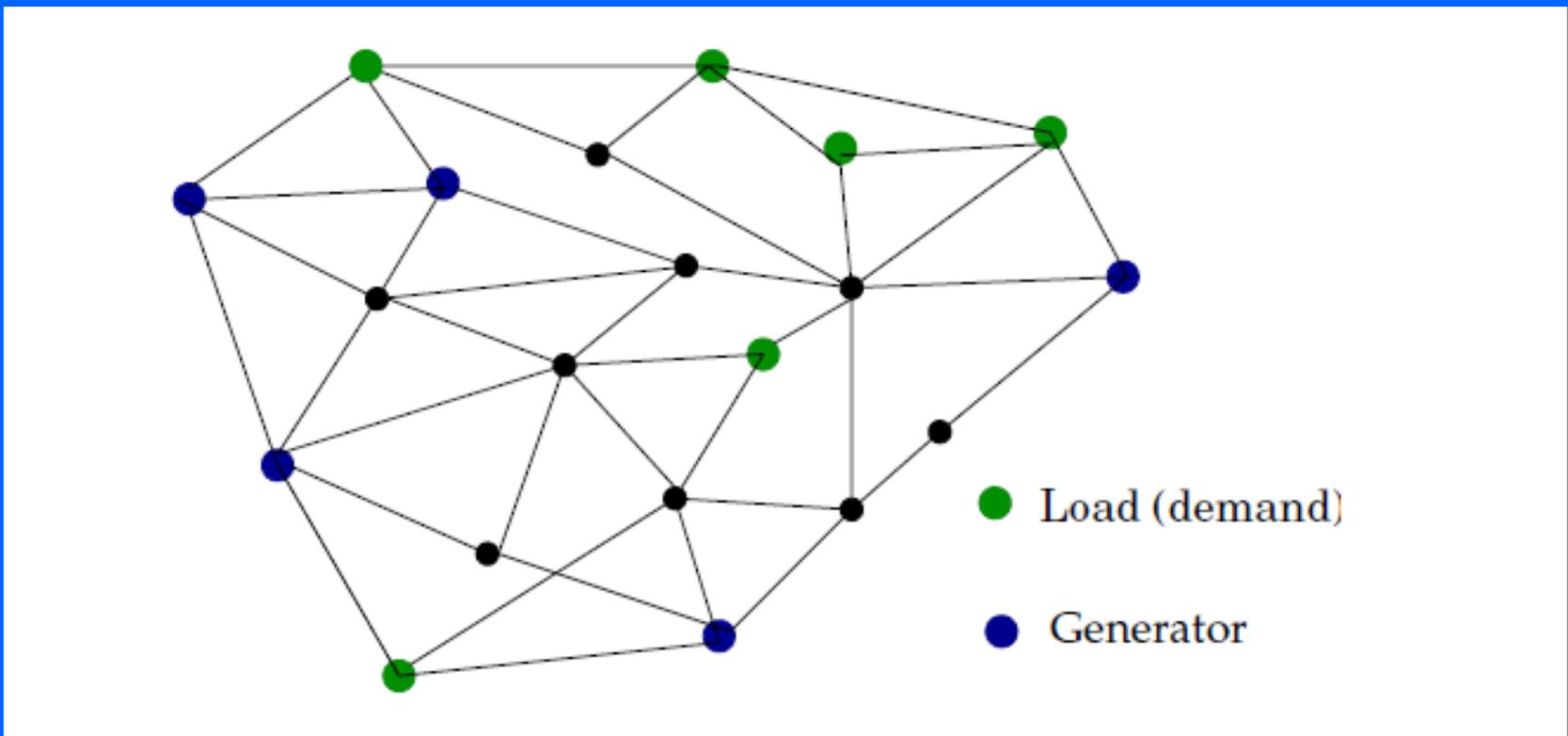
Grid Resilience:

- Because of the rerouting, flows on parallel paths are increased.
- This could cause an overload in a distant transmission line.
- So failures can take place non-locally.

Cascading Outages in the Power Grid

Grid Resilience:

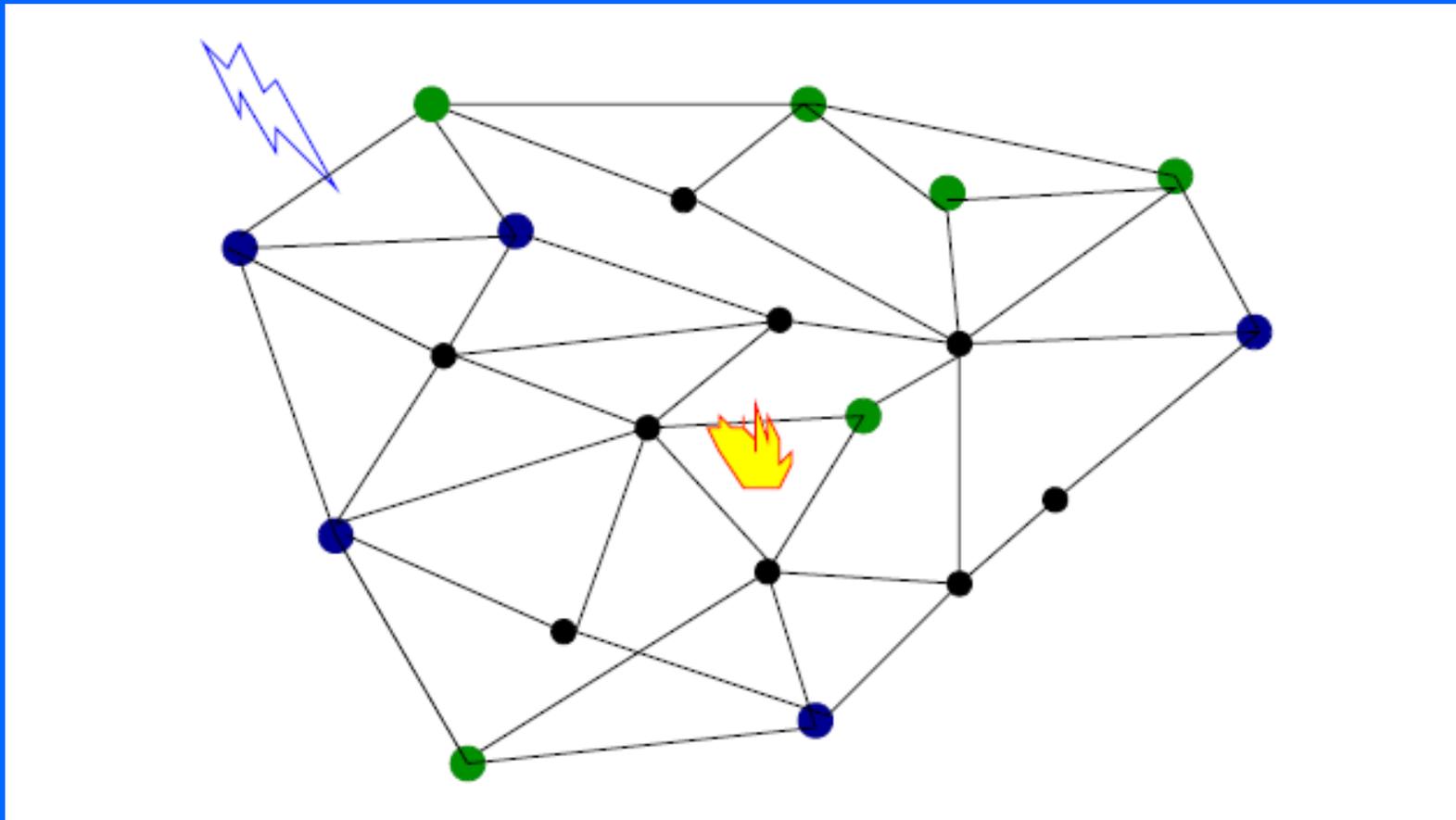
Cascade Model (Dobson, et al.)



Cascading Outages in the Power Grid

Grid Resilience:

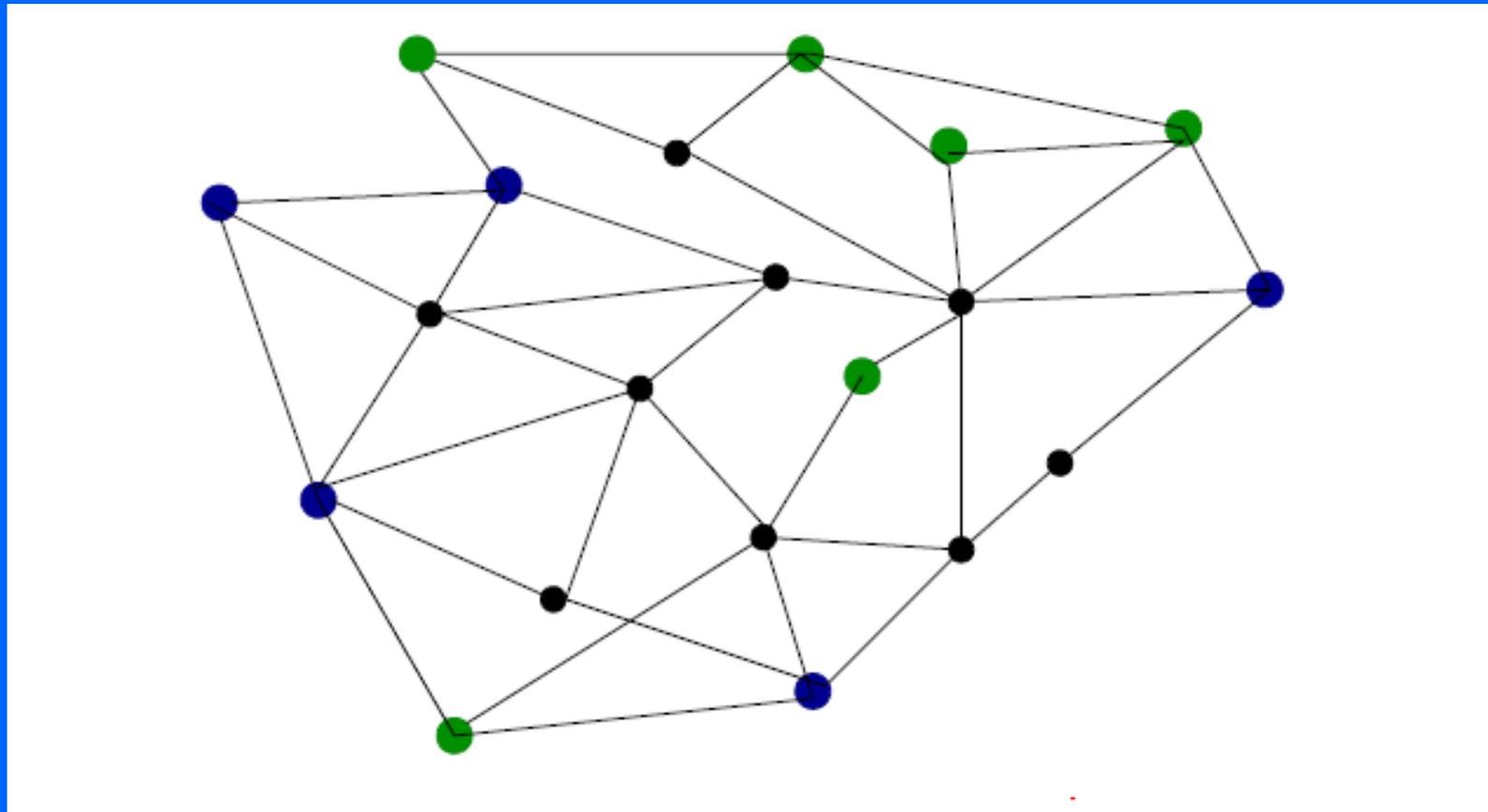
Cascade Model (Dobson, et al.)



Cascading Outages in the Power Grid

Grid Resilience:

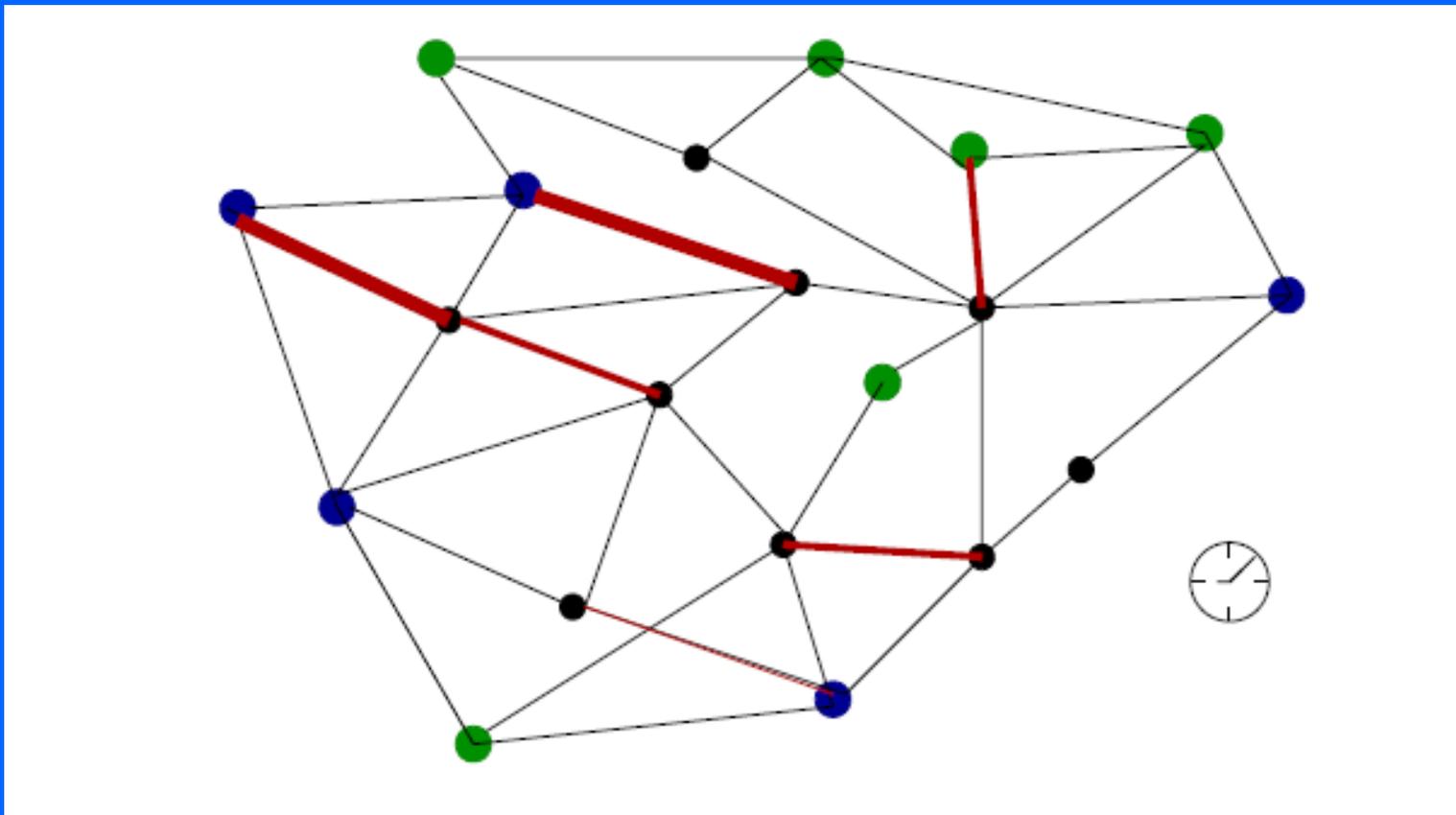
Cascade Model (Dobson, et al.)



Cascading Outages in the Power Grid

Grid Resilience:

Cascade Model (Dobson, et al.)



Credit: Daniel Bienstock

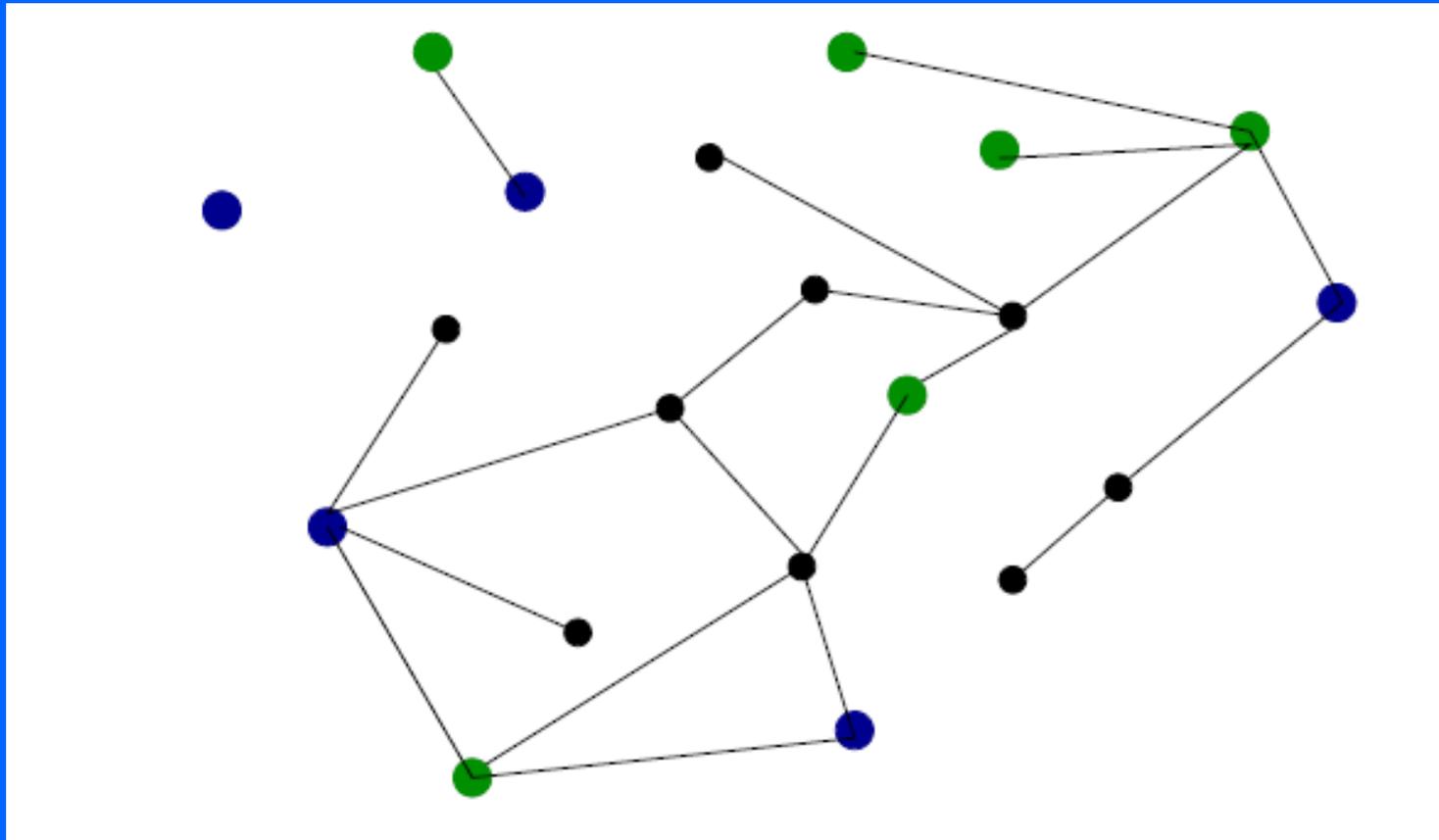
Increased flows on some lines



Cascading Outages in the Power Grid

Grid Resilience:

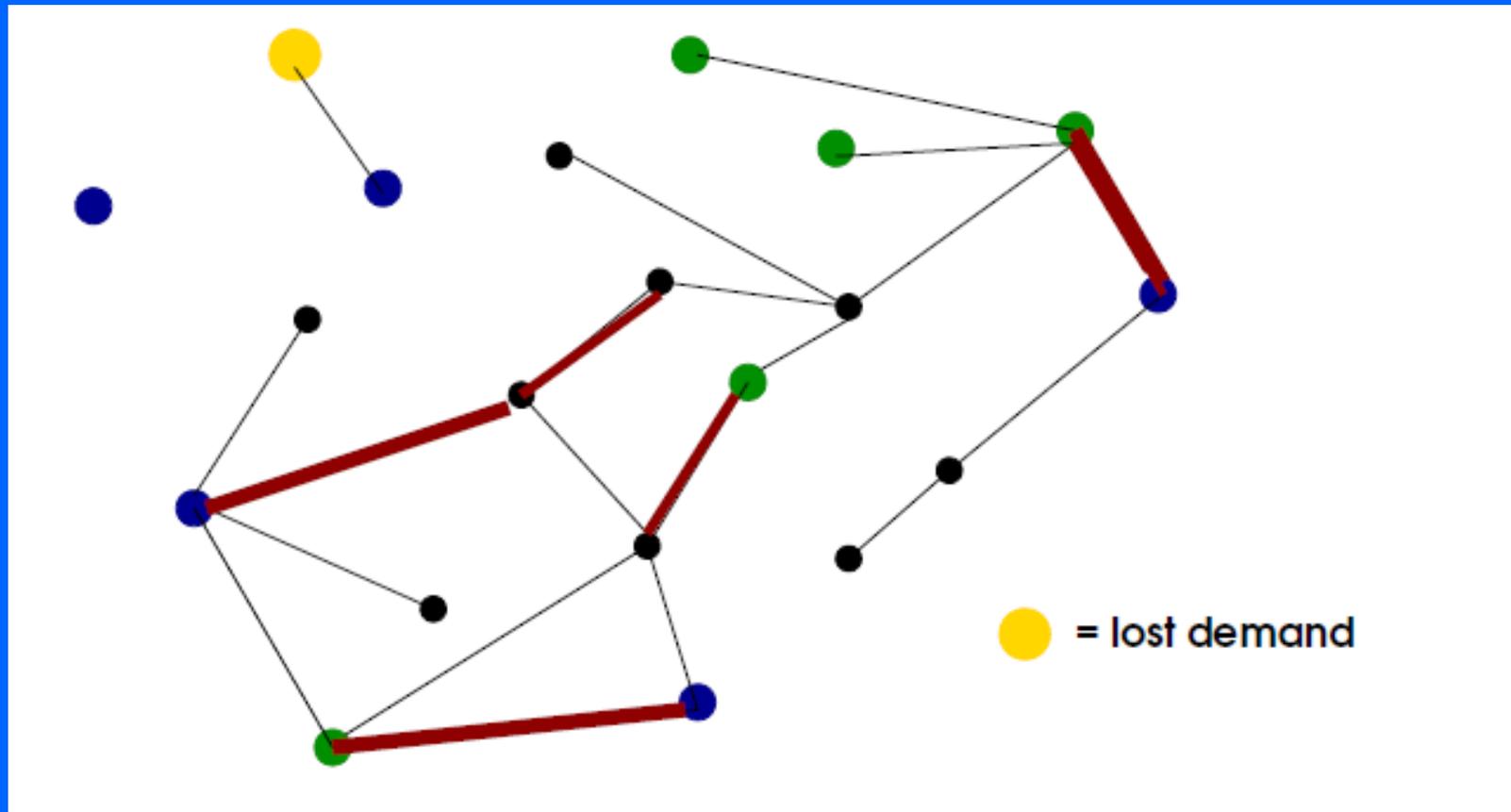
Cascade Model (Dobson, et al.)



Cascading Outages in the Power Grid

Grid Resilience:

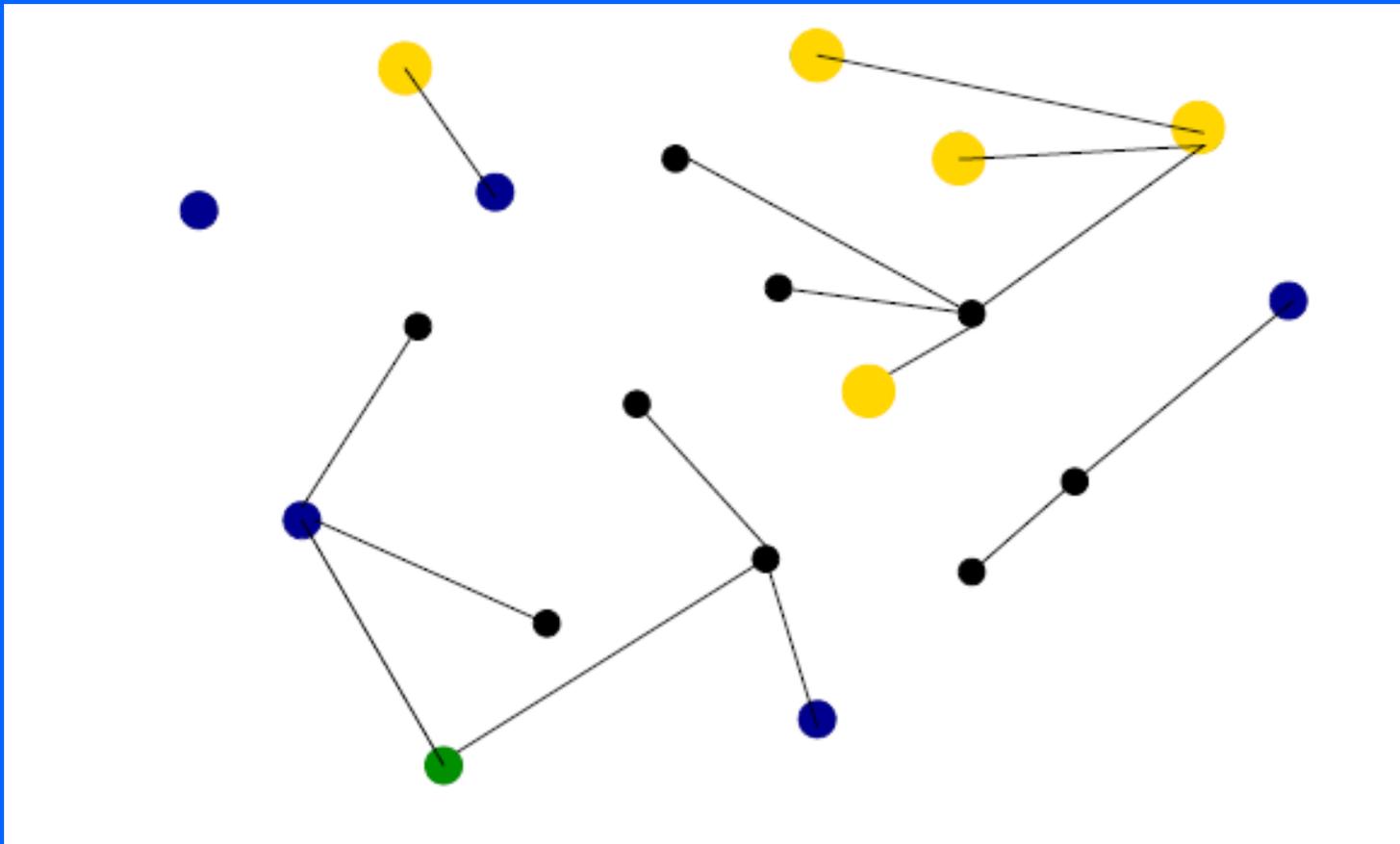
Cascade Model (Dobson, et al.)



Cascading Outages in the Power Grid

Grid Resilience:

Cascade Model (Dobson, et al.)



Cascading Outages in the Power Grid

Grid Resilience:

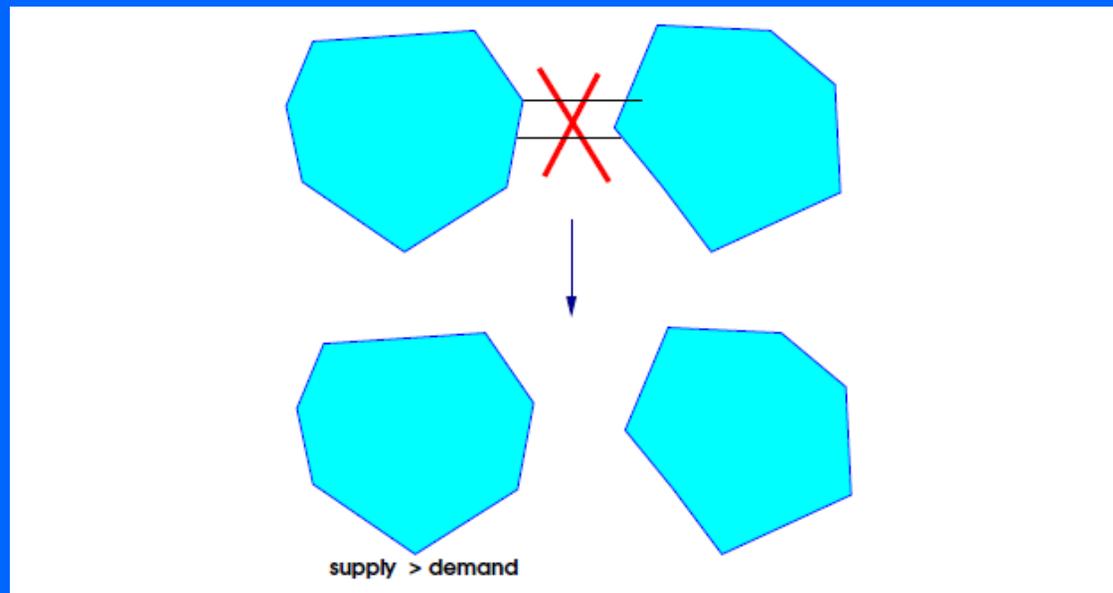
- Cascade model of Dobson, et al.: Exercising “Control”
 - An initial “event” takes place
 - Reconfigure demands and generator output levels
 - New power flows are instantiated
 - Instead of waiting for the next set of faults to take place according to some stochastic process, use the cascade model to learn how to:
 - Take measurements and apply control to shed demand.
 - Reconfigure generator outputs; get new power flows

Cascading Outages in the Power Grid

Grid Resilience:

Cascade Model (Dobson, et al.)

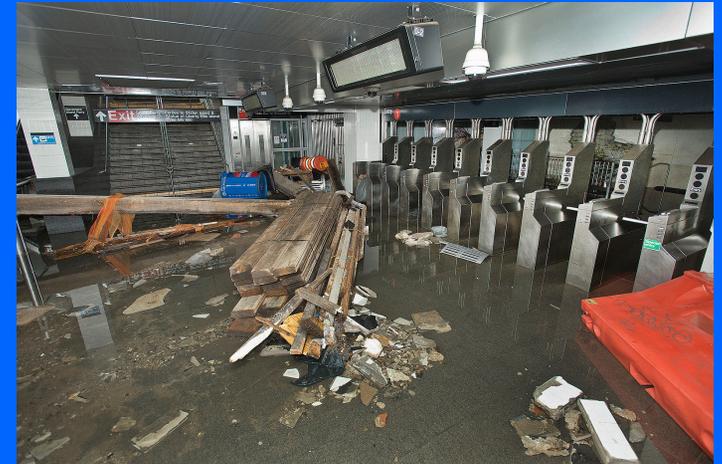
- Use Model to Learn how Best to Create Islands to Protect Part of the Grid
- Hopefully the islands are small and in the rest of the grid, supply > demand.



Example IV: Infrastructure Resilience

- Critical infrastructure systems include:

- Transportation systems
- Telecom
- Water supply systems
- Wastewater systems
- Electric power systems



- After a disruption, system begins to restore service until returning to performance level *at or below* the level before the disruption.

Image credit: Metropolitan Transportation Authority of the State of New York via Wikimedia commons, no changes made.

Example IV: Infrastructure Resilience



Hurricane Sandy, NJ

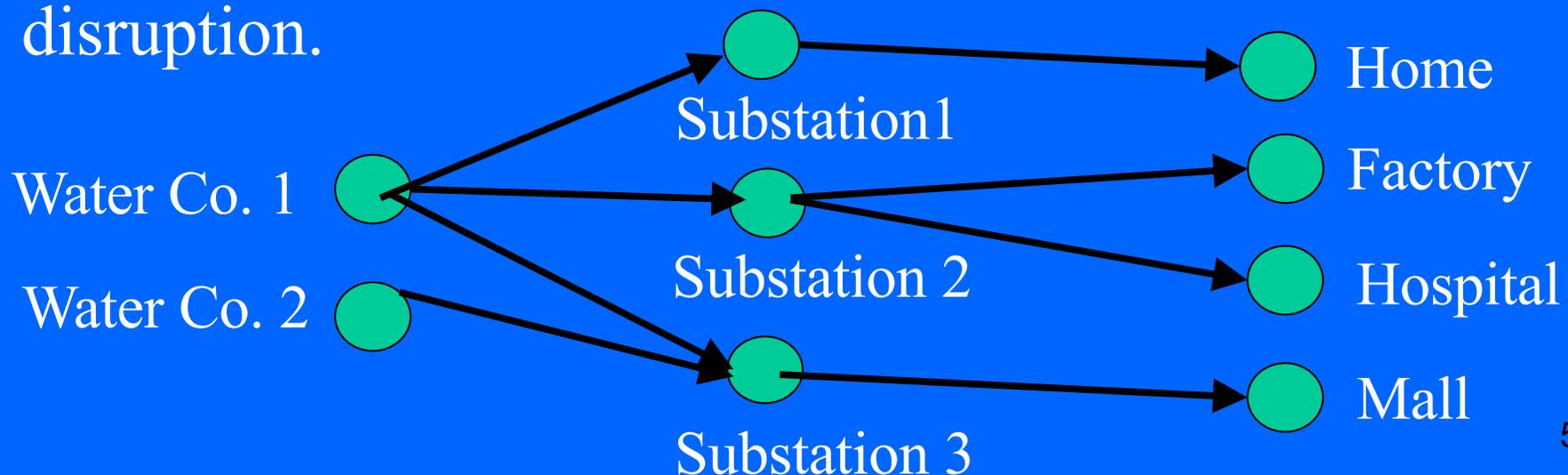
- The following models were developed by Sharkey and Pinkley (2019).
- Service is modeled by flows in networks.

Infrastructure Resilience

- Network now has vertices and directed edges (called *arcs*).
- Flow can only go from vertex i to vertex j along an arc directed from i to j .
- Vertices represent:
 - Components that generate services (supply vertices)
 - Alter the routes of the services (transshipment vertices)
 - Consume services (demand vertices)
- Arcs move the services from one vertex to another.

Infrastructure Resilience

- Example: Water supply system
 - Supply vertices = water companies
 - Transshipment vertices = substations
 - Demand vertices are at households, factories, hospitals, malls, etc.
- Pipes are the arcs, and water is the flow.
- Meeting as much demand as possible is modeled as the classical maximum flow problem – both before and after a disruption.



Infrastructure Resilience

Maximum Flow Problem

- Consider a network $G = (V, A)$
- $V =$ set of vertices, $A =$ set of arcs.
- The arc i to j has a *capacity* u_{ij} .
- Fix one supply vertex s and one demand vertex t .
- There is a *supply* $A(s)$ at s and a *demand* $B(t)$ at t .
- We seek to assign a *flow* x_{ij} to the arc from i to j .
- The flow along that arc must at most the capacity:

$$x_{ij} \leq u_{ij}.$$

Infrastructure Resilience

Maximum Flow Problem

- *Flow conservation*: the sum of flows on arcs into a vertex =- the sum of flows out of the vertex.
- If $A(i) = \{j: (i,j) \in A\}$, then this says:
$$\sum_{j \in A(i)} x_{ij} = \sum_{j: i \in A(j)} x_{ji}$$
- The total flow out of s cannot exceed the supply $A(s)$ and the total flow into t cannot exceed the demand $B(t)$.
- We seek to maximize the total flow that reaches t .

Infrastructure Resilience

Maximum Flow Problem

The *Maximum Flow Problem* seeks to determine the largest amount of flow that can reach t while:

- Keeping the flow on each arc at most the capacity
- Not exceeding total supply and demand
- Satisfying the flow conservation requirement at each vertex.

Infrastructure Resilience

Maximum Flow Problem

- The famous augmenting path algorithm (Ford-Fulkerson Algorithm) finds the maximum flow.
- Note: the maximum flow problem is a simplification.
- It assumes that there are no other constraints on flow.
- This might apply to supply chain networks:
 - E.g., physical goods move through intermediate warehouses and distribution centers.

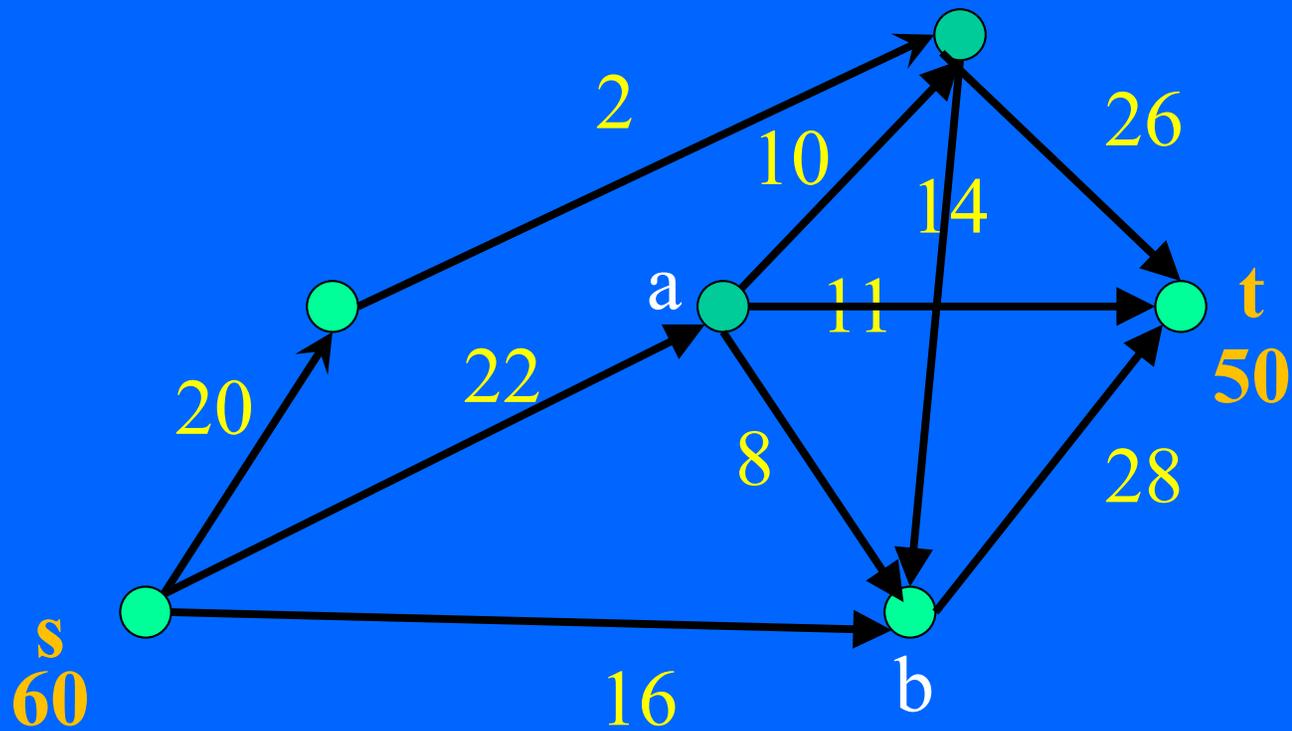
Infrastructure Resilience

Maximum Flow Problem

- For more complicated infrastructure, there are things like physical laws offering additional constraints.
- Example: Kirchhoff's and Ohm's Laws for power grid networks.
- Example: water distribution networks involve constraints involving the relation between flow of water and pressure.

Infrastructure Resilience

Maximum Flow Problem



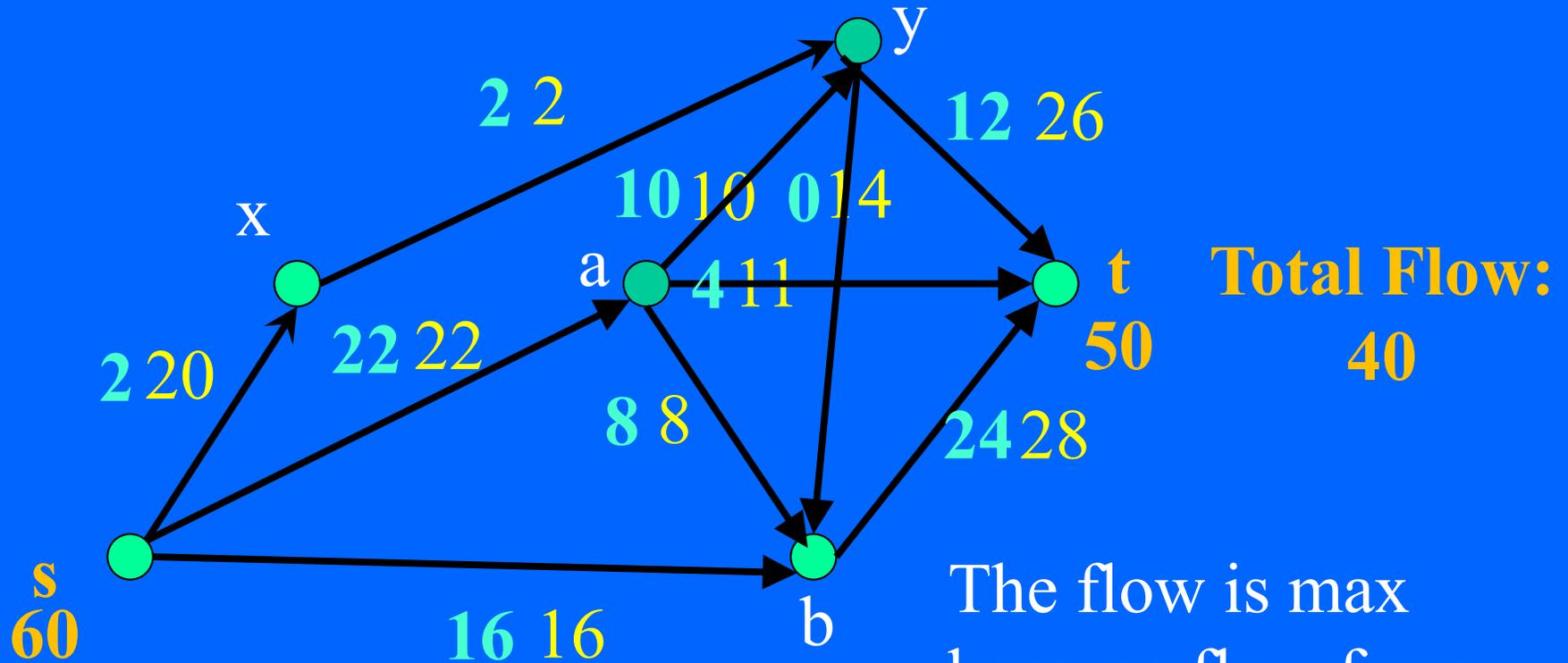
Key:

Orange: supply and demand

Yellow: Capacity

Infrastructure Resilience

Maximum Flow Problem



Key:

Orange: supply and demand

Yellow: Capacity

Green: Flow

The flow is max because flow from x to y is at most 2, so s to x is at most 2, so out of s is at most 40.

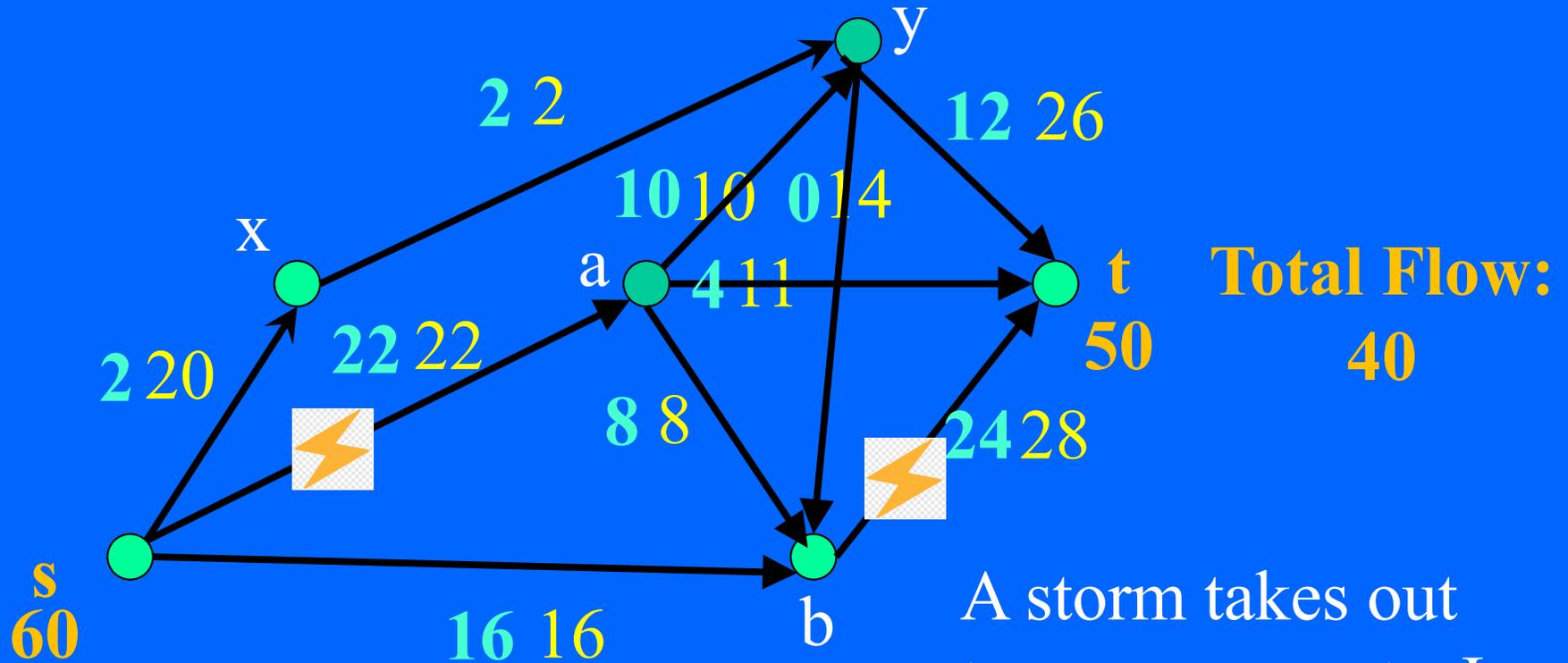
Infrastructure Resilience

Maximum Flow Problem

- If some of the arcs are destroyed, in what order should we reopen them?
- One goal: get closest to original maximum flow as early as possible.

Infrastructure Resilience

Maximum Flow Problem



Key:

Orange: supply and demand

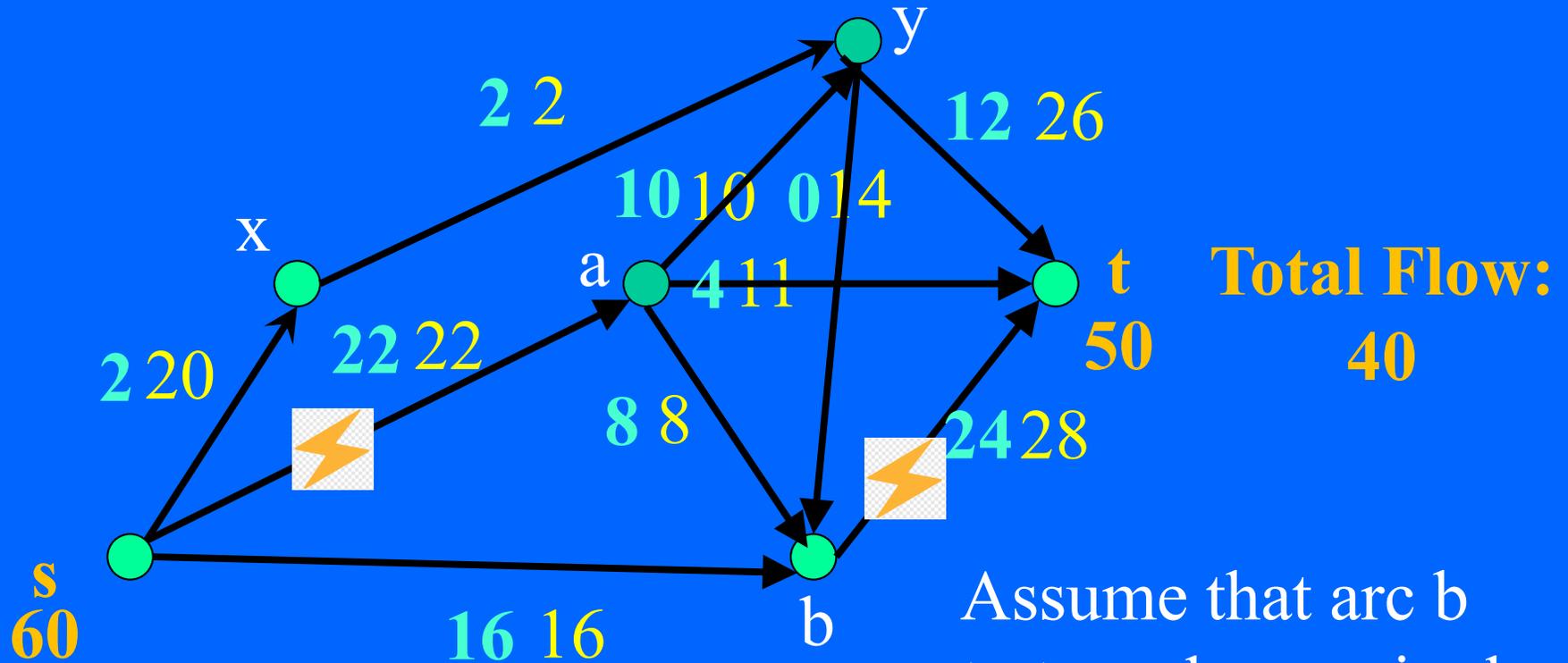
Yellow: Capacity

Green: Flow

A storm takes out two components. In what order should we bring them back online?

Infrastructure Resilience

Maximum Flow Problem



Key:

Orange: supply and demand

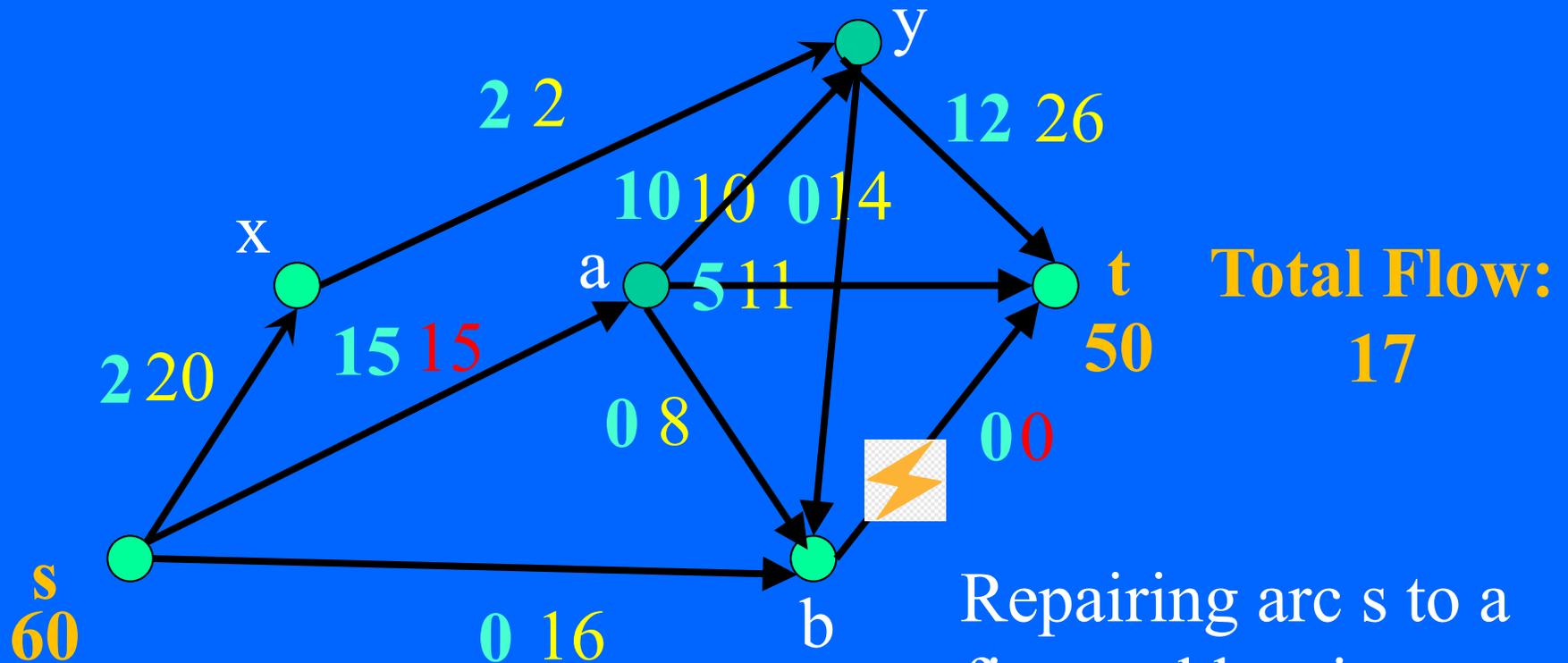
Yellow: Capacity

Green: Flow

Assume that arc b to t can be repaired fully but arc s to a can only be repaired to a lower capacity of 15.

Infrastructure Resilience

Maximum Flow Problem



Key:

Orange: supply and demand

Yellow: Capacity

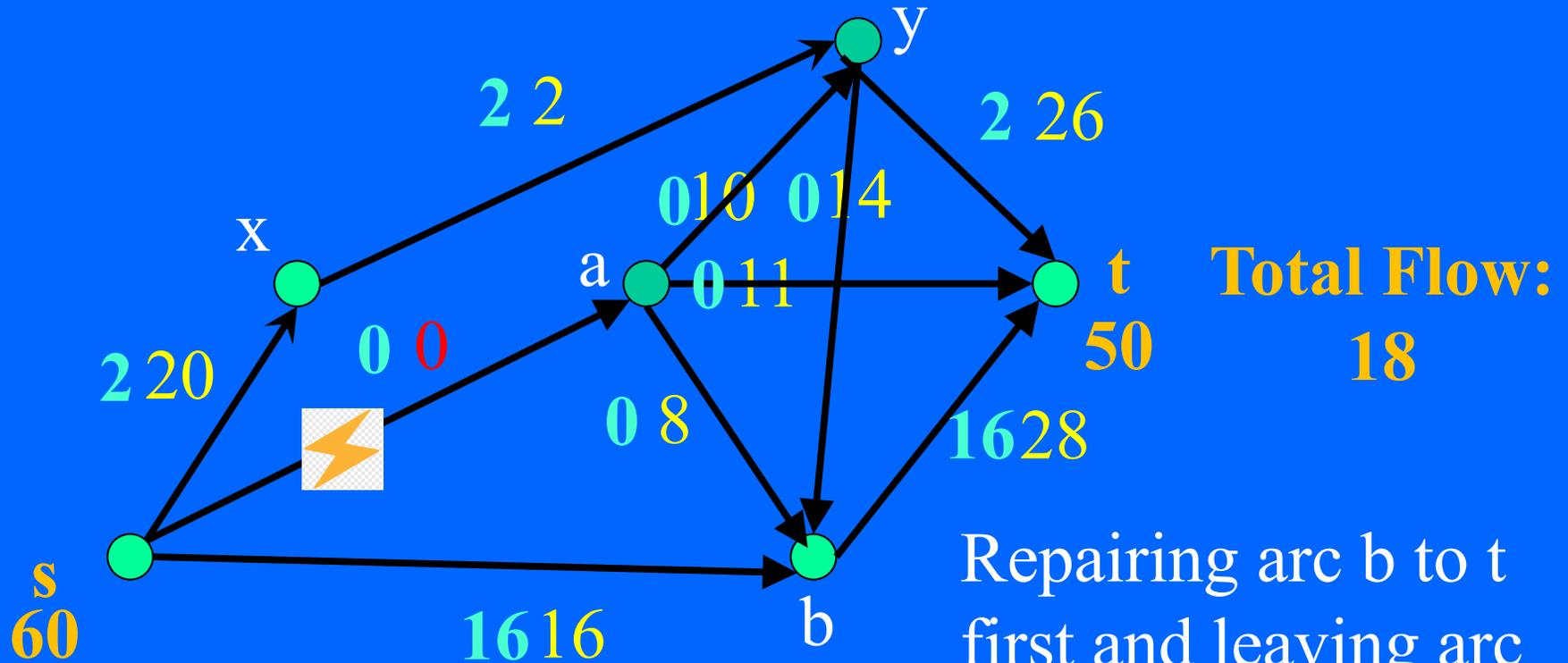
Green: Flow

Red: Reduced Capacity

Repairing arc s to a first and leaving arc b to t unrepaired gives a max flow of 17.

Infrastructure Resilience

Maximum Flow Problem



Key:

Orange: supply and demand

Yellow: Capacity

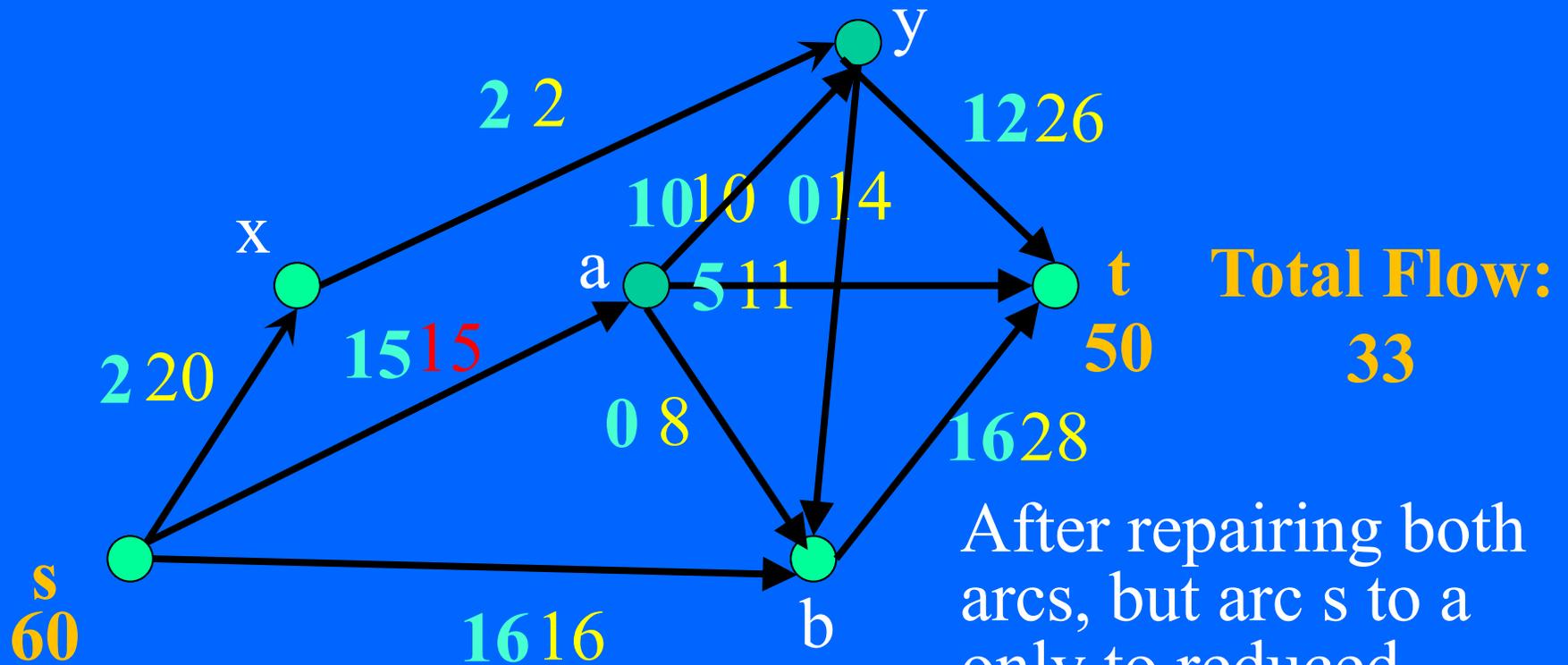
Green: Flow

Red: Reduced Capacity

Repairing arc b to t first and leaving arc s to a unrepaired gives a max flow of 18. So repair this arc first.

Infrastructure Resilience

Maximum Flow Problem



Key:

Orange: supply and demand

Yellow: Capacity

Green: Flow

Red: Reduced Capacity

After repairing both arcs, but arc s to a only to reduced capacity, means we only regained a reduced max flow of 33 – didn't fully restore flow of 40

Infrastructure Resilience

Maximum Flow Problem

- We made the simplifying assumption that there was one supply vertex and one demand vertex.
- In practice, there are many supply vertices s_1, s_2, \dots , and demand vertices t_1, t_2, \dots , with supply $A(s_i)$ at s_i and $B(t_i)$ at t_i .
- But we can reduce this to a single supply and demand vertex by adding a supply vertex S with supply $A(S) = \sum A(s_i)$ and an arc from S to each s_i with capacity $A(s_i)$ and a demand vertex T with demand $B(T) = \sum B(t_i)$ and an arc from each t_i to T with capacity $B(t_i)$.

Infrastructure Resilience

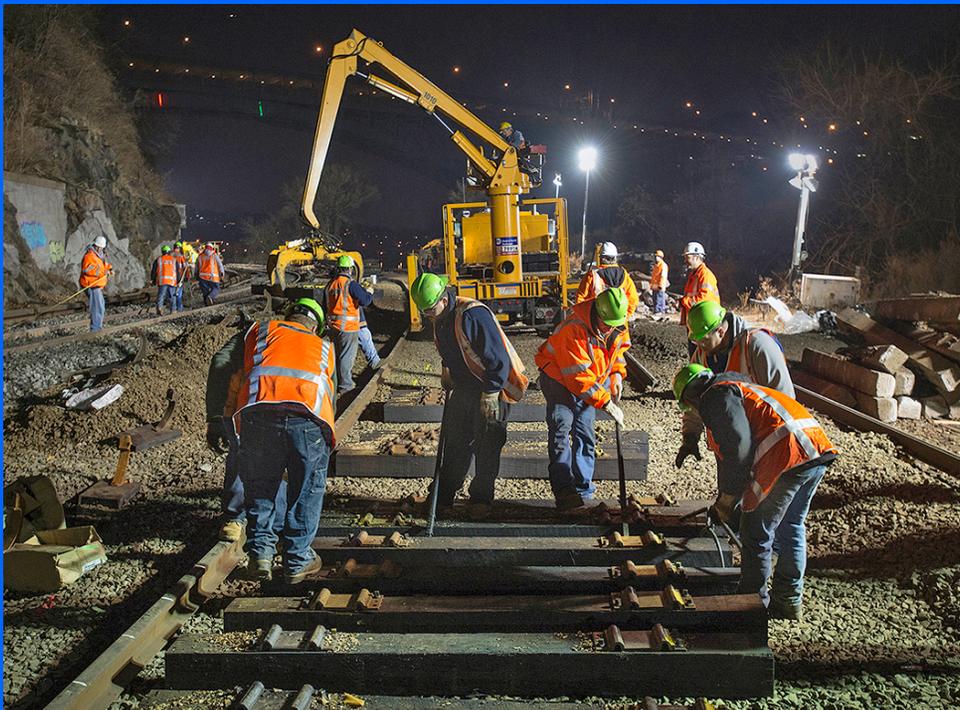
Maximum Flow Problem

- As different components of a network are repaired (to the extent possible), the maximum flow increases.
- How far off it is from the original max flow when repairs are done is one metric for resilience.
- How long it takes to complete the repairs is another metric for resilience.
- We turn next to the repair process.

Infrastructure Resilience

The Repair Process

- A different approach to reopening damaged components uses the theory of machine scheduling.
- After a disruption, repairs are made so services can be restored.
- Repairs use scarce *resources*: work crews, equipment.



Credit: Patrick Cashin / MTA. Edited and cropped slightly by Daniel Case; via Wikimedia commons.

Infrastructure Resilience

The Repair Process

- Simplifying assumption: can only repair one component at a time (one vertex or arc).
- Need a schedule for when a resource is repairing a component.
- In the scheduling literature, we talk about jobs on a set of machines, and processing them.
- Jobs* here correspond to damaged components.
- Machines* correspond to work crews.

Infrastructure Resilience

The Repair Process

- Each job (damaged component) k has a different level of *importance* w_k .
- Each job also has a *duration* p_k .
- In the scheduling literature, each job k is assigned to a machine (work crew) m_k .
- The jobs assigned to a machine (work crew) m are given an order.
- So the *completion time* C_k of job k is the sum of the durations of all jobs assigned to the machine (crew) m_k that precede job k plus the duration of job k .

Infrastructure Resilience

The Repair Process

- There are various objectives for a good repair schedule.
- One is to minimize the weighted average completion time over all jobs, with the weight measuring the importance of the job.

$$\min \sum_k w_k C_k$$

- This is sometimes called the *restoration performance*.
- If there is just one work crew, a greedy algorithm minimizes this: Repair component k in non-increasing order of the ratio w_k/p_k .

Infrastructure Resilience

The Repair Process

- A similar algorithm works if there many machines but each has the same processing time for repairing a given component.
- However, in general, most such scheduling problems are hard: NP-hard.

Infrastructure Resilience

Repairing Multiple Interdependent Infrastructures

- In a complex city, there are many infrastructures.
- They have interdependencies.
- Examples:
 - A subway (transportation infrastructure) needs power (electrical infrastructure) before it can be reopened.
 - A hospital needs both power and water before it can be reopened.
- This is modeled by studying a collection of networks, one for each infrastructure.
- A given infrastructure cannot operate until there is sufficient level of service (flow) on certain specific vertices in other infrastructures.

Infrastructure Resilience

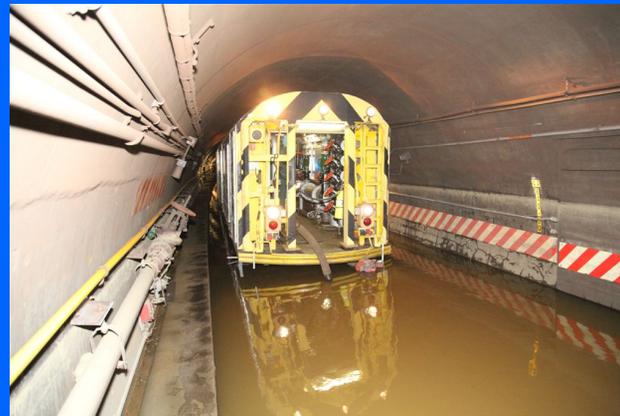
Repairing Multiple Interdependent Infrastructures

- Scheduling repair of different infrastructures will therefore depend on these interdependencies.
- There is a considerable literature on this topic.
- Another complication: interdependencies among repair jobs – sometimes in different infrastructures.

Infrastructure Resilience

Repairing Multiple Interdependent Infrastructures

- Example:
 - To reopen subway lines, you need to repair a line.
 - Once you repair the lines, you need to run a test train on the line to check for safety and quality of the repair.
 - But power to the line must be restored before you can run a test train.



Subway tunnel pump train

Image credits:
Flood: ---=XEON=---
Pump train: Metropolitan
Transportation Authority
of the State of New York

Infrastructure Resilience

Repairing Multiple Interdependent Infrastructures

- Example:
 - Trees bring power lines down on a road.
 - First need to do a safety inspection to make sure it's safe to enter the road.
 - Then clear debris from the road.
 - Then repair downed power lines.



Closing Comment

- I have presented several simple examples of how to generate responses to disruptive events.
- Even these simple examples lead to problems that are “hard” in a precise sense.
- Another approach is to study ways to design graphs or networks so as to make them more resilient in case of disruption.
- That is a topic for another day.