# On Weighted Graphs Yielding Facets of the Linear Ordering Polytope 

Gwenaël Joret

Université Libre de Bruxelles, Belgium

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## Definition

For any finite set $Z$,

- for $R \subseteq Z \times Z$, the vector $x^{R}$ is the characteristic vector of $R$, that is,

$$
x_{i, j}^{R}= \begin{cases}1 & \text { if }(i, j) \in R \\ 0 & \text { otherwise }\end{cases}
$$

- the linear ordering polytope $P_{L O}^{Z} \subset \mathbb{R}^{Z \times Z}$ is

$$
P_{L O}^{Z}=\operatorname{conv}\left\{x^{L}: L \text { linear order on } Z\right\}
$$

## Definition

For a vertex-weighted graph $(G, \mu)$ and $S \subseteq V(G)$,

- $\mu(S):=\sum_{v \in S} \mu(v)$
- $\mathrm{w}(S):=\mu(S)-|E(G[S])|$
(weight of $S$ )
(worth of $S$ )
- $\alpha(G, \mu):=\max _{S \subseteq V(G)} \mathrm{w}(S)$
- $S$ is tight if $\mathrm{w}(S)=\alpha(G, \mu)$

- weight $=4$
- worth $=2$
- tight

Suppose

- $(G, \mu)$ is any weighted graph
- $Y$ is a set s.t. $|Y|=|V(G)|$ and $Y \cap V(G)=\emptyset$
- $f: V(G) \rightarrow Y$ is a bijection
- $Z$ is a finite set s.t. $V(G) \cup Y \subseteq Z$


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## Definition

- The graphical inequality of $(G, \mu)$, which is valid for $P_{L O}^{Z}$, is

$$
\sum_{v \in V(G)} \mu(v) \cdot x_{v, f(v)}-\sum_{\{v, w\} \in E(G)}\left(x_{v, f(w)}+x_{f(v), w}\right) \leq \alpha(G, \mu)
$$

- $(G, \mu)$ is facet-defining if its graphical inequality defines a facet of $P_{L O}^{Z}$


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- $(G, \mu)$ is facet-defining if its graphical inequality defines a facet of $P_{L O}^{Z}$
N.B. $(G, \mu)$ being facet-defining is a property of the graph solely, i.e. it is independent of the particular choice of $Y, f$ and $Z$


## A characterization of facet-defining graphs

Definition

- For any tight set $T$ of $(G, \mu)$, a corresponding affine equation is defined:

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\sum_{v \in T} y_{v}+\sum_{e \in E(T)} y_{e}=\alpha(G, \mu)
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- The system of $(G, \mu)$ is obtained by putting all these equations together


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Theorem (Christophe, Doignon and Fiorini, 2004)
( $G, \mu$ ) is facet-defining $\Leftrightarrow$ the system of $(G, \mu)$ has a unique solution

- Basically rephrases the fact that the dimension of the face of $P_{L O}^{Z}$ defined by the graphical inequality must be high enough
- We lack a 'good characterization' of these graphs...


## A few results

(assuming from now on that all graphs have at least 3 vertices)

## Definition

$G$ is stability critical if $G$ has no isolated vertex and $\alpha(G \backslash e)>\alpha(G)$ for all $e \in E(G)$

Theorem (Koppen, 1995)
$(G, \mathbb{1})$ is facet-defining $\Leftrightarrow G$ is connected and stability critical

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$(G, \mathbb{1})$ is facet-defining $\Leftrightarrow G$ is connected and stability critical
Theorem (Christophe, Doignon and Fiorini, 2004)
$(G, \mu)$ is facet-defining $\Leftrightarrow$ its 'mirror image' $(G, \operatorname{deg}-\mu)$ is facet-defining


## Definition

- The defect of $G$ is $|V(G)|-2 \alpha(G)$

a stability critical graph

$$
\begin{aligned}
& |V(G)|=12 \\
& \alpha(G)=3 \\
& \rightarrow \text { defect }=6
\end{aligned}
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## Definition

- The defect of $G$ is $|V(G)|-2 \alpha(G)$
- The defect of $(G, \mu)$ is $\mu(V(G))-2 \alpha(G, \mu)$

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$$


a facet-defining graph

$$
\begin{aligned}
& \mu(V(G))=7 \\
& \alpha(G, \mu)=2 \\
& \rightarrow \text { defect }=3
\end{aligned}
$$

## Theorem

- The defect $\delta$ of a connected stability critical graph $G$ is always positive (Erdős and Gallai, 1961)
- Moreover, $\delta \geq \operatorname{deg}(v)-1$ for all $v \in V(G) \quad$ (Hajnal, 1965)


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Theorem (Doignon, Fiorini, J.)

- The defect $\delta$ of any facet-defining graph $(G, \mu)$ is positive
- $(G, \mu)$ and $(G, \operatorname{deg}-\mu)$ have the same defect
- For all $v \in V(G)$, we have

$$
\delta \geq \operatorname{deg}(v)-\mu(v) \geq 1
$$

and, because of the mirror image, also

$$
\delta \geq \mu(v) \geq 1
$$

## Odd subdivision

Here is an extension of a classical operation on stability-critical graphs:

odd subdivision

inverse of odd subdivision


Theorem (Christophe, Doignon and Fiorini, 2004)
The odd subdivision operation and its inverse keep both a graph facet-defining. Moreover, the defect does not change

## Lemma

An inclusionwise minimal cutset of a facet-defining graph cannot span "○" or "O-○"

Thus when we have $\stackrel{1}{1-1}^{(1)}$ we can always contract both edges by using the inverse of odd subdivision operation

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## Definition

A facet-defining graph is minimal if no two adjacent vertices have degree 2

## Classification of stability critical graphs

Theorem (Lovász, 1978)
For every positive integer $\delta$, the set $\mathcal{S}_{\delta}$ of minimal connected stability critical graphs with defect $\delta$ is finite

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## Research problem

Is there a finite number of minimal facet-defining graphs with defect $\delta$, for every $\delta \geq 1$ ?

- It turns out to be true for $\delta \leq 3$
$\rightarrow$ an overview of the proofs is given in the next few slides
- The problem is wide open for $\delta \geq 4$

Notice first that the only minimal facet-defining graph with defect
$\delta=1$ is (1), because $\delta \geq \mu(v) \geq 1$

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Let's look at another operation:

subdivision of a star


Theorem
The subdivision of a star operation keeps a graph facet-defining.
Moreover, the defect does not change

## Definition

$\left(G_{1}, \mu_{1}\right)$ and $\left(G_{2}, \mu_{2}\right)$ are equivalent if one can be obtained from the other by using the

- odd subdivision
- inverse of odd subdivision
- subdivision of a star
operations finitely many times.


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operations finitely many times.
Notice
- two equivalent graphs have the same defect
- $(G, \mu)$ and $(G, \operatorname{deg}-\mu)$ are equivalent:



## Facet-defining graphs with defect 2

Recall

$$
\left\{\begin{array}{l}
\delta \geq \mu(v) \geq 1 \\
\delta \geq \operatorname{deg}(v)-\mu(v) \geq 1
\end{array}\right.
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for any vertex $v$ of a facet-defining graph with defect $\delta$
$\Rightarrow \operatorname{deg}(v) \leq 2 \delta$

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$\operatorname{deg}(v) \leq 2 \delta-1$ for any vertex $v$ of a facet-defining graph with defect $\delta \geq 2$

Thus, every vertex of a facet-defining graph with defect 2 is either $\overbrace{\text { or }}^{1}$

$\Rightarrow$ Any facet-defining graph with defect 2 is equivalent to some stability critical graph

Theorem (Andrásfai, 1967)
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$\rightarrow$ we derive:
Theorem
There are exactly five minimal facet-defining graphs with defect 2 :


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By previous bounds, any vertex falls in one of these cases when $\delta=3$ :


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- We would like to show that the number of vertices $v$ of $(G, \mu)$ with $\operatorname{deg}(v) \geq 3$ is bounded by some absolute constant


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- By the subdivision of a star operation, w.l.o.g. $\nexists(2,3)$-, $(3,4)$-, or $(3,5)$-vertices in $(G, \mu)$
- Main issue: how to get rid of the $(2,4)$-vertices and $(2,5)$-vertices?

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$\rightarrow$ defines a graph on the neighborhood $N(v)$ of $v$, denoted $H_{v}$ :


## Expanding a vertex

Assume $\exists a, b, c, d \in V\left(H_{v}\right)$ s. t. $\{a, b\} \in E\left(H_{v}\right)$ and $\{c, d\} \notin E\left(H_{v}\right)$

expanding $v$


## Lemma

- Expanding v keeps $(G, \mu)$ facet-defining and does not change the defect
- Any $(2,5)$-vertex of $(G, \mu)$ is expandable


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## Lemma

- Expanding $v$ keeps $(G, \mu)$ facet-defining and does not change the defect
- Any $(2,5)$-vertex of $(G, \mu)$ is expandable
$\rightarrow$ w.l.o.g. $(G, \mu)$ has no expandable vertices, as expanding a vertex increases the number of vertices with degree at least 3


## Splitting a vertex

Suppose that $v$ is a (2,4)-vertex and that $\{a, b\},\{c, d\} \notin E\left(H_{v}\right)$

splitting $v$



Lemma

- Splitting $v$ keeps $(G, \mu)$ facet-defining and does not change the defect
- Every nonexpandable (2,4)-vertex is splittable

Assume now that $v$ is a nonexpandable (2,4)-vertex. As $v$ is splittable, $H_{v}$ is isomorphic to one of these 3 graphs:


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$v$ is "thin"
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- $(G, \mu)$ has at most 5 thick vertices

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- $(G, \mu)$ has at most 5 thick vertices
$\rightarrow$ it remains to show that $(G, \mu)$ has not too many thin vertices...

Key lemma
$(G, \mu)$ has at most $\frac{3}{2} N$ thin vertices, where $N$ is the number of vertices with weight 1 and degree at least 3

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- So, the number of vertices with degree at least 3 in $(G, \mu)$ is at most $N+\frac{3}{2} N+5=\frac{5}{2} N+5 \leq \frac{5}{2} c+5$

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Thus we obtain:
Theorem
There is a finite number of minimal facet-defining graphs with defect 3


## As a (brief) conclusion

Graphical inequalities for the linear ordering polytope give rise to a new family of weighted graphs with interesting structural properties

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Thank you!

