

# Sieving for shortest lattice vectors using fast search algorithms

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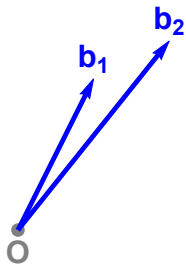
# Lattices

What is a lattice?



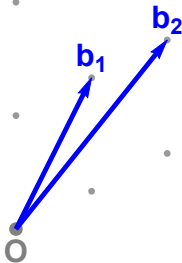
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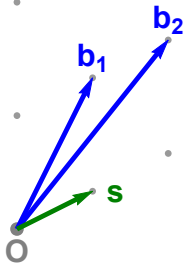
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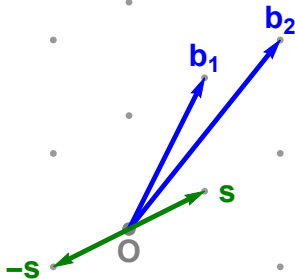
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## Shortest Vector Problem (SVP)



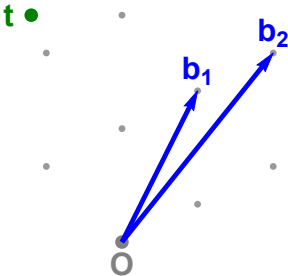
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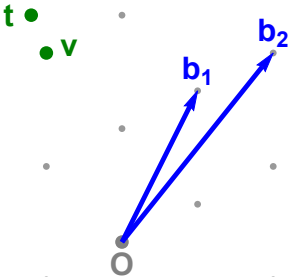
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- “Constructive cryptography”: Lattice-based cryptosystems
  - ▶ Based on hard lattice problems (SVP, CVP, LWE, SIS)
  - ▶ NTRU cryptosystem [HPS98, ...]
  - ▶ Fully Homomorphic Encryption [Gen09, ...]
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How hard are hard lattice problems such as SVP?

# Nguyen-Vidick sieve



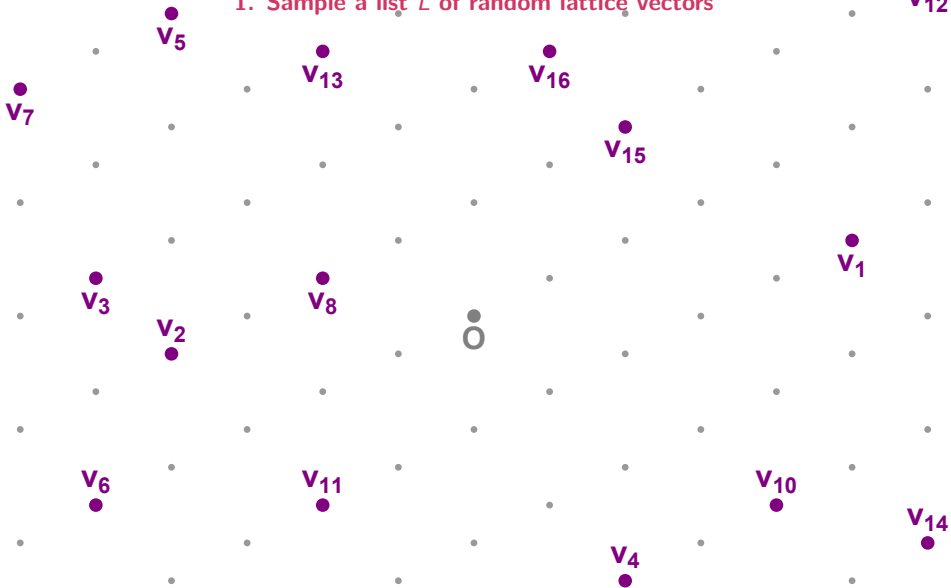
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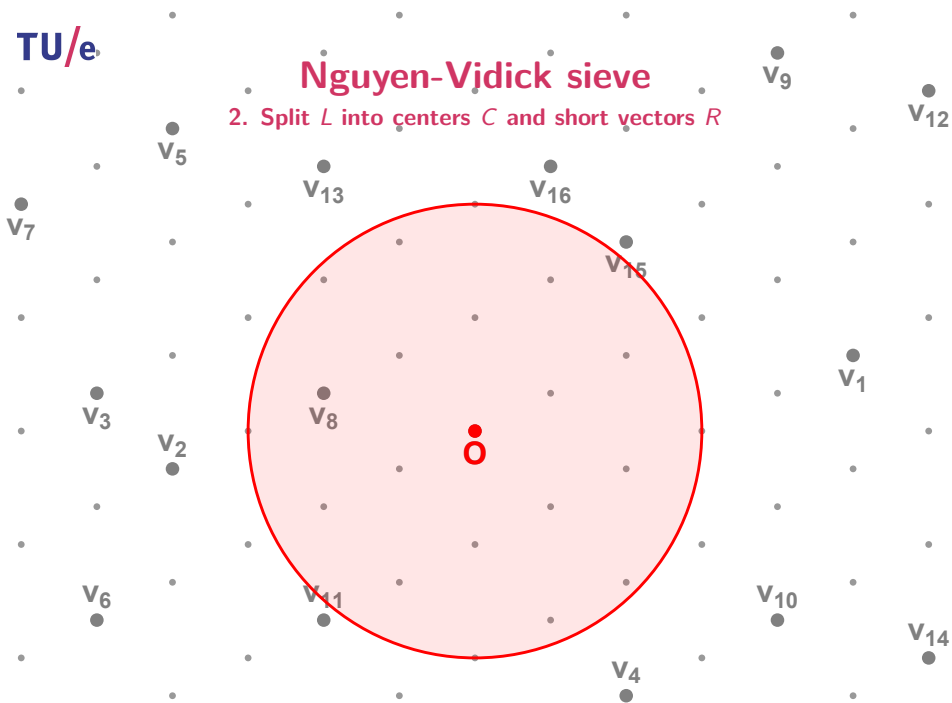
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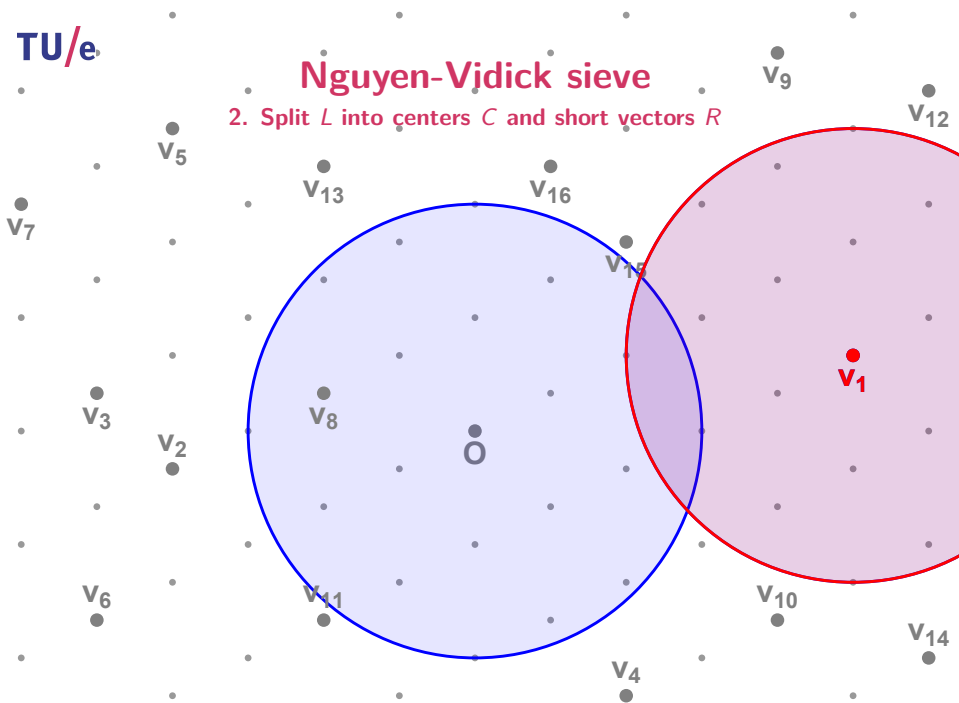
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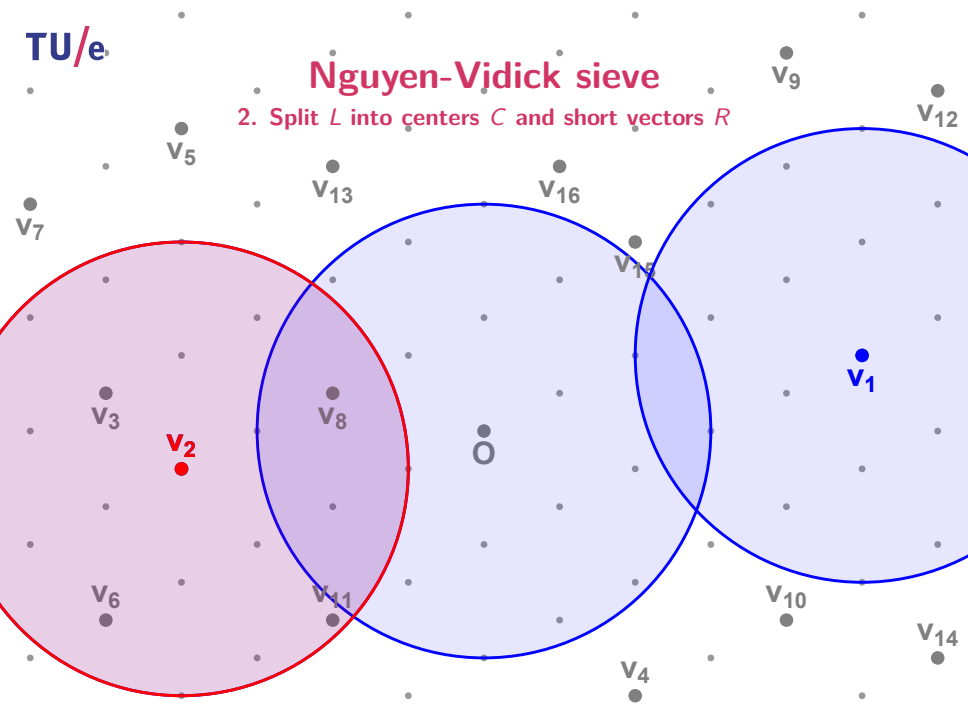
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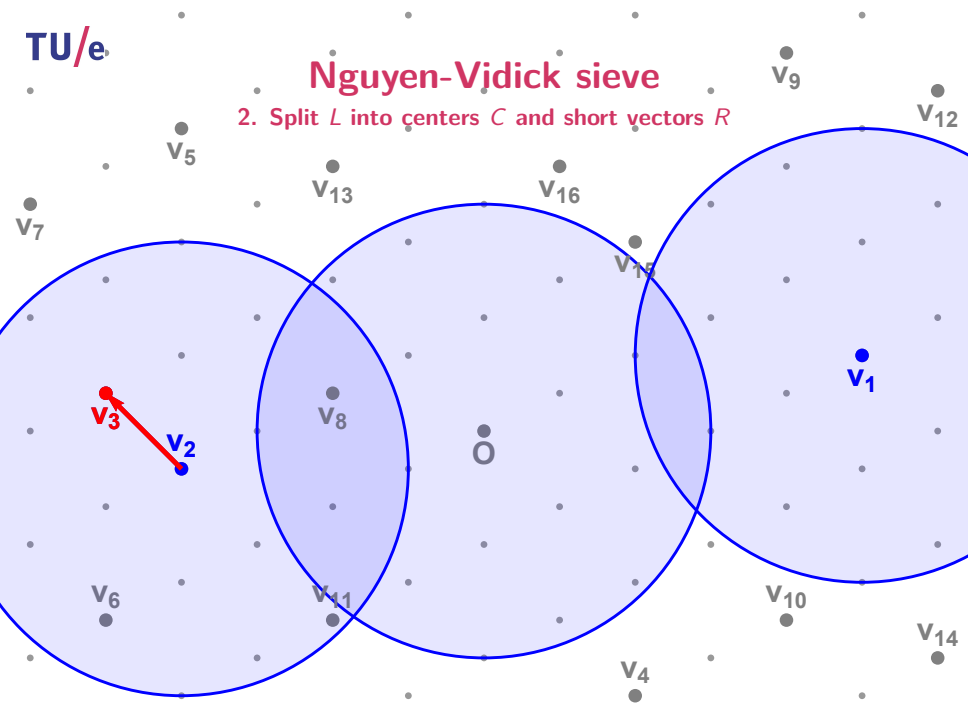
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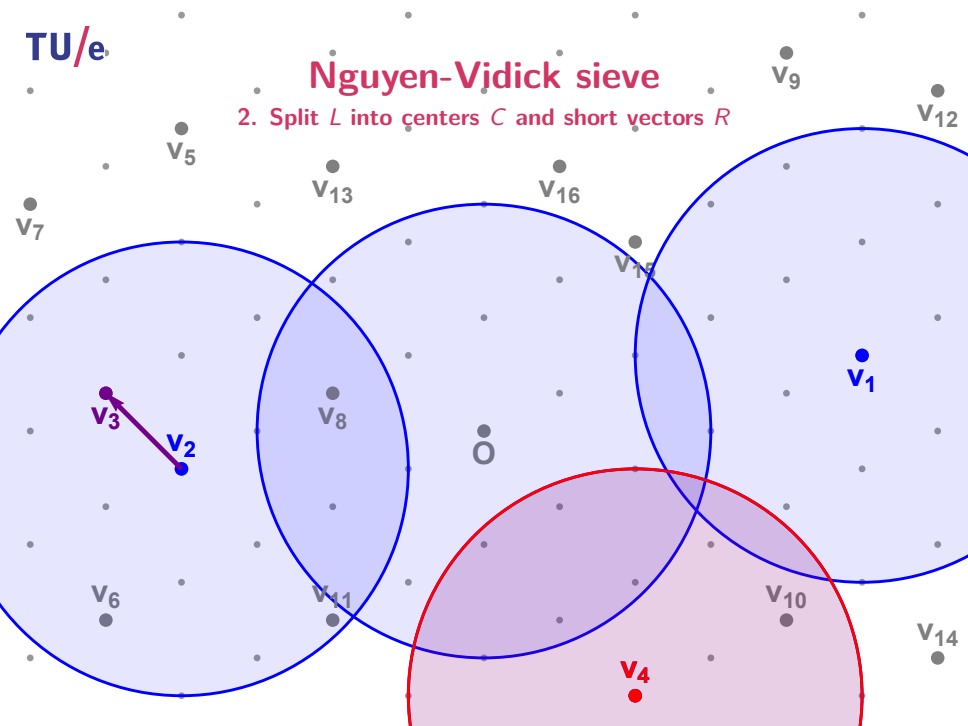
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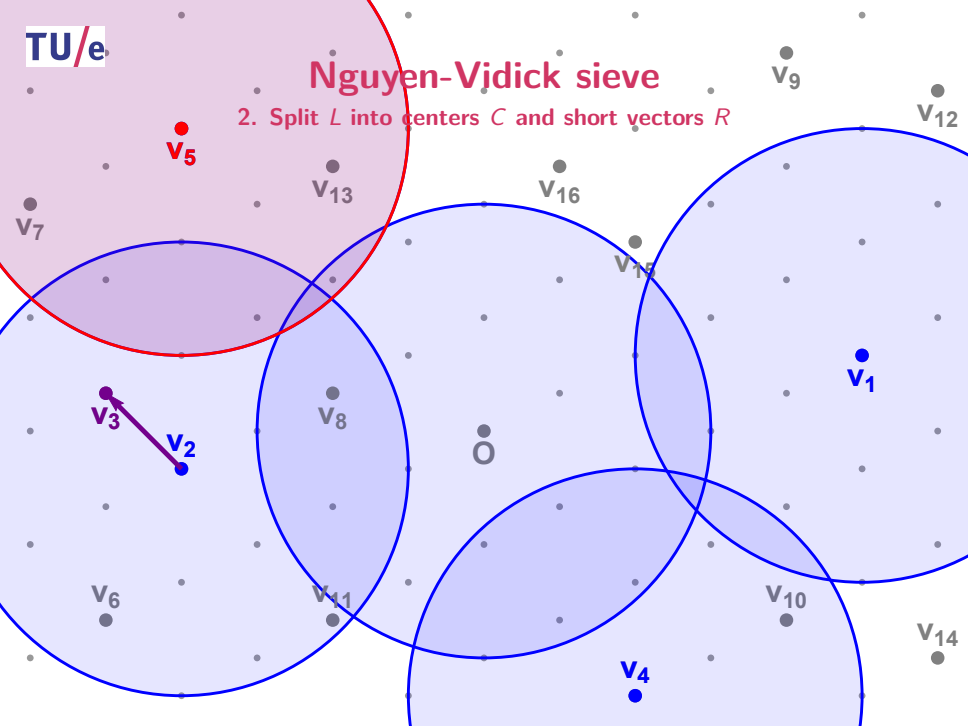


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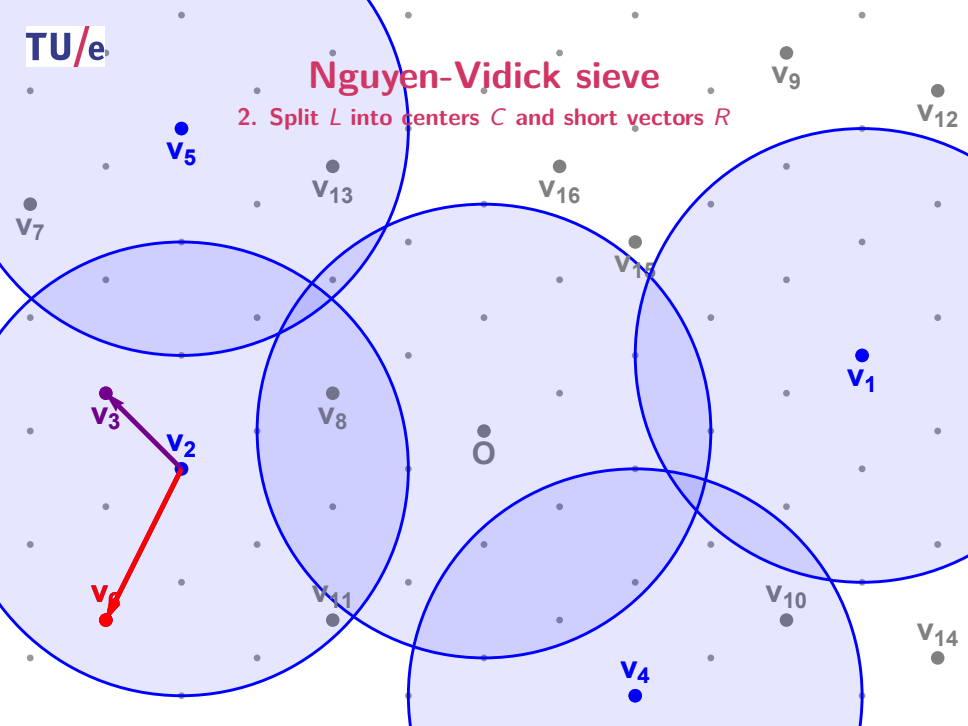


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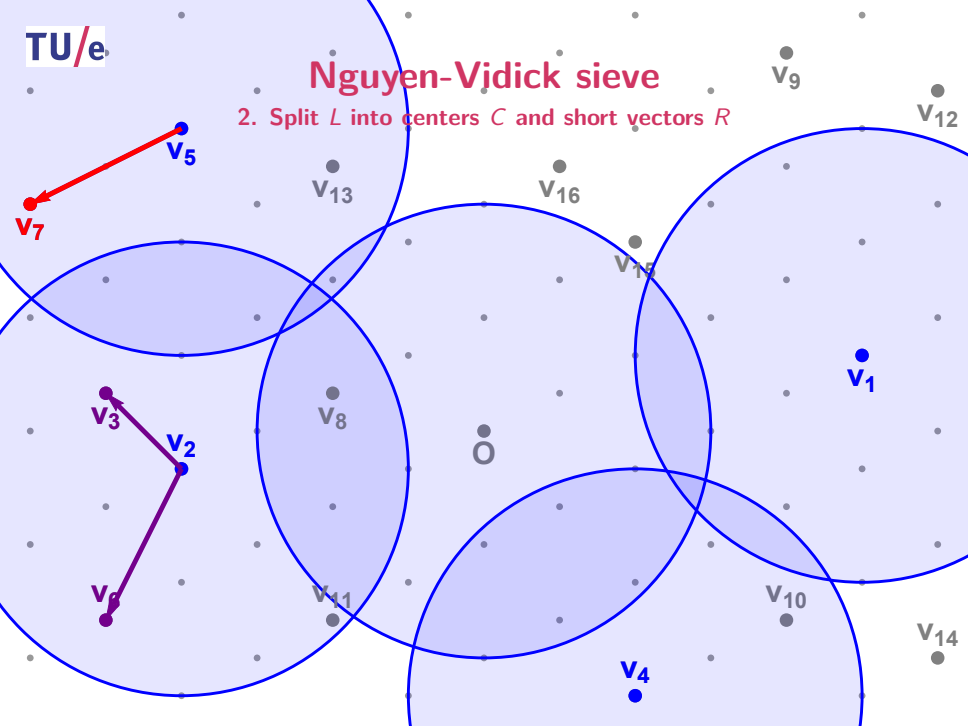
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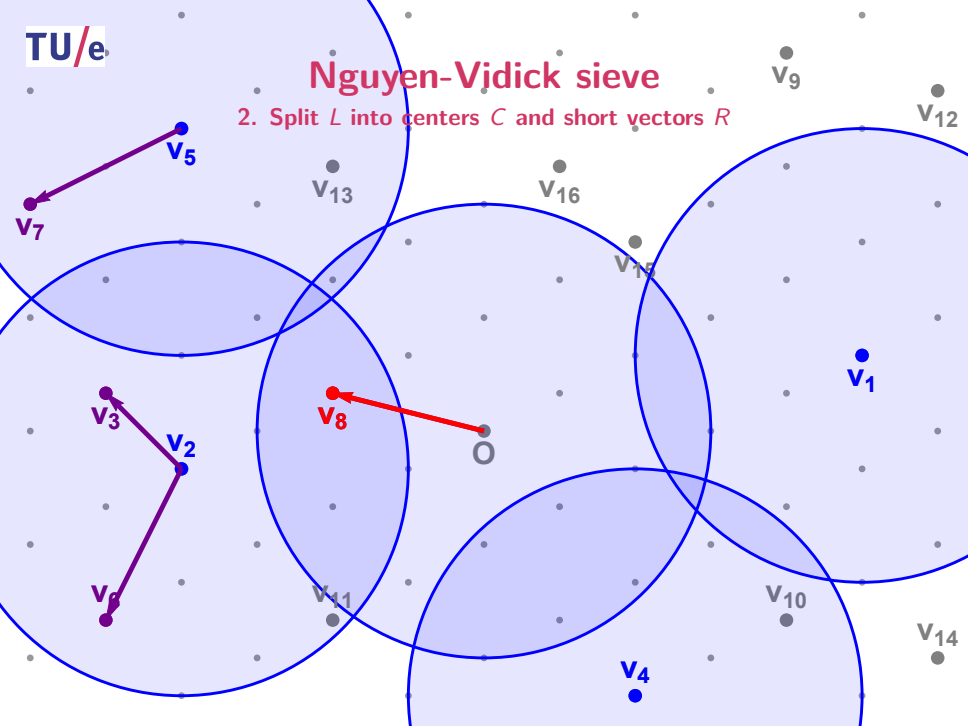
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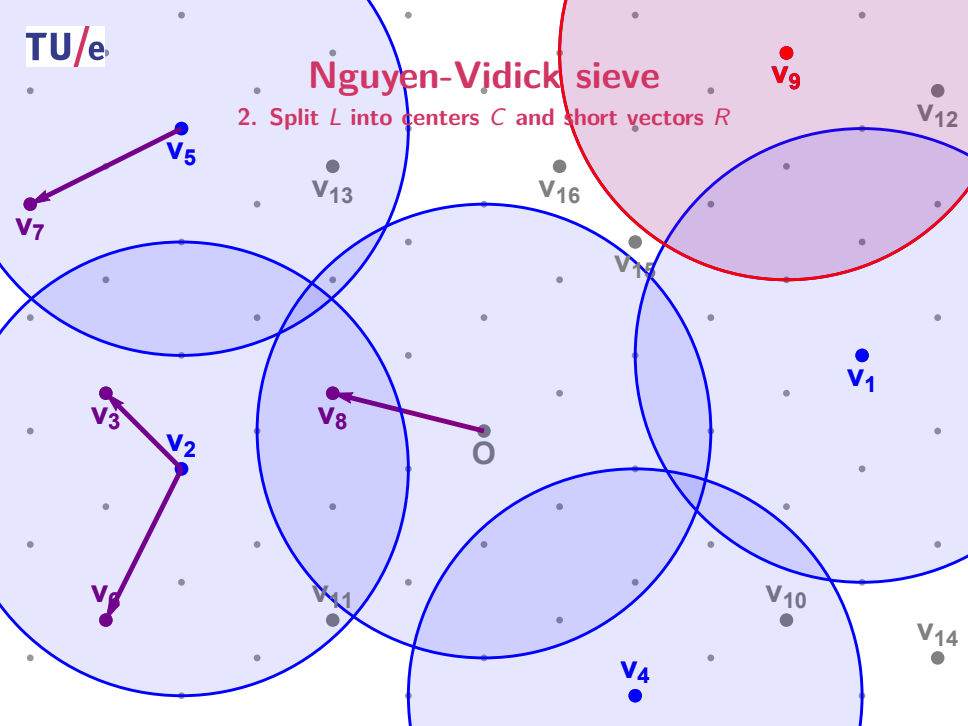




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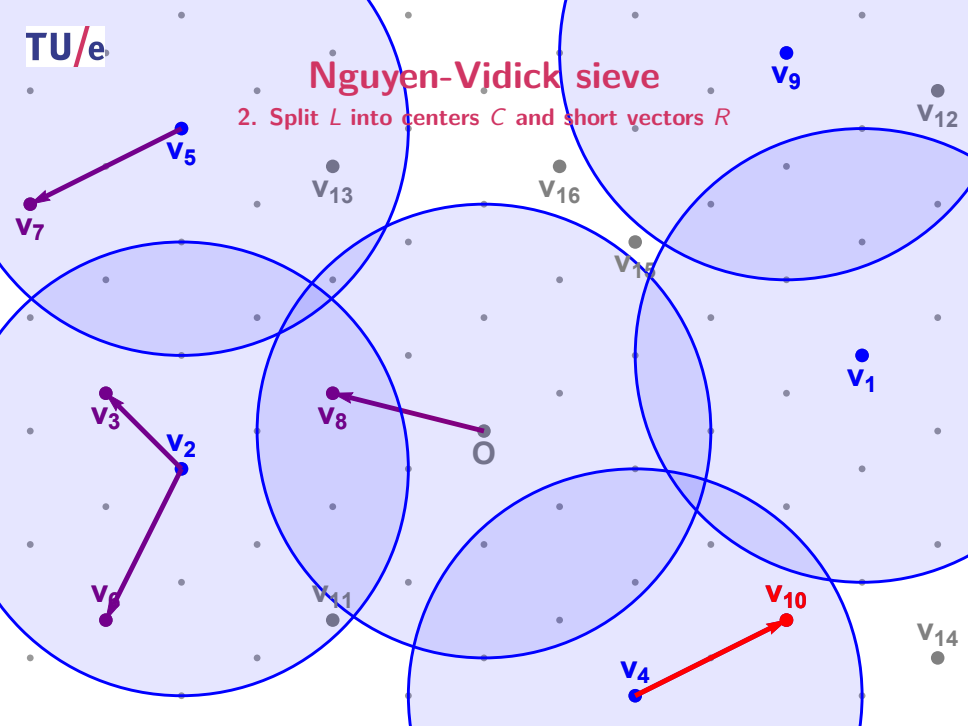
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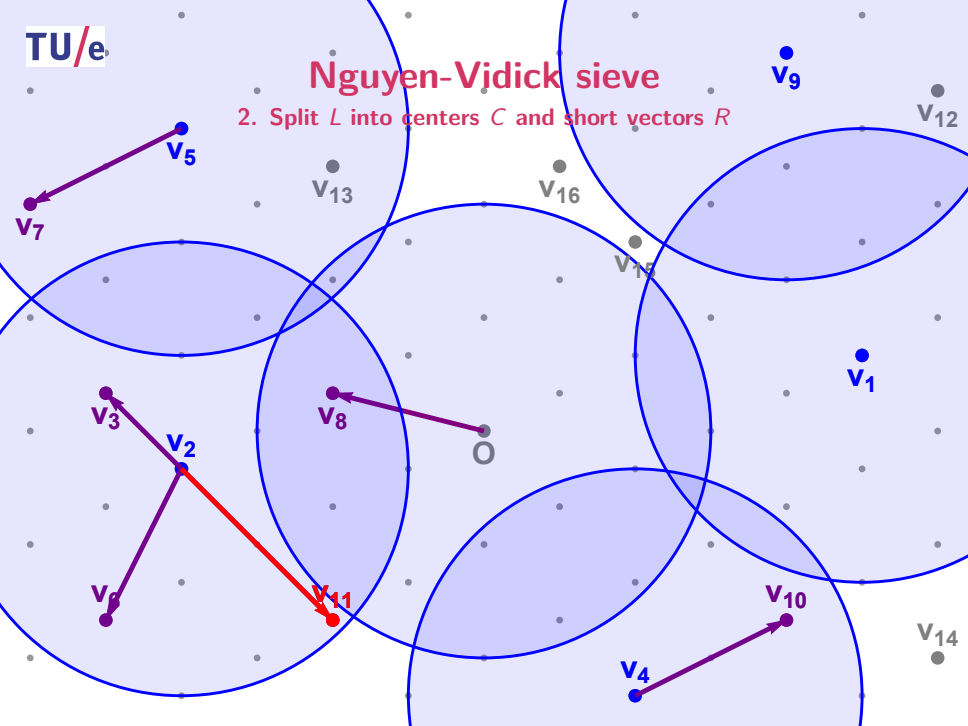
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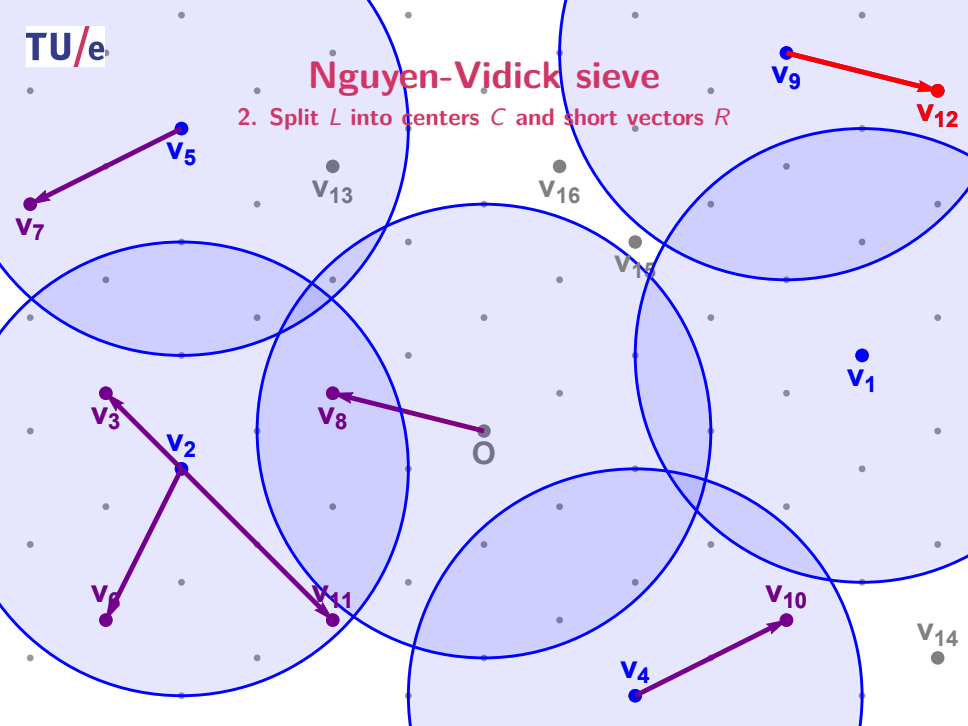
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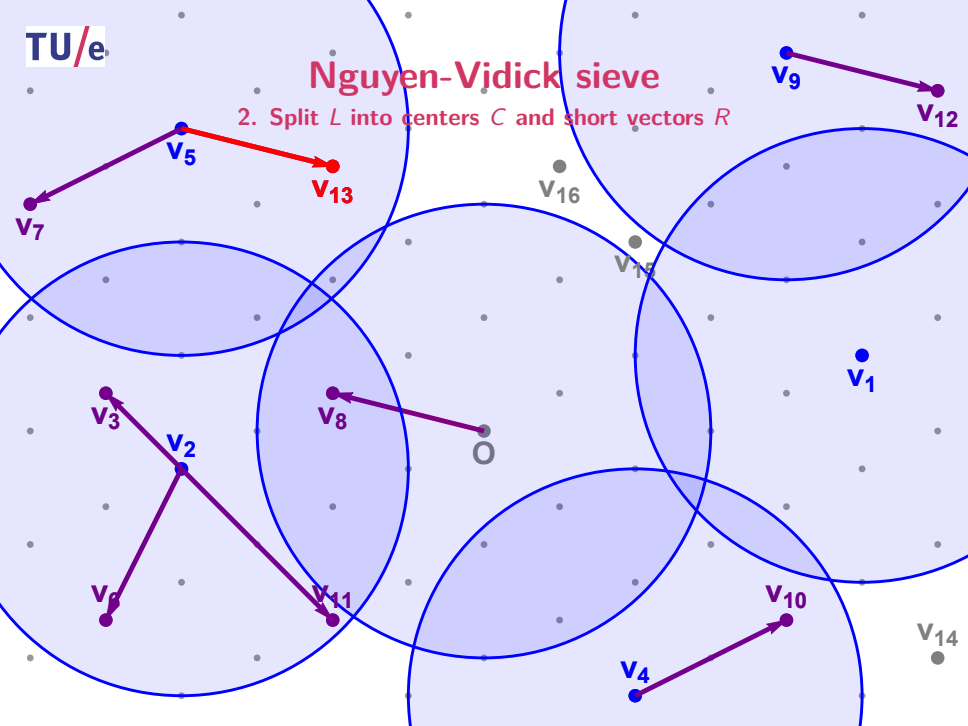


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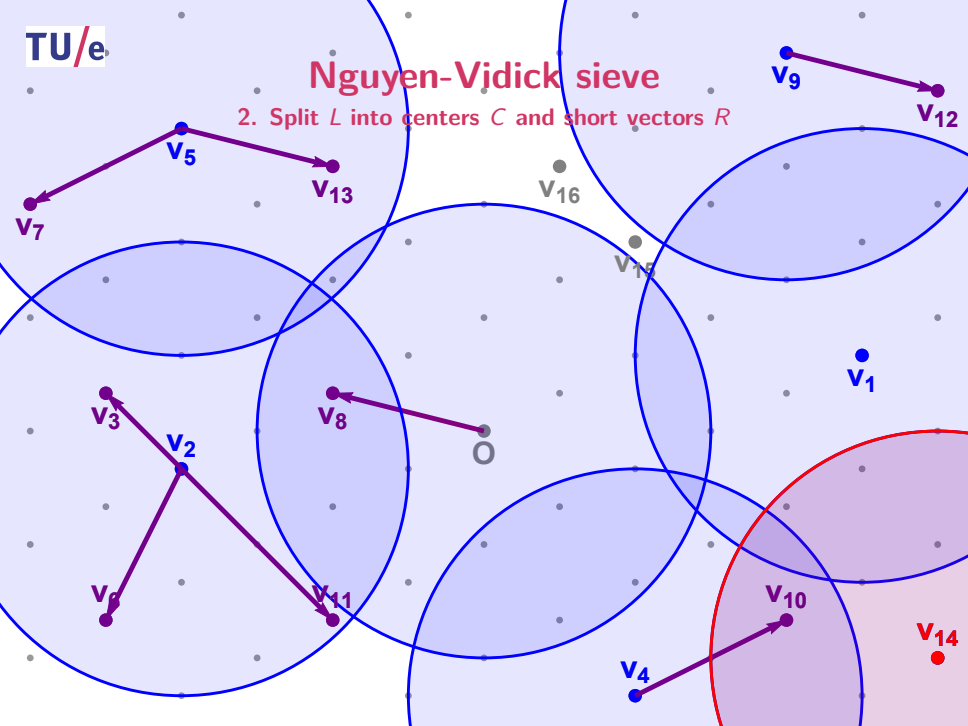


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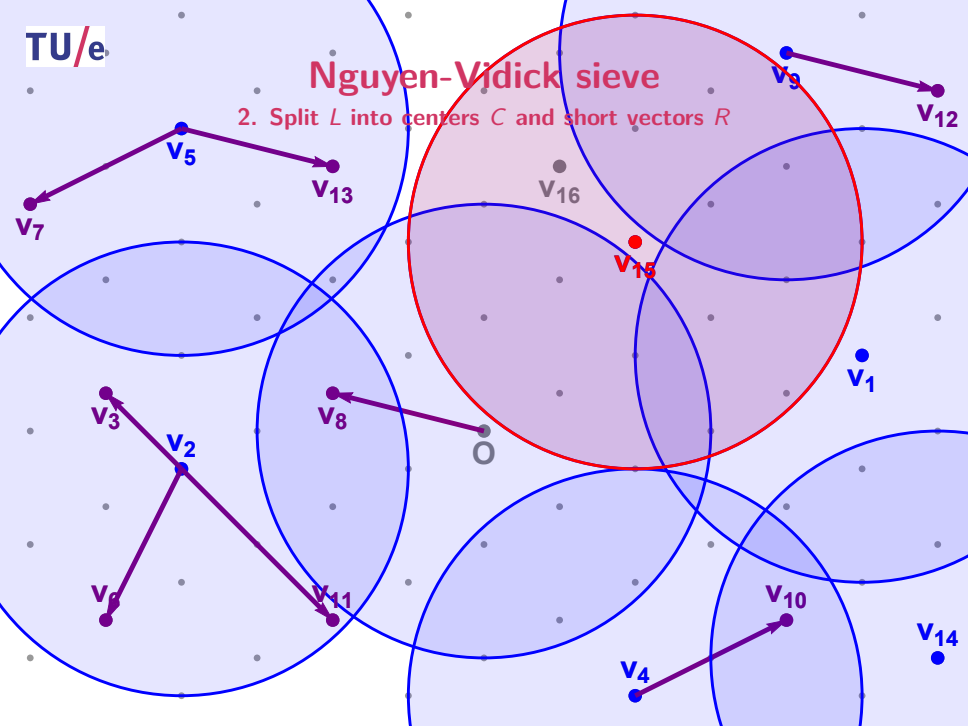
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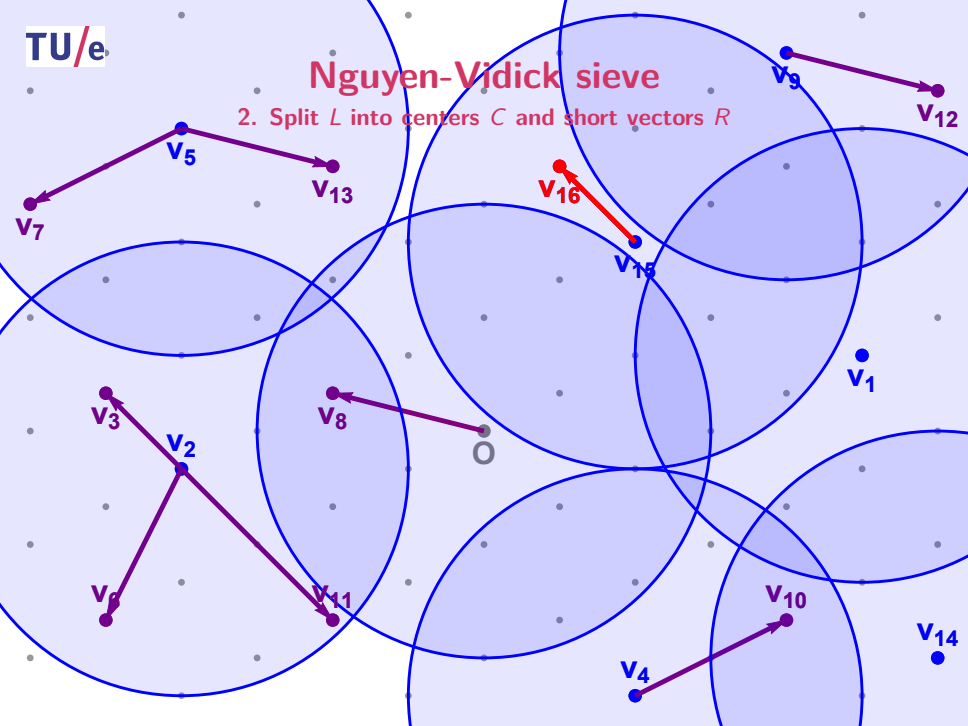
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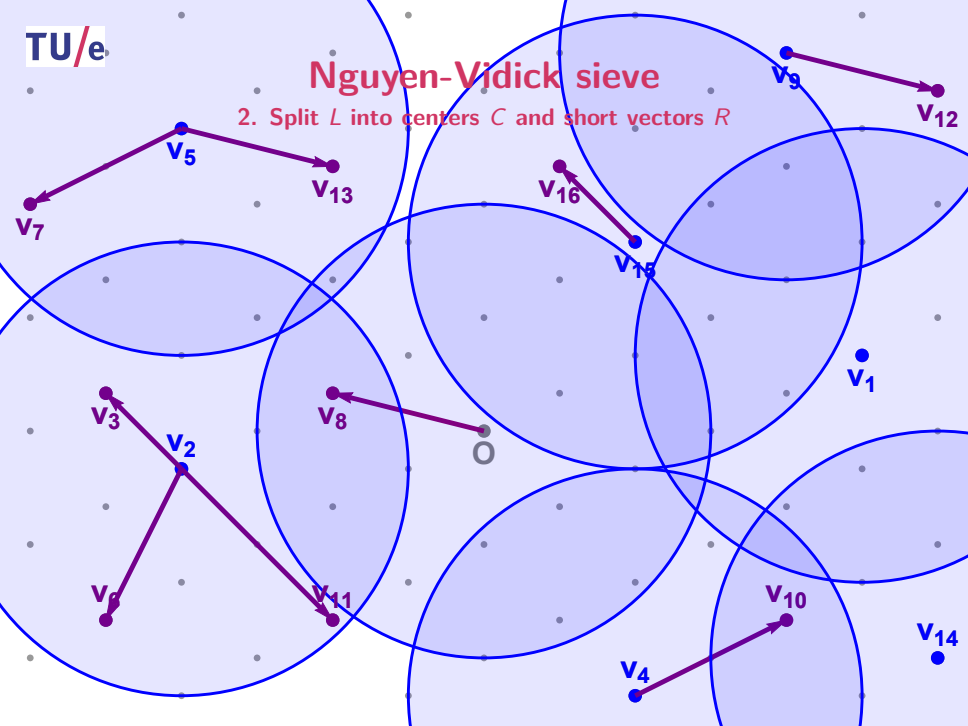
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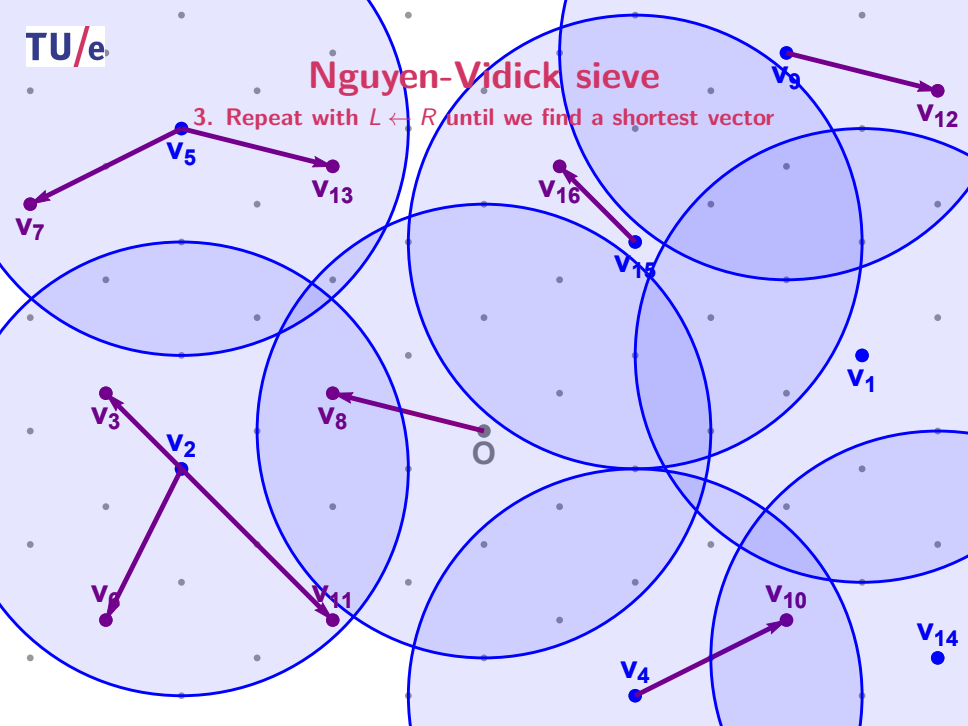
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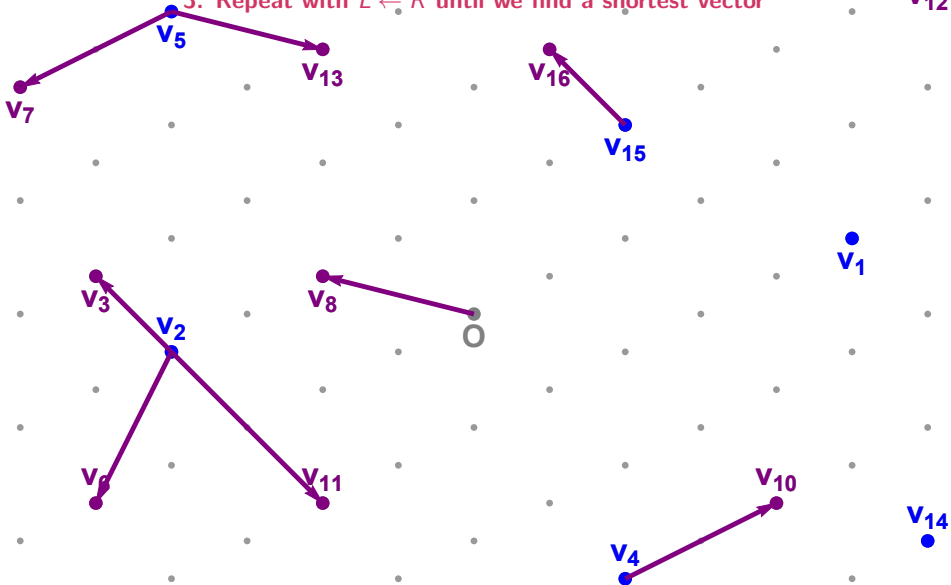
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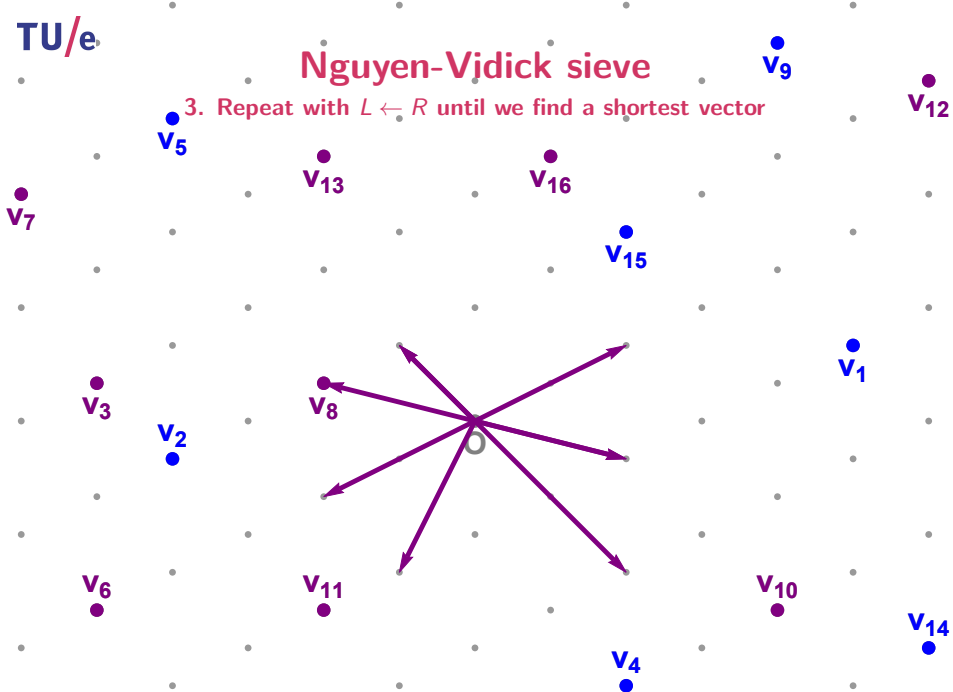
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Overview



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Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The Nguyen-Vidick sieve runs in time  $(4/3)^n$  and space  $\sqrt{4/3}^n$ .

# Nguyen-Vidick sieve

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Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The Nguyen-Vidick sieve runs in time  $2^{0.415n}$  and space  $2^{0.208n}$ .



# Two-level sieve

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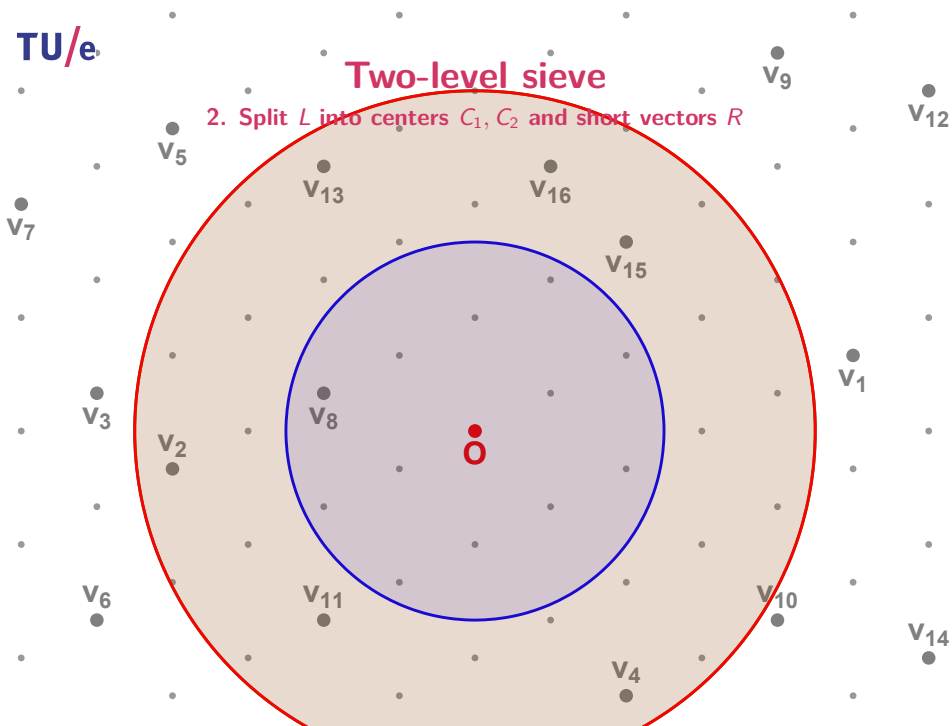
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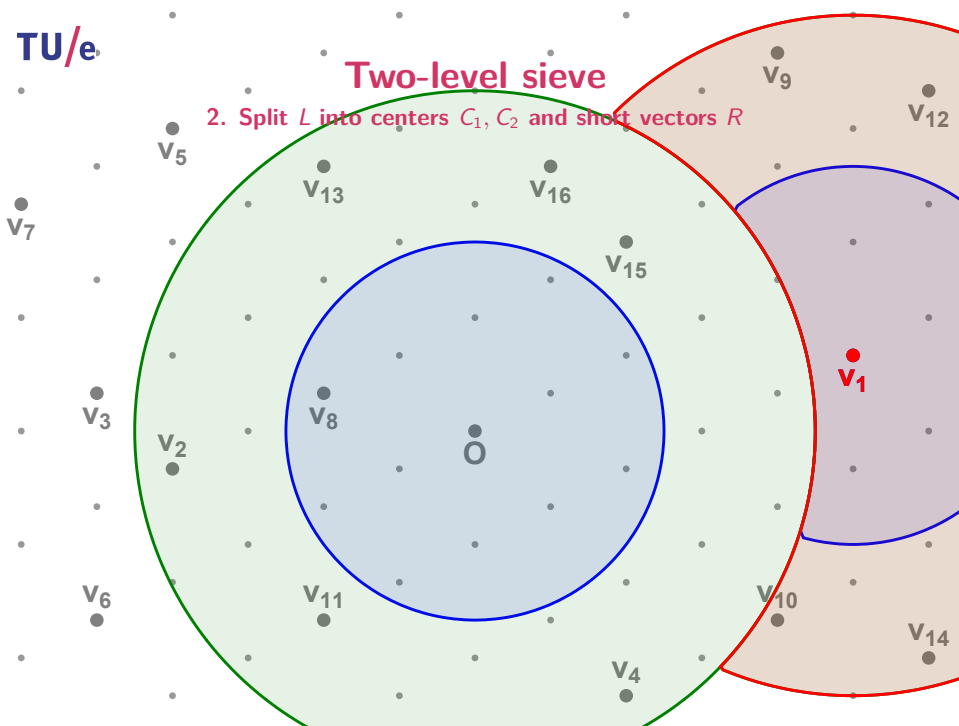
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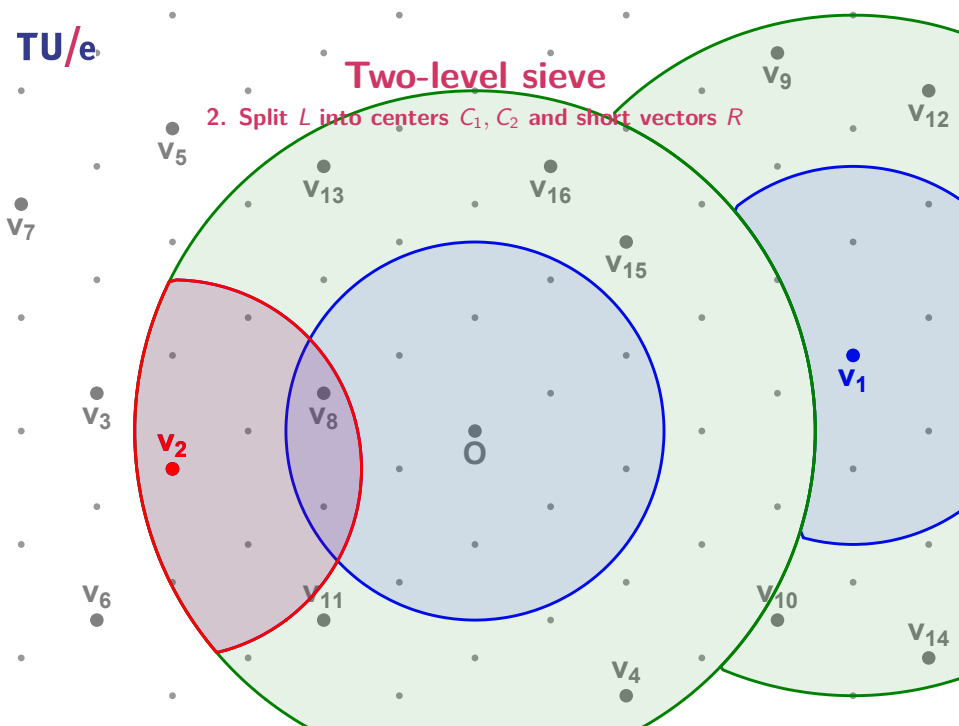
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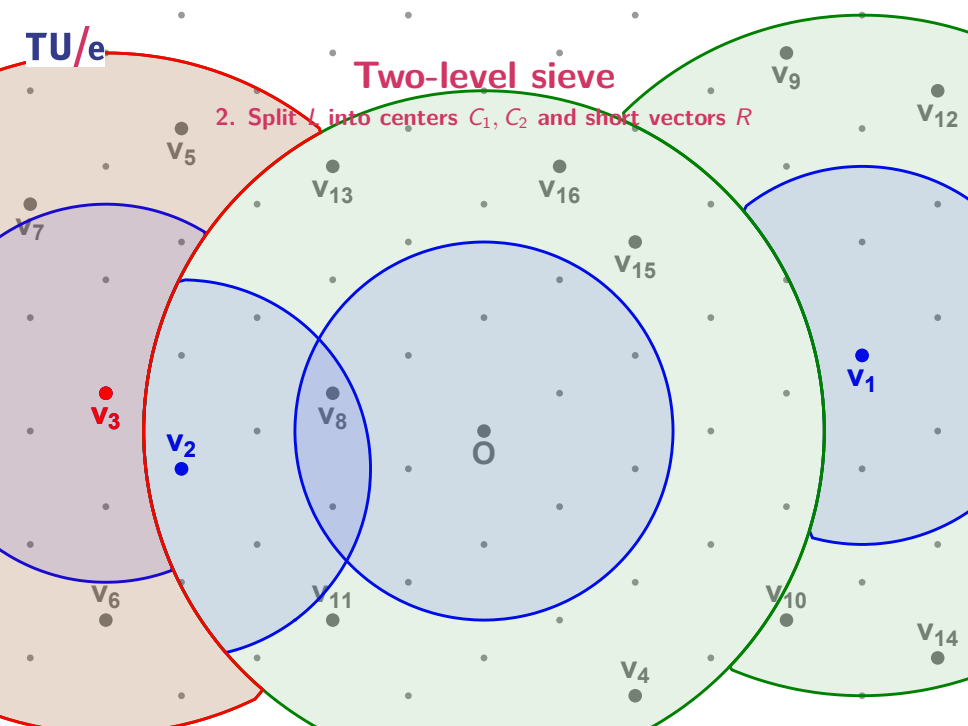
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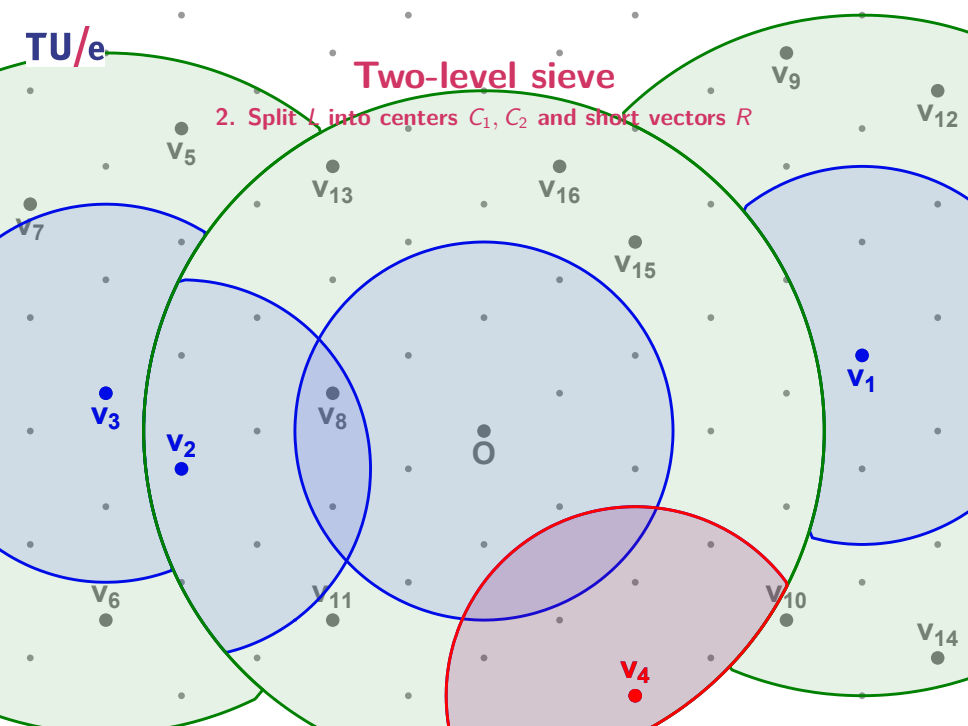
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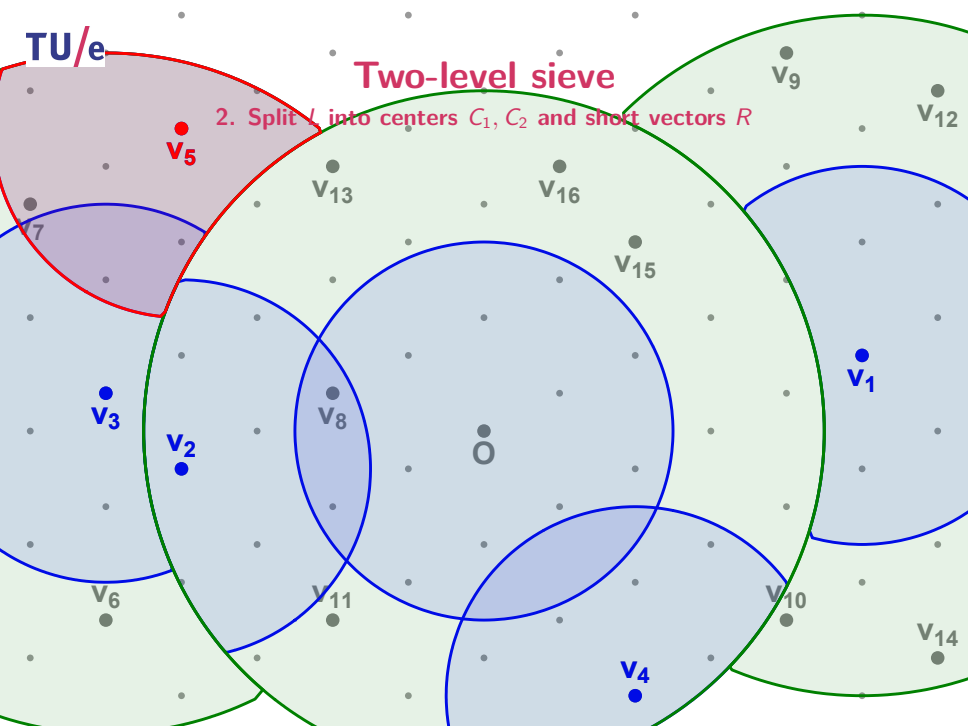




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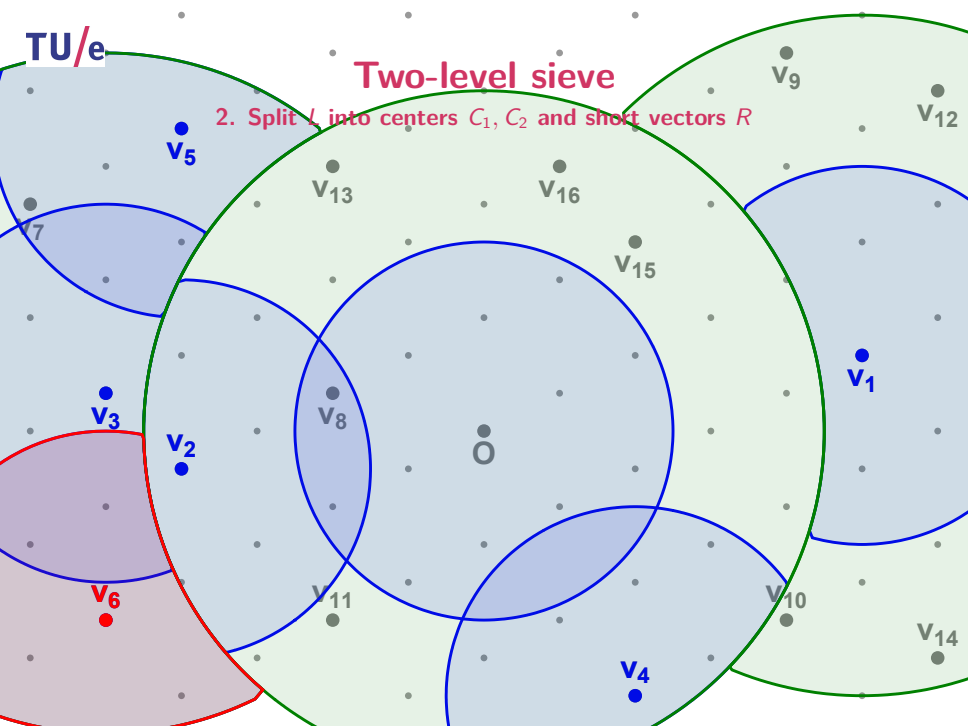
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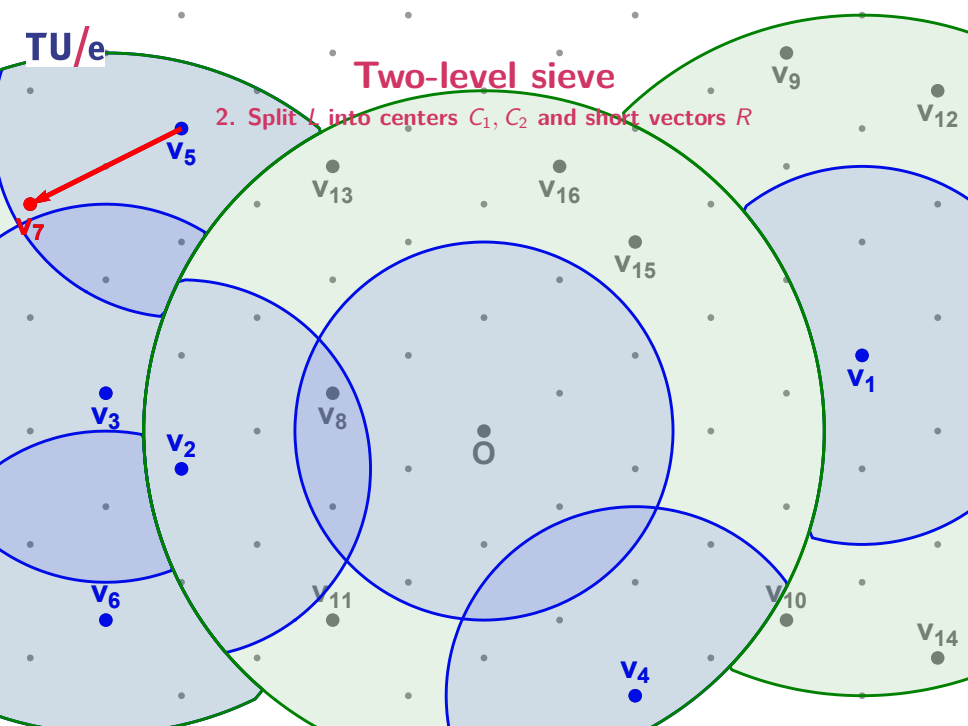
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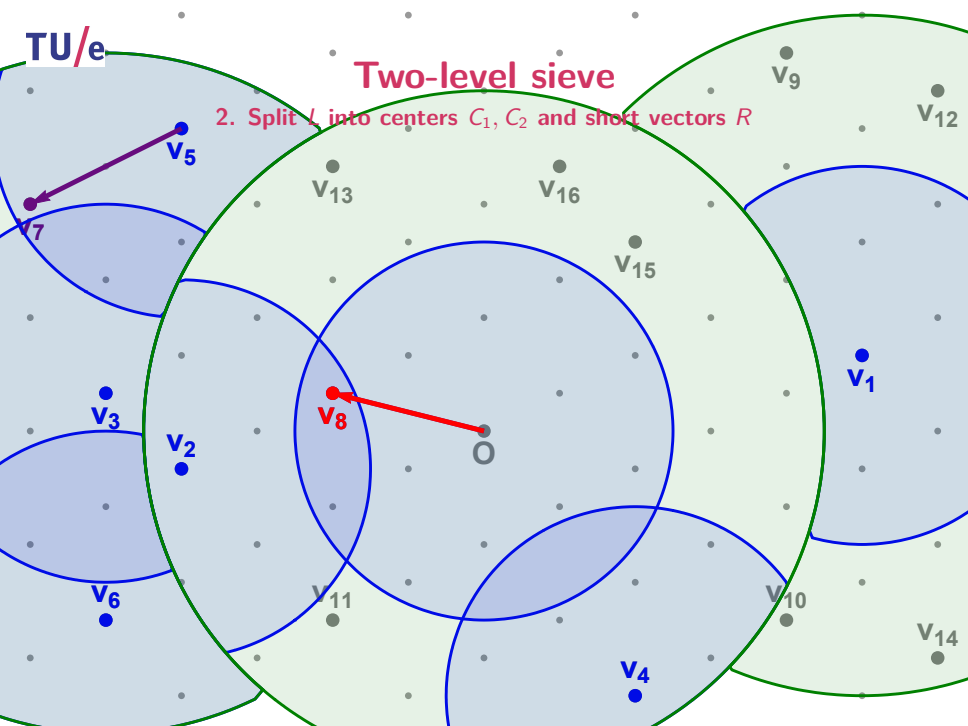
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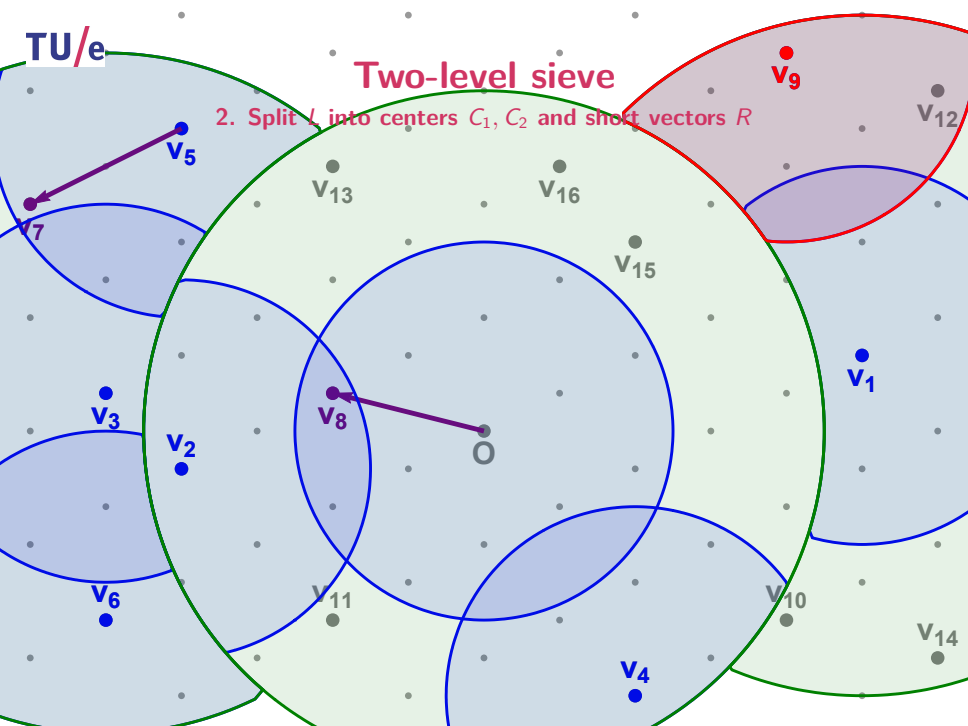
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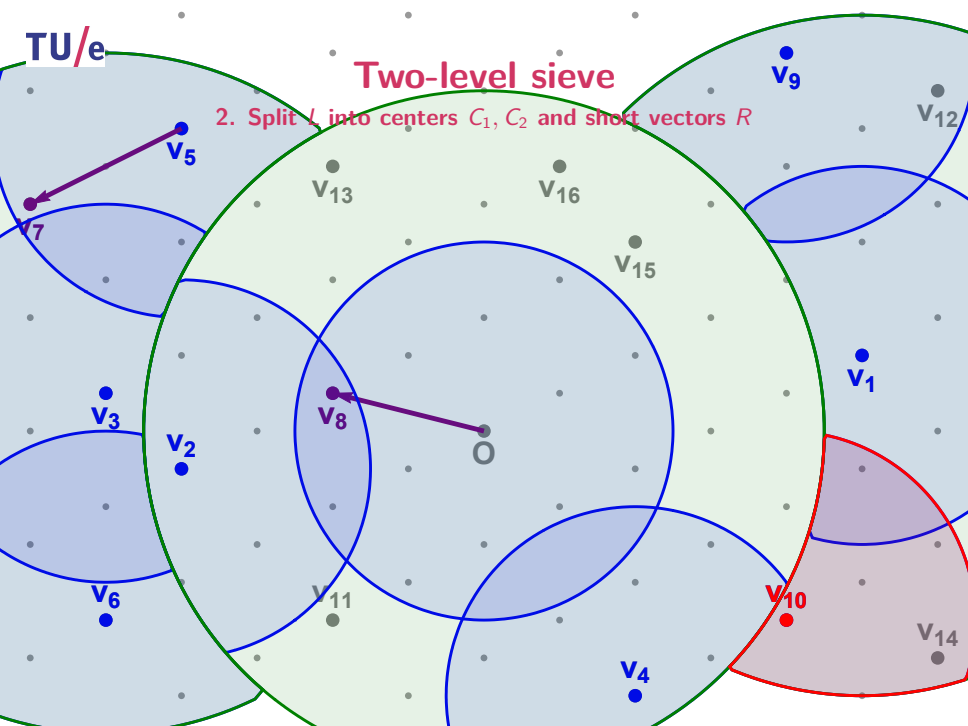
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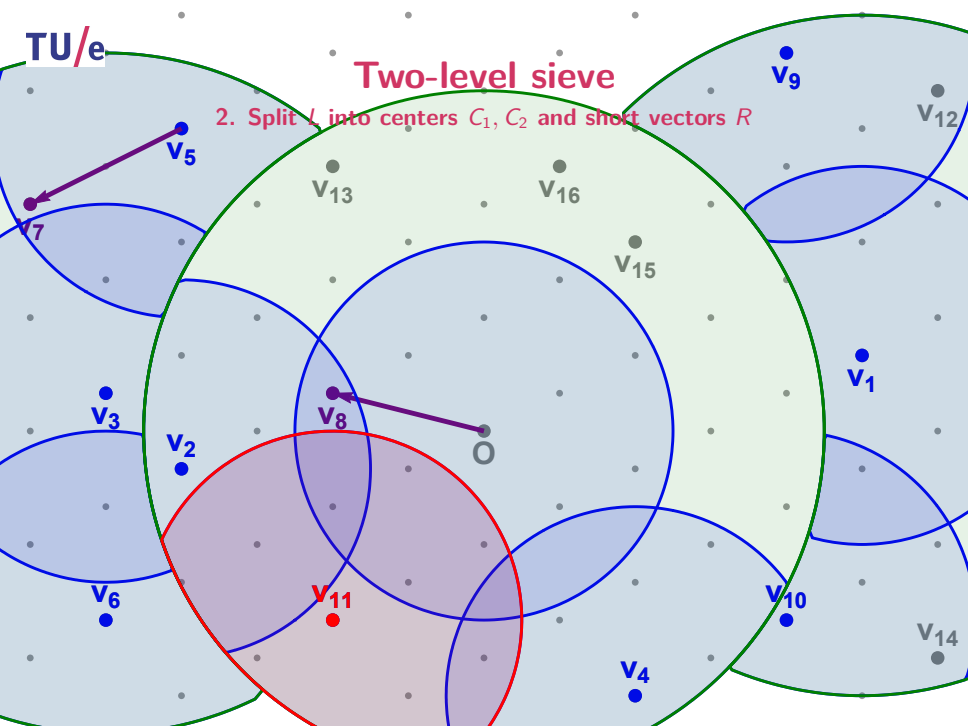
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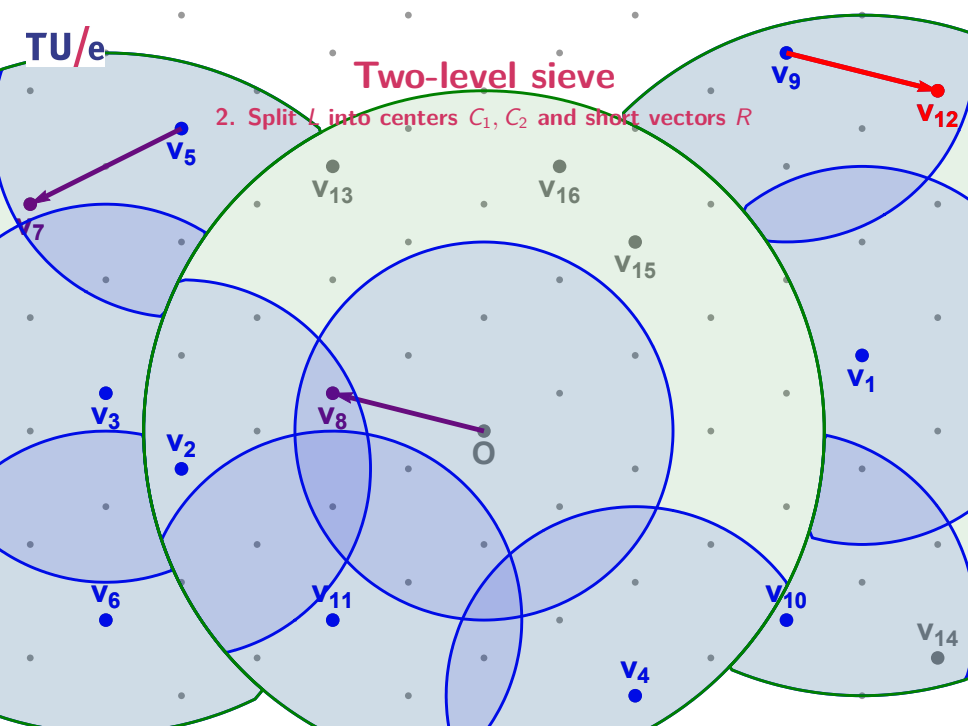
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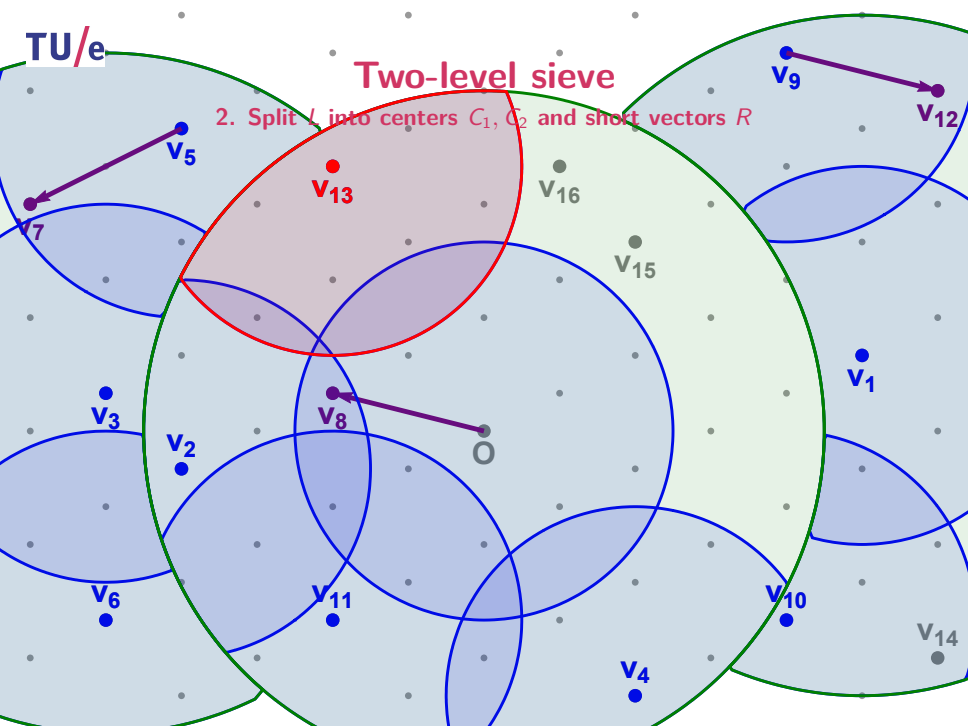




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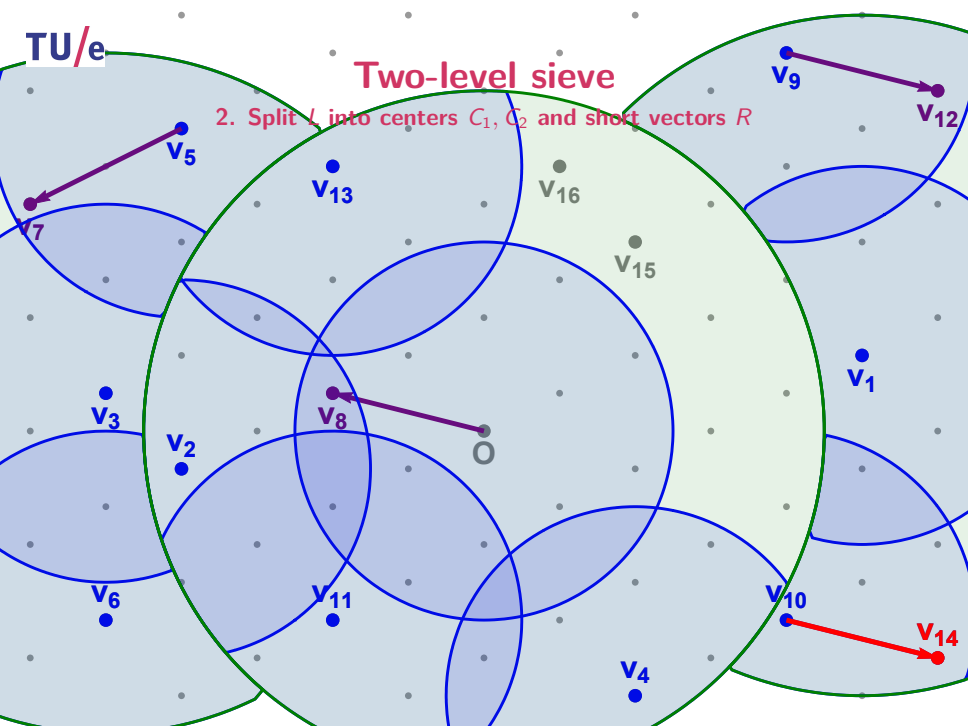
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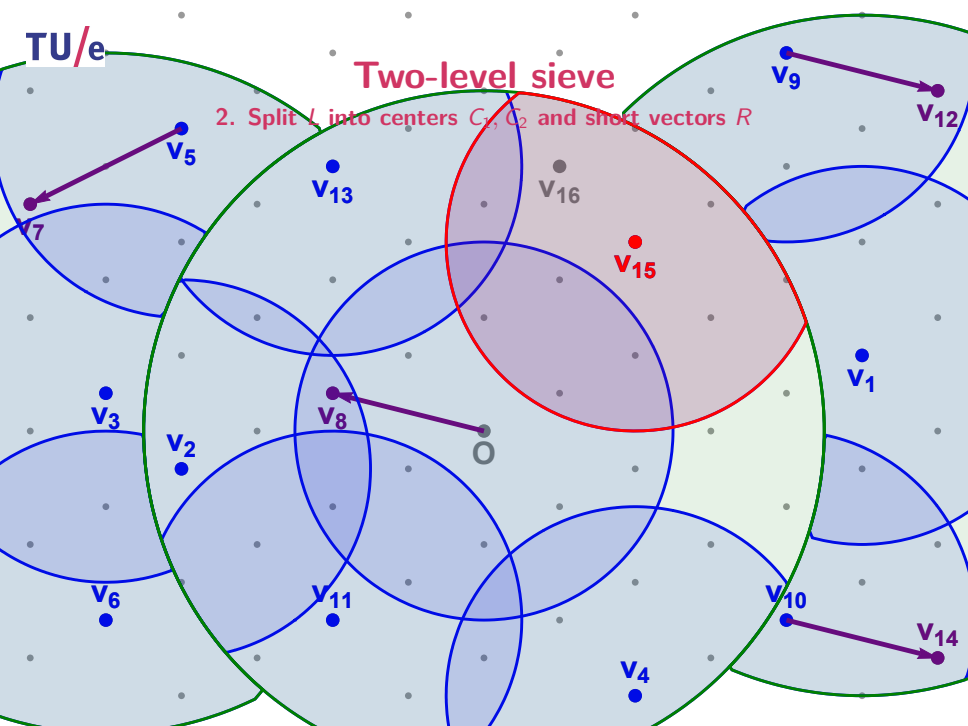
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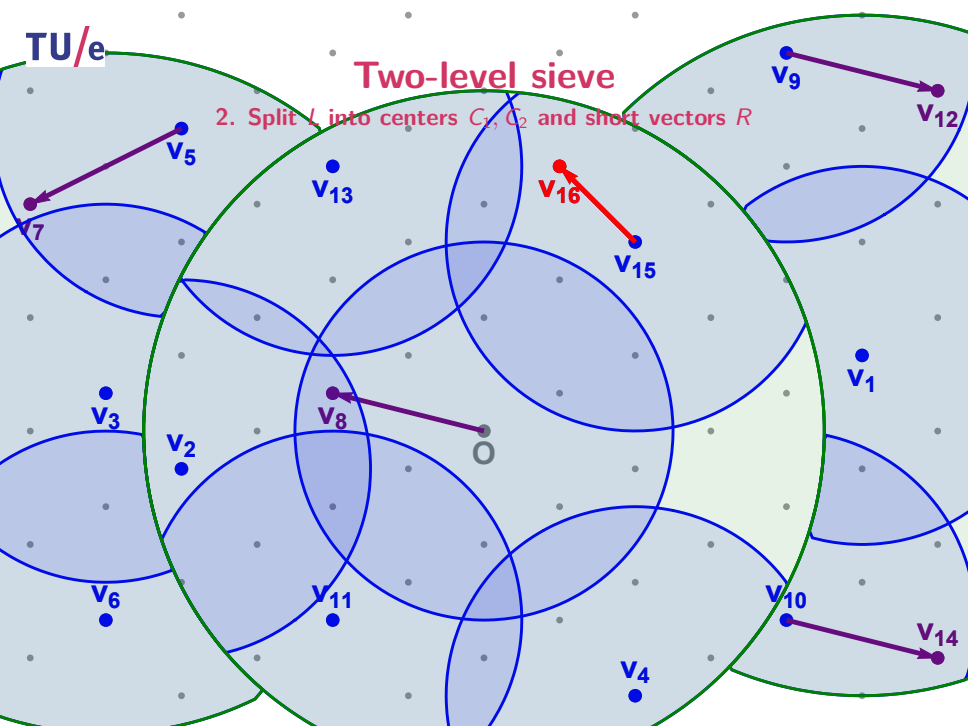
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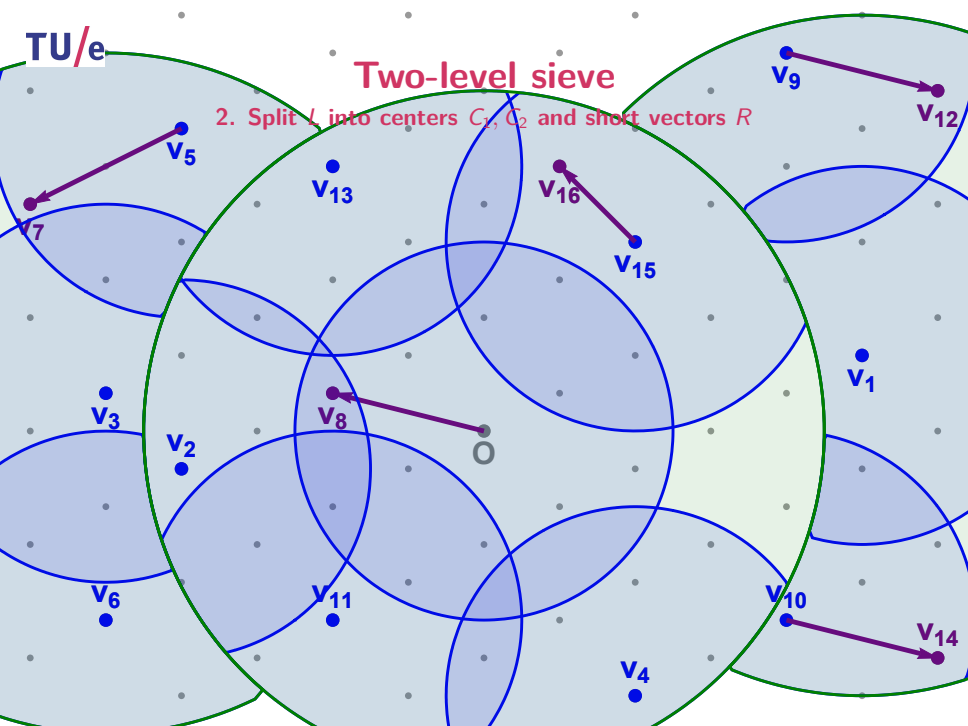
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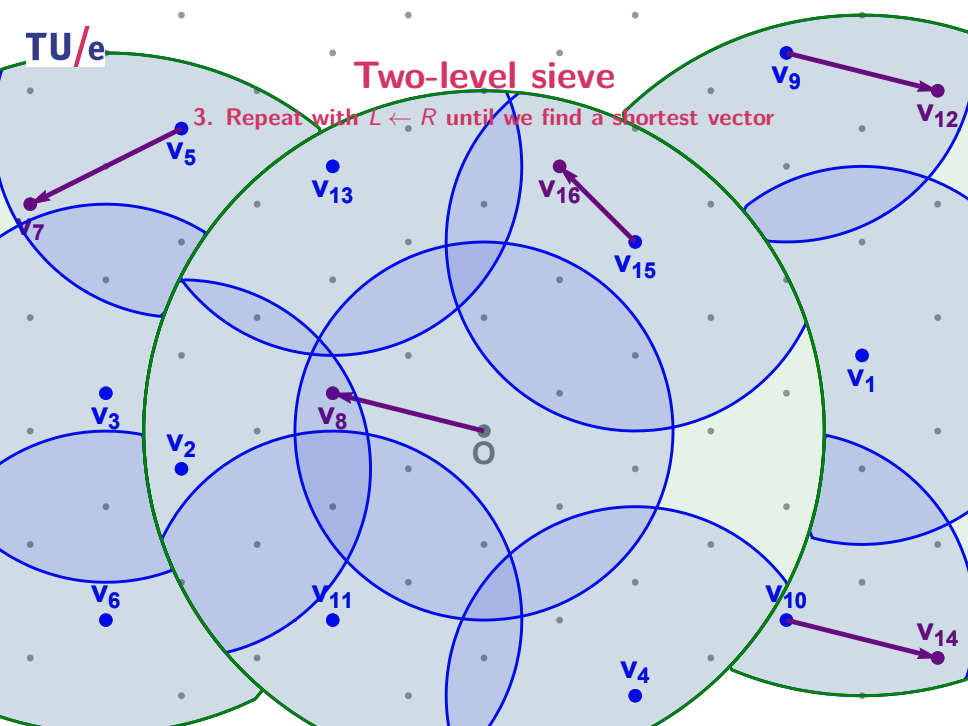
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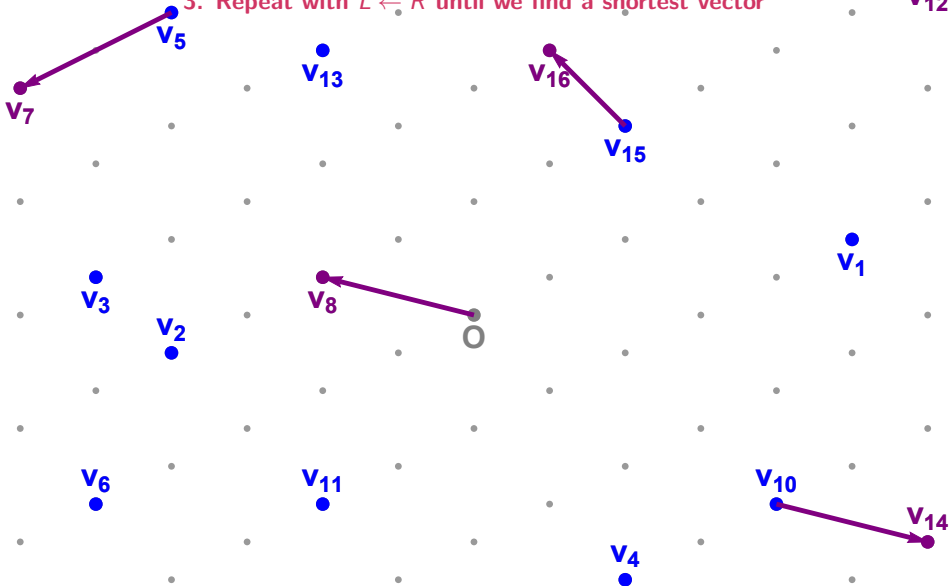
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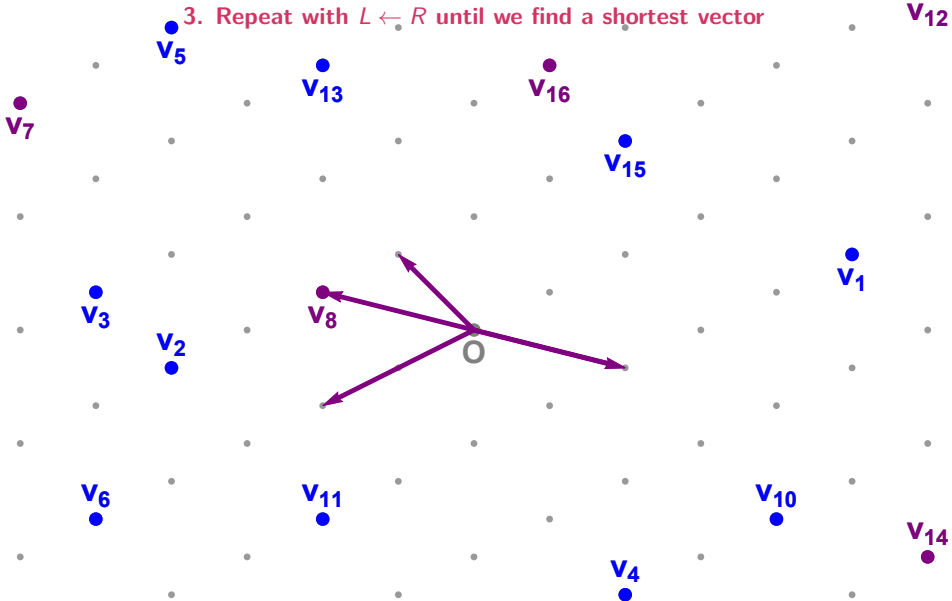
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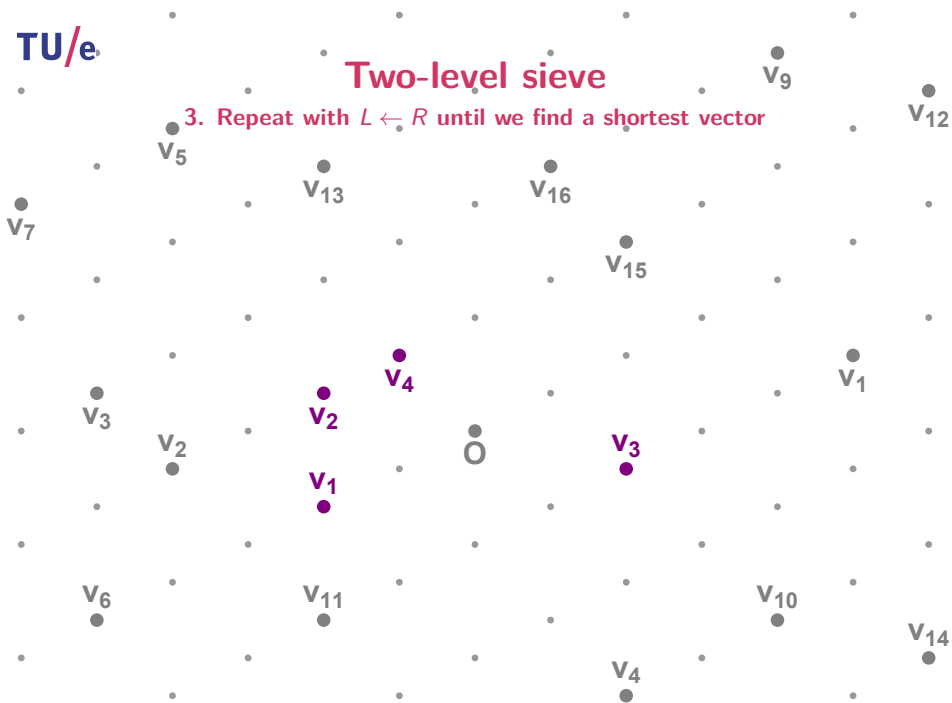
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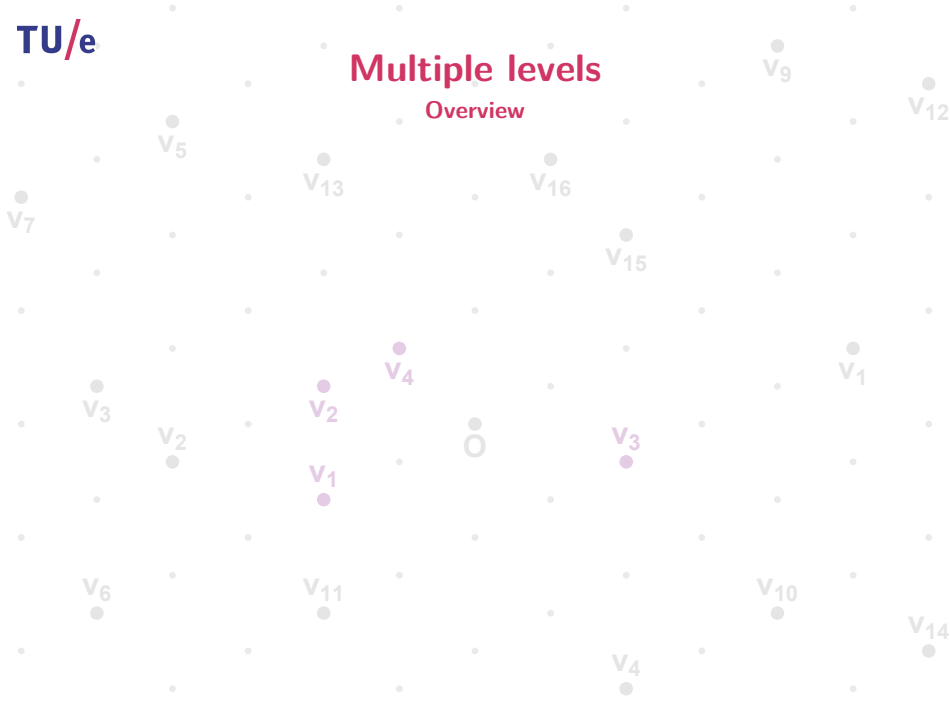
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# Multiple levels

## Overview



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Heuristic (Nguyen and Vidick, J. Math. Crypt. '08)

The one-level sieve runs in time  $2^{0.4150n}$  and space  $2^{0.2075n}$ .



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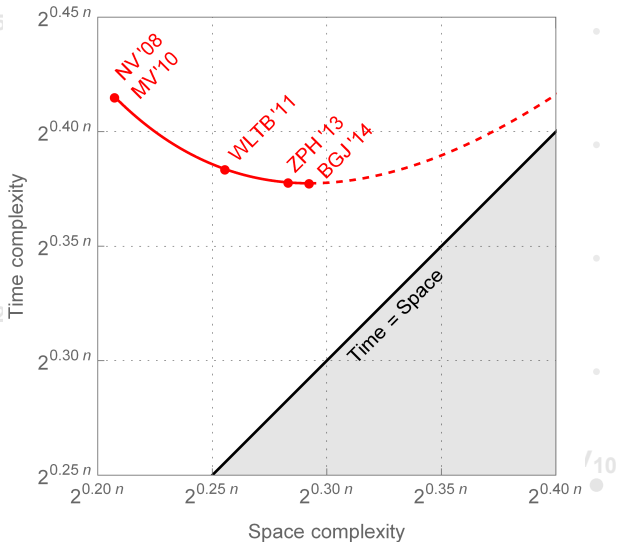
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### Conjecture

The four-level sieve runs in time  $2^{0.3774n}$  and space  $2^{0.2925n}$ , and higher-level sieves are not faster than this.

## Sieving

Space/time trade-off



# Quantum Search

## Classical form

**Problem:** Given a list  $L$  of size  $N$ , and a function  $f : L \rightarrow \{0, 1\}$  such that there is exactly one element  $e \in L$  with  $f(e) = 1$ . Find this element  $e$ .

- Classical search:  $\Theta(N)$  time
- Quantum search:  $\Theta(\sqrt{N})$  time [Gro96]



# Quantum Search

## General form

**Problem:** Given a list  $L$  of size  $N$ , and a function  $f : L \rightarrow \{0, 1\}$  such that there are  $c = O(1)$  elements  $e \in L$  with  $f(e) = 1$ . Find one such element  $e$ .

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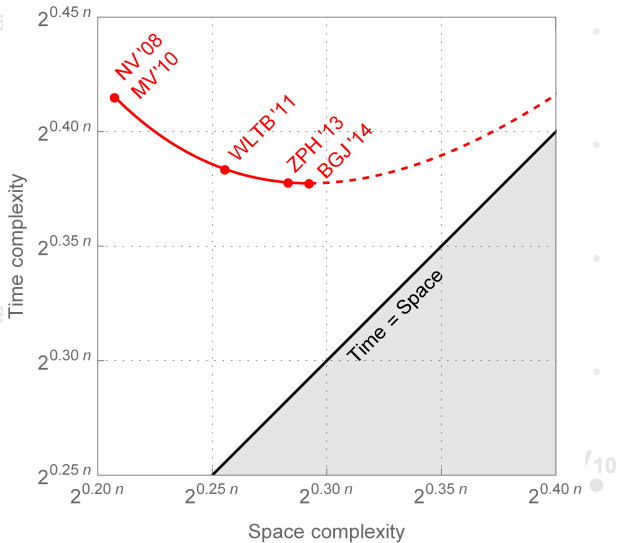
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Potentially speed up search subroutines in sieving

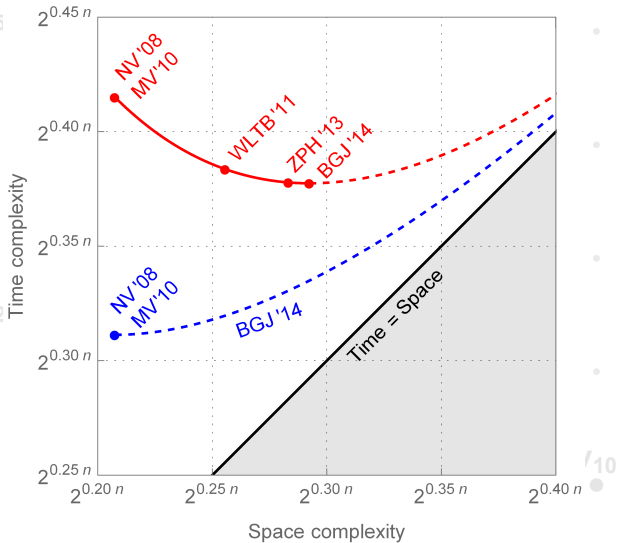
## Sieving

Space/time trade-off



## Sieving

Space/time trade-off



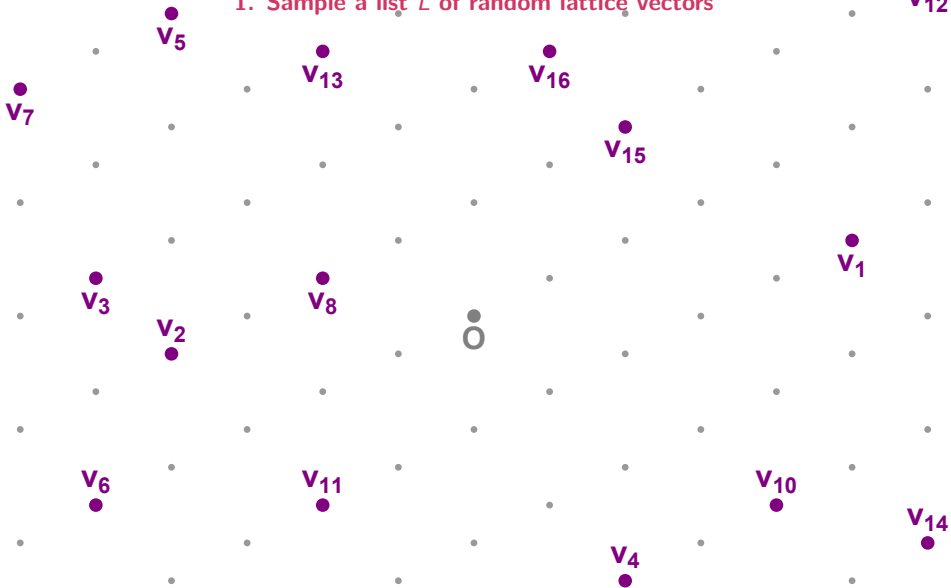
# Nguyen-Vidick sieve with LSH

1. Sample a list  $L$  of random lattice vectors



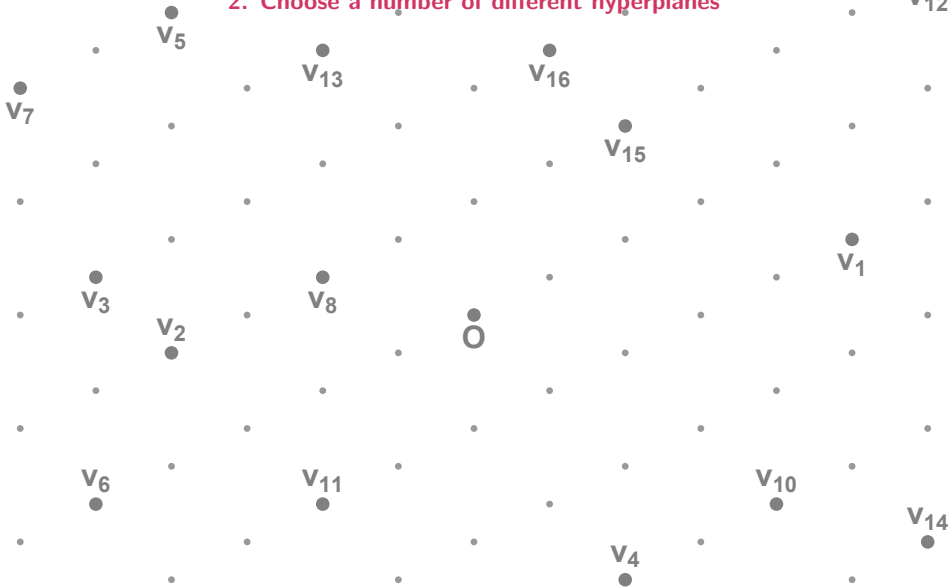
# Nguyen-Vidick sieve with LSH

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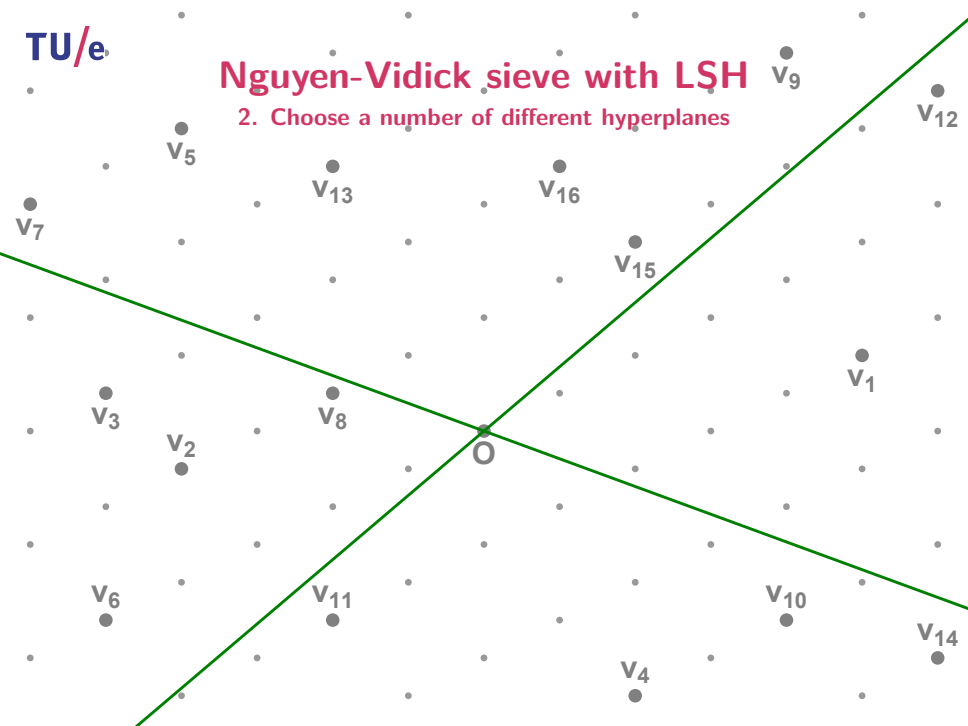
# Nguyen-Vidick sieve with LSH

2. Choose a number of different hyperplanes



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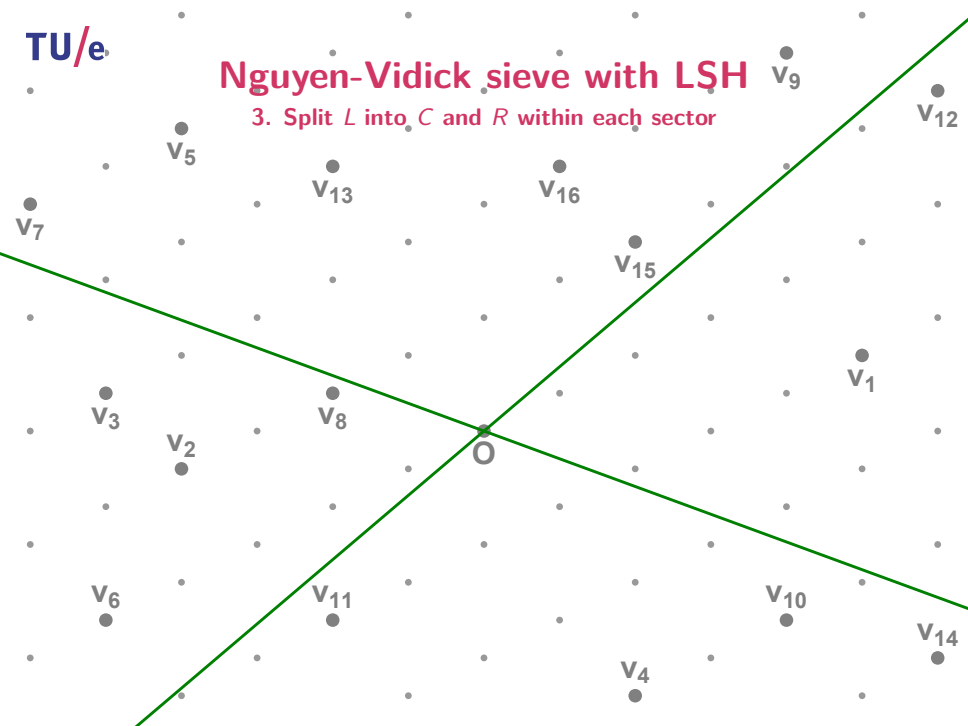
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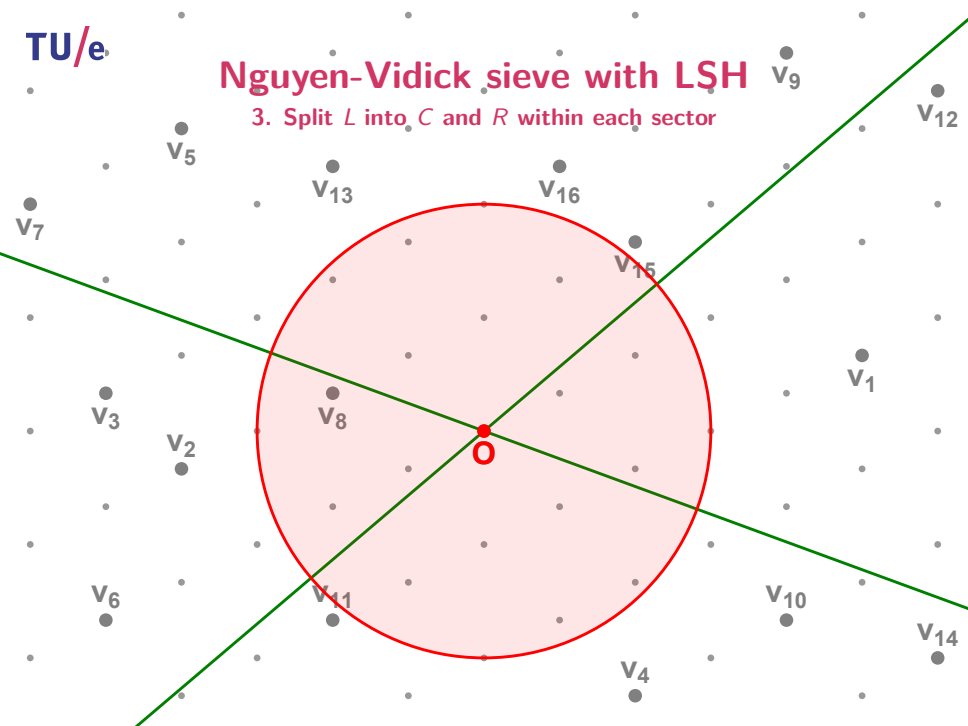
# Nguyen-Vidick sieve with LSH

3. Split  $L$  into  $C$  and  $R$  within each sector

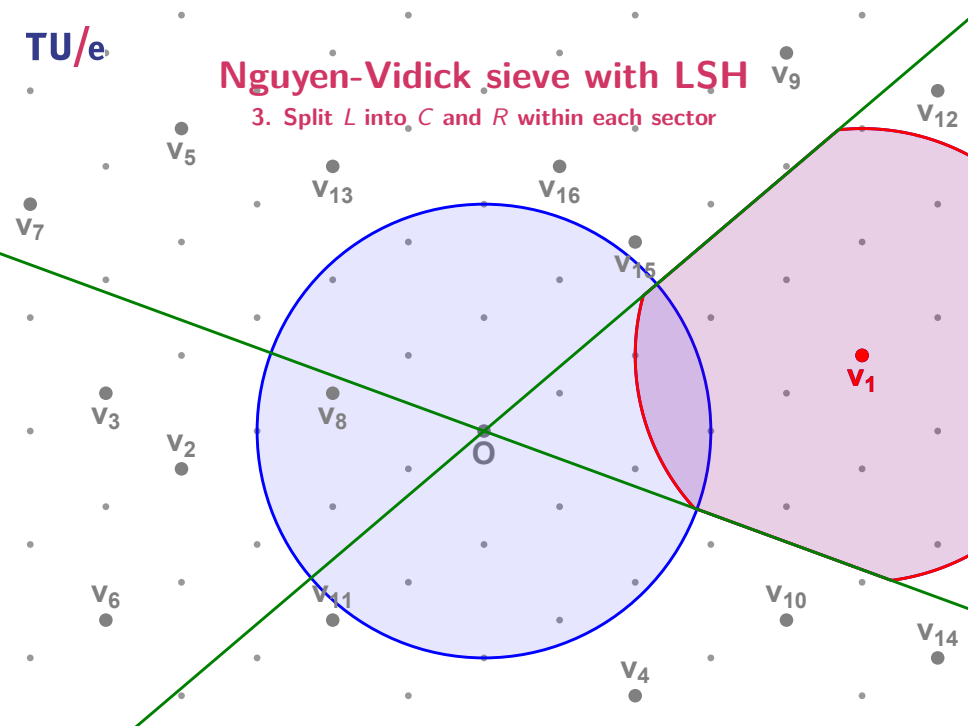


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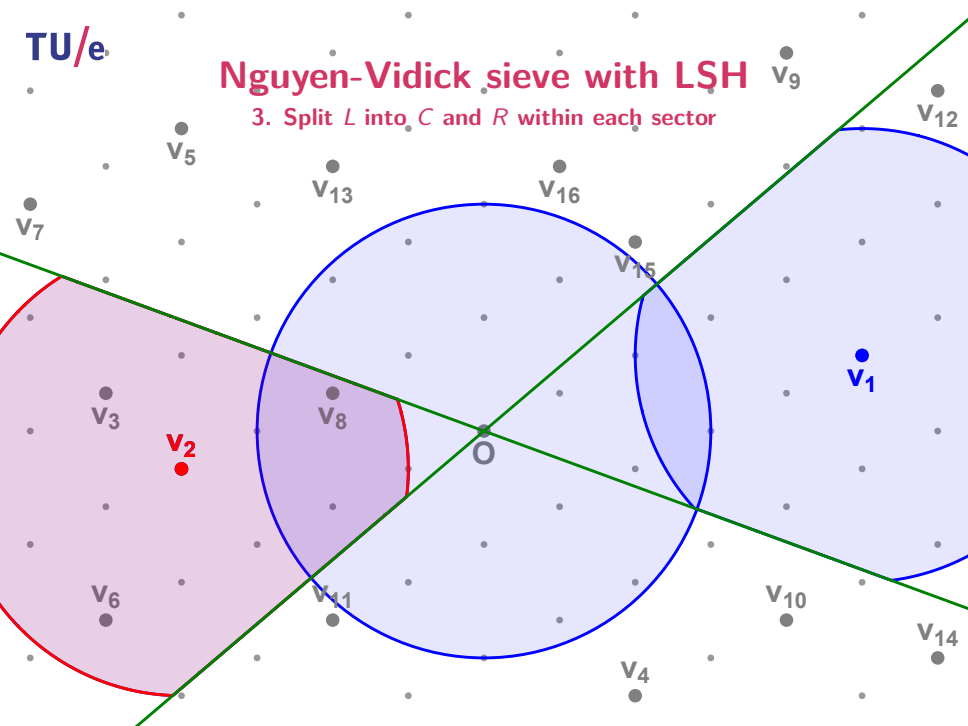
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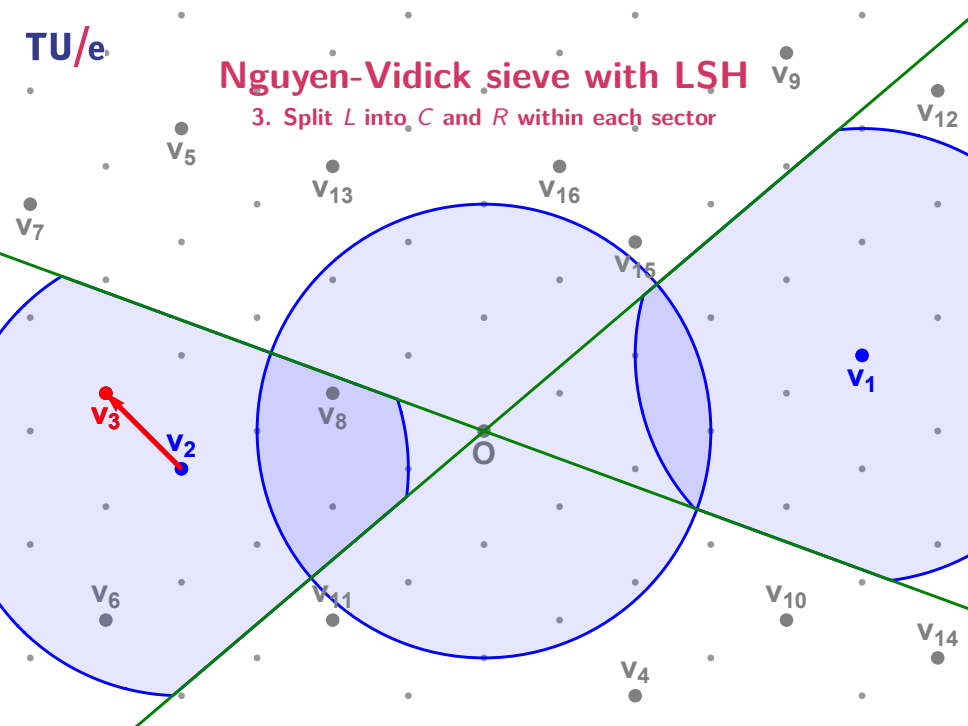
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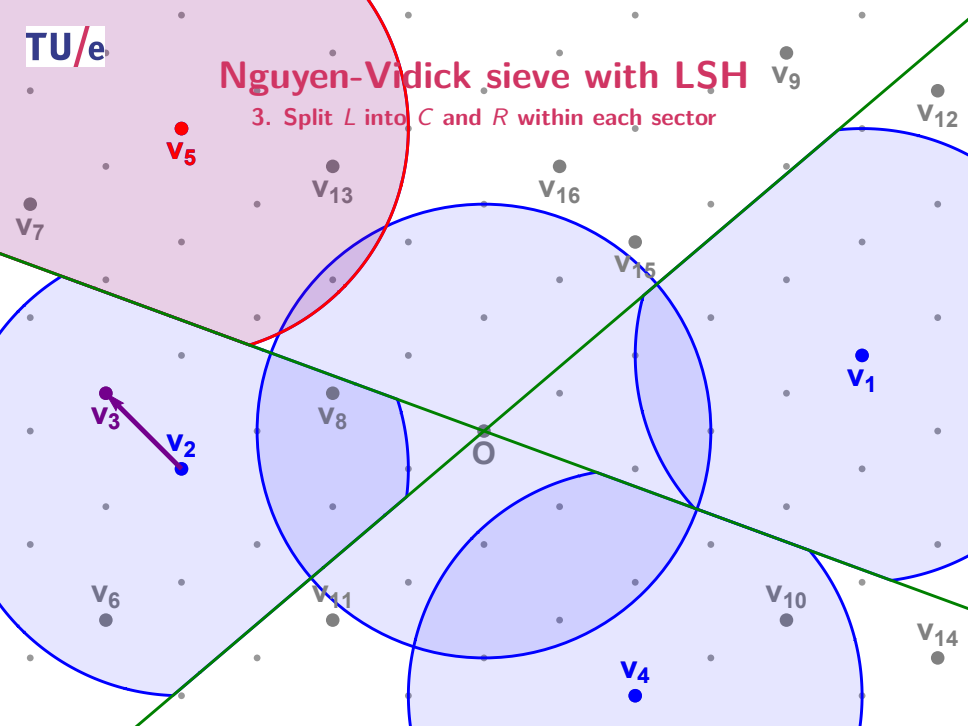
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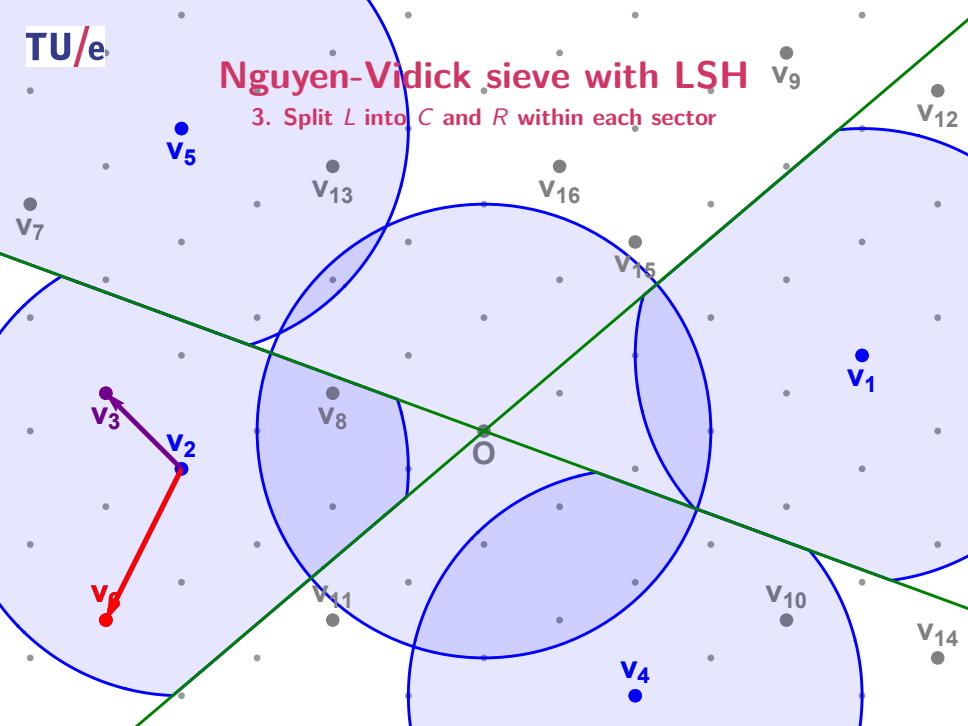


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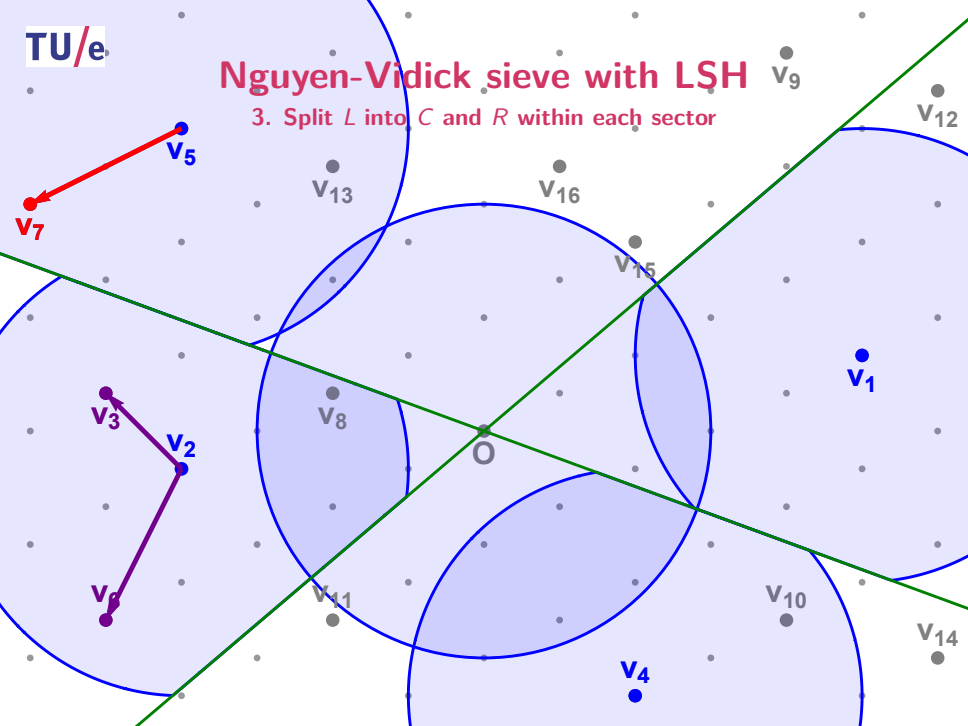




TU/e

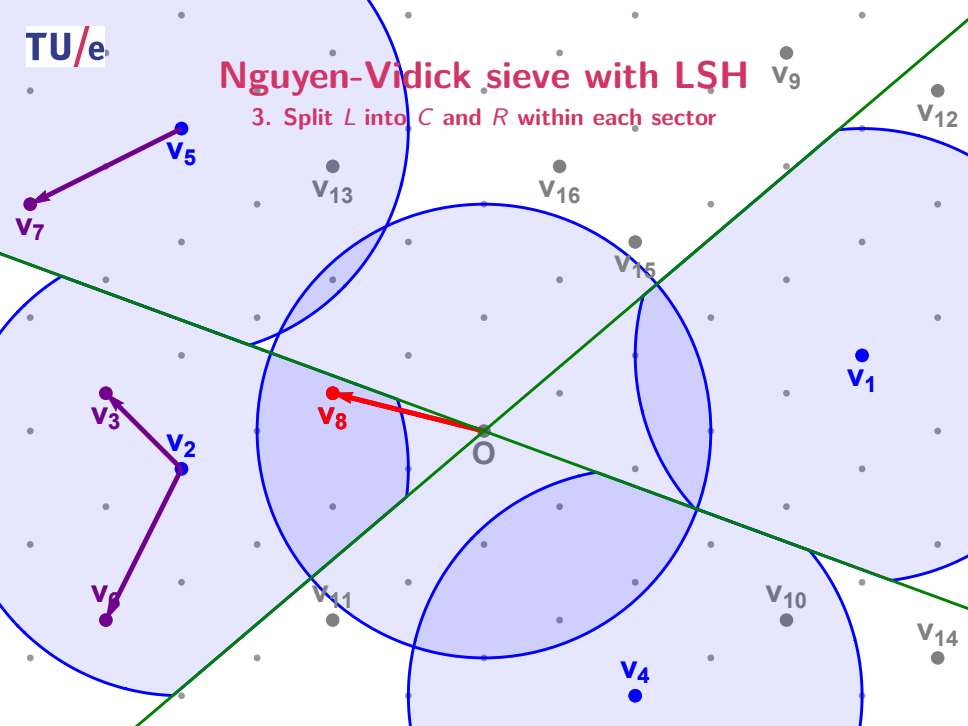
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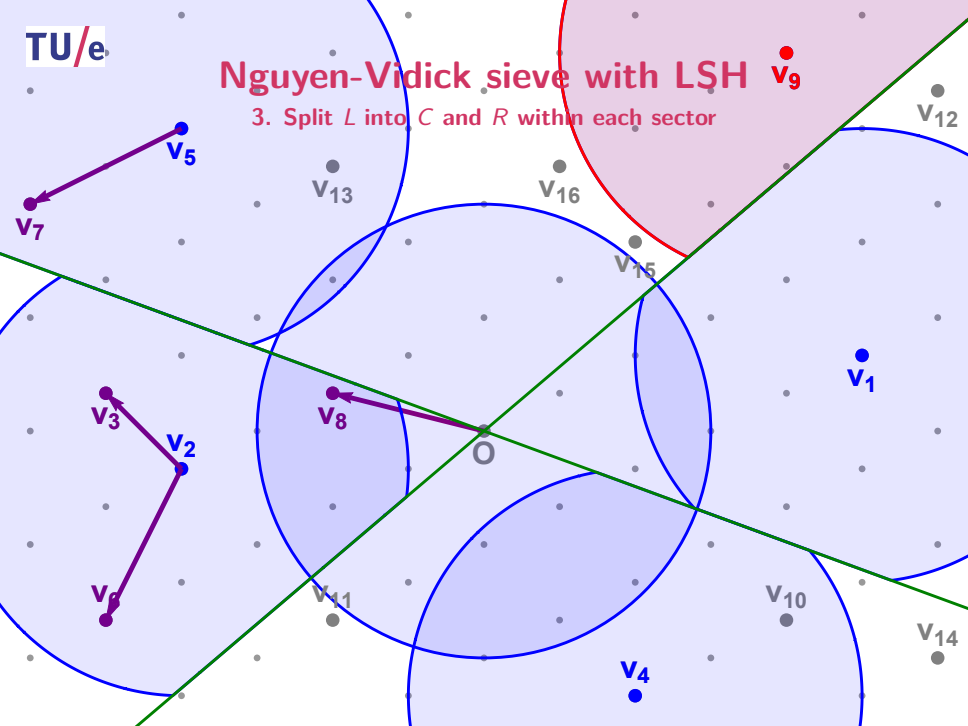
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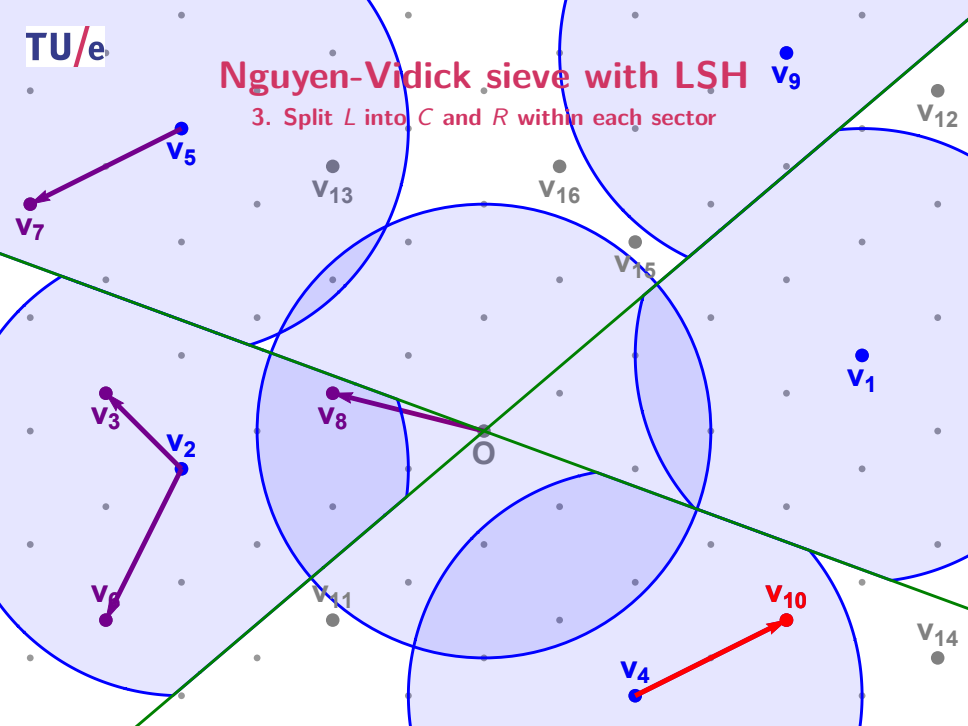
TU/e

# Nguyen-Vidick sieve with LSH $v_9$

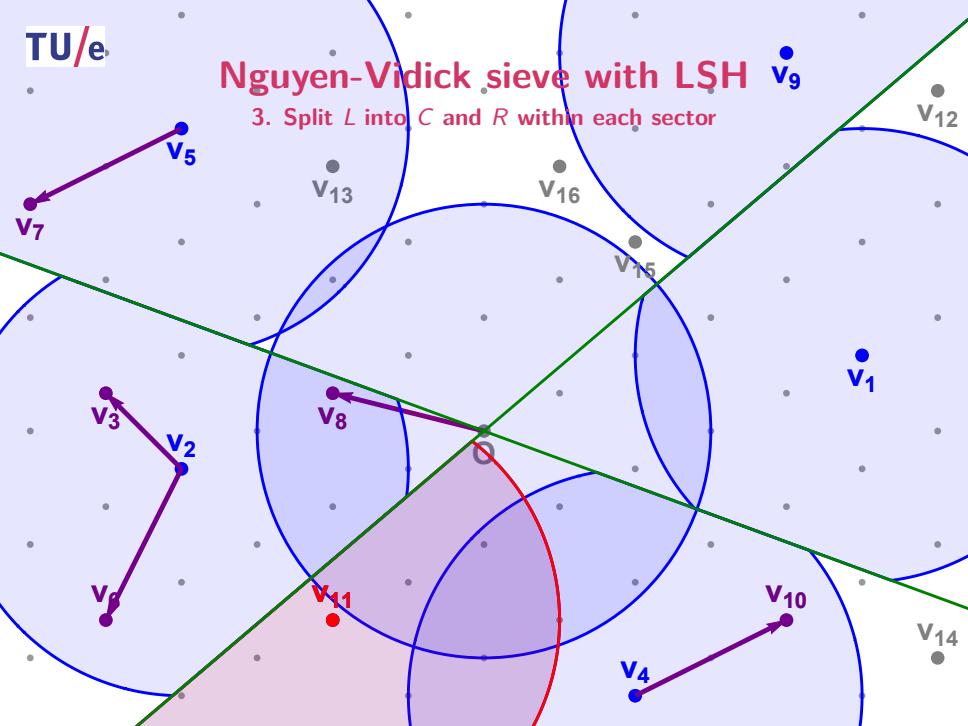
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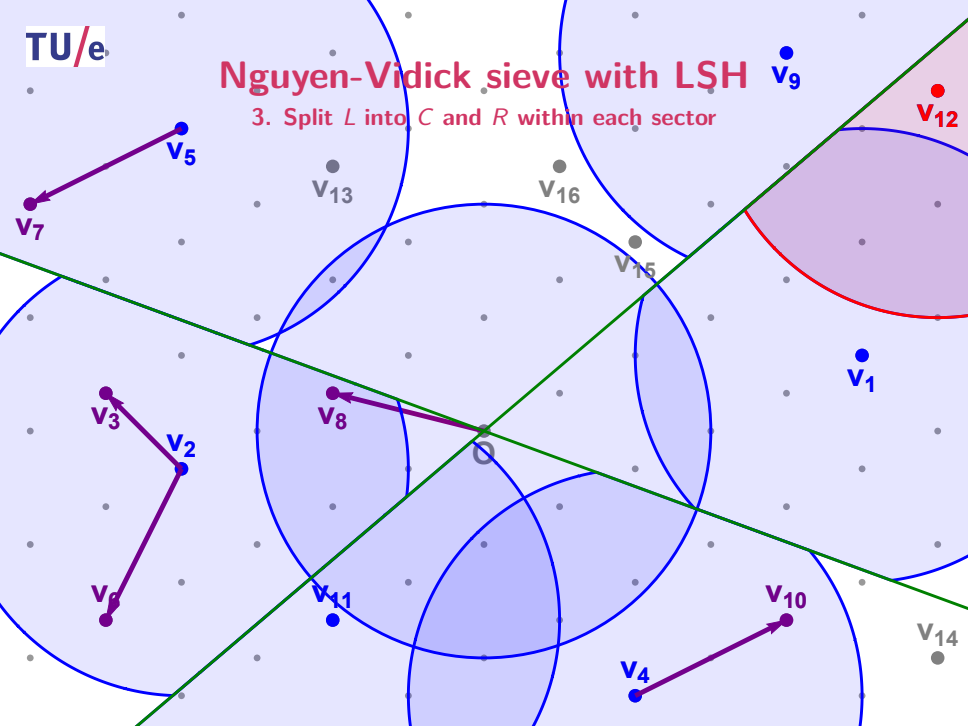
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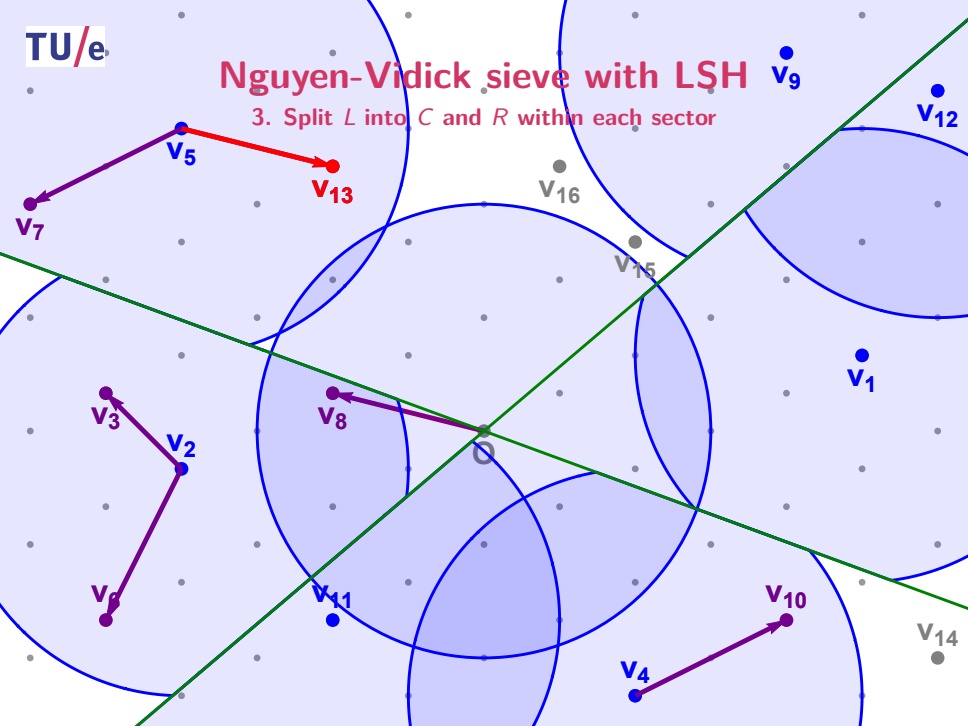
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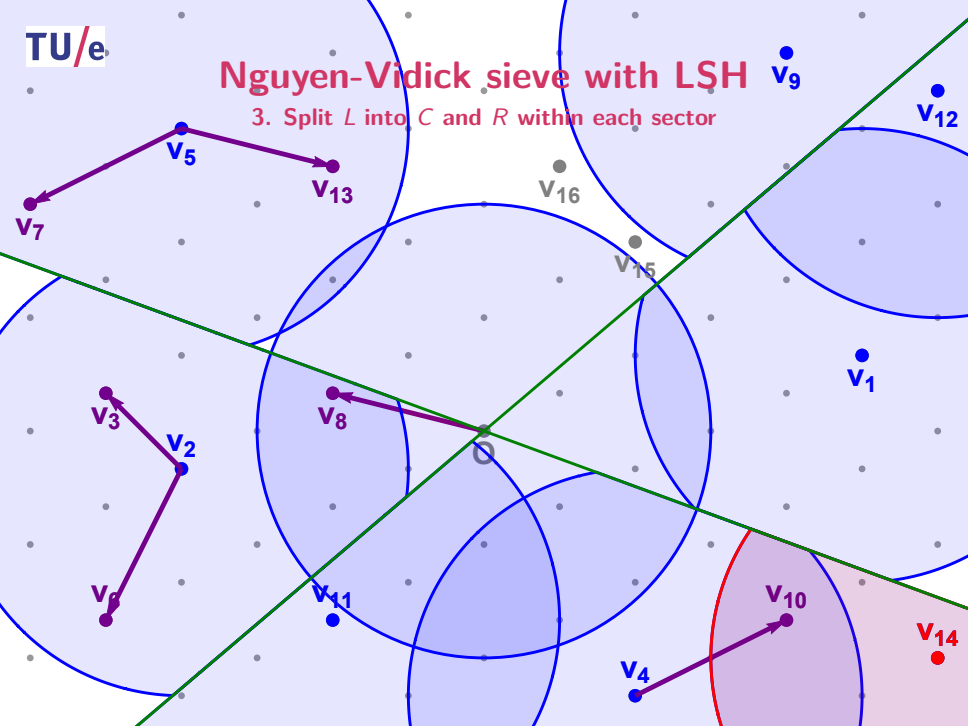
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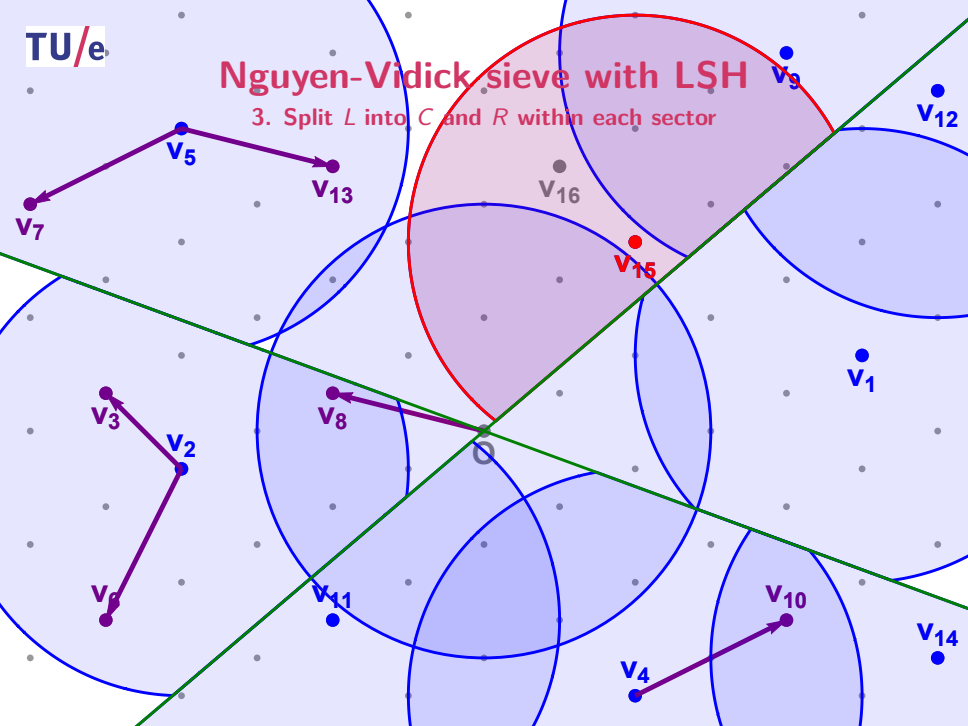
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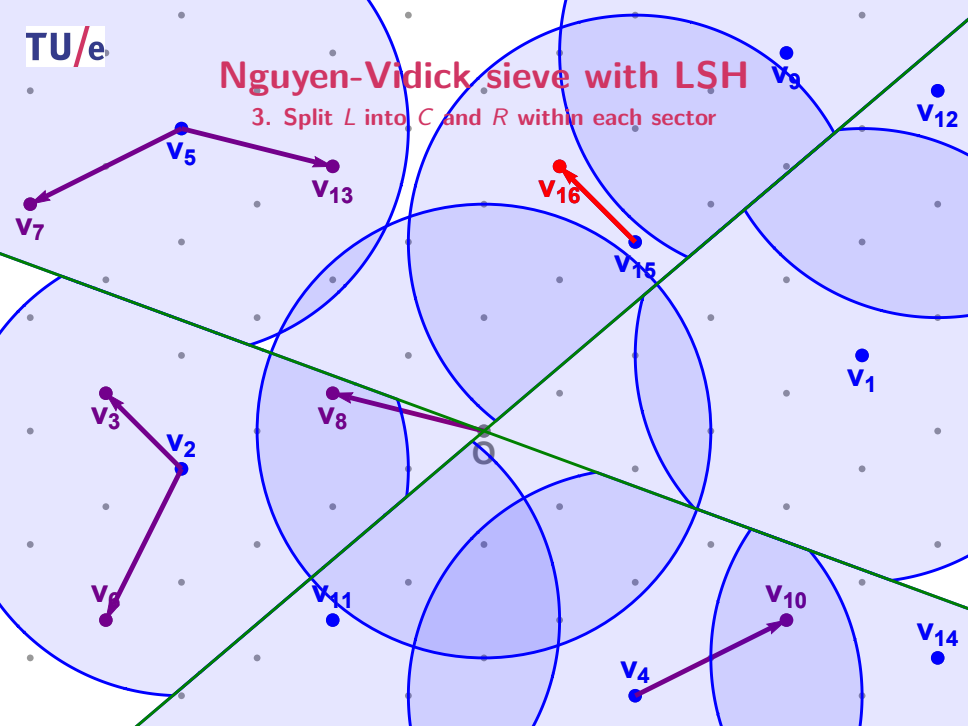


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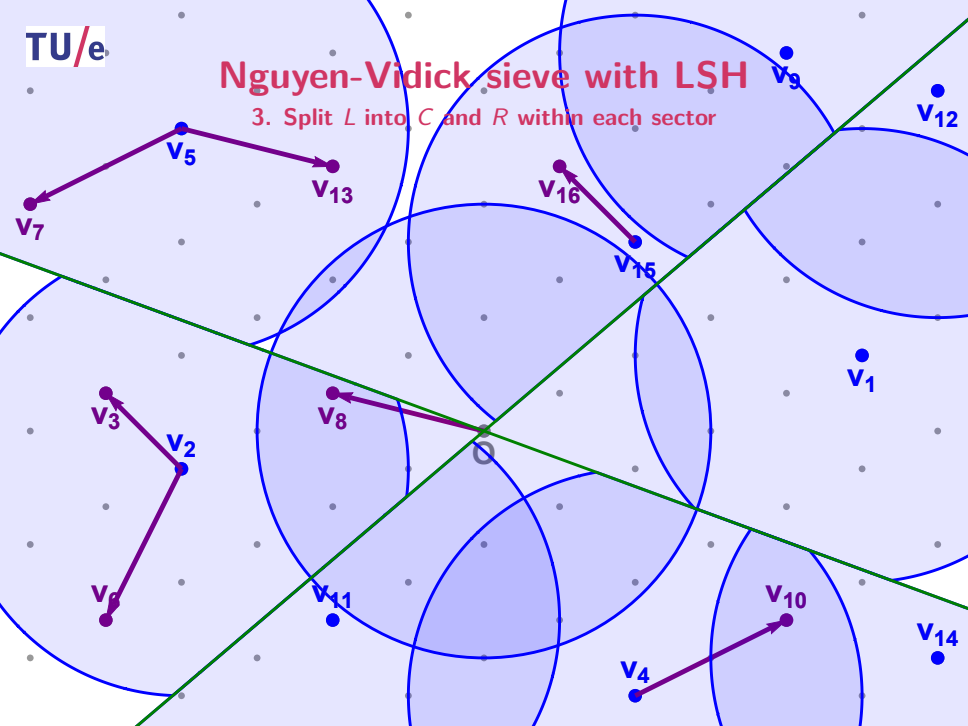
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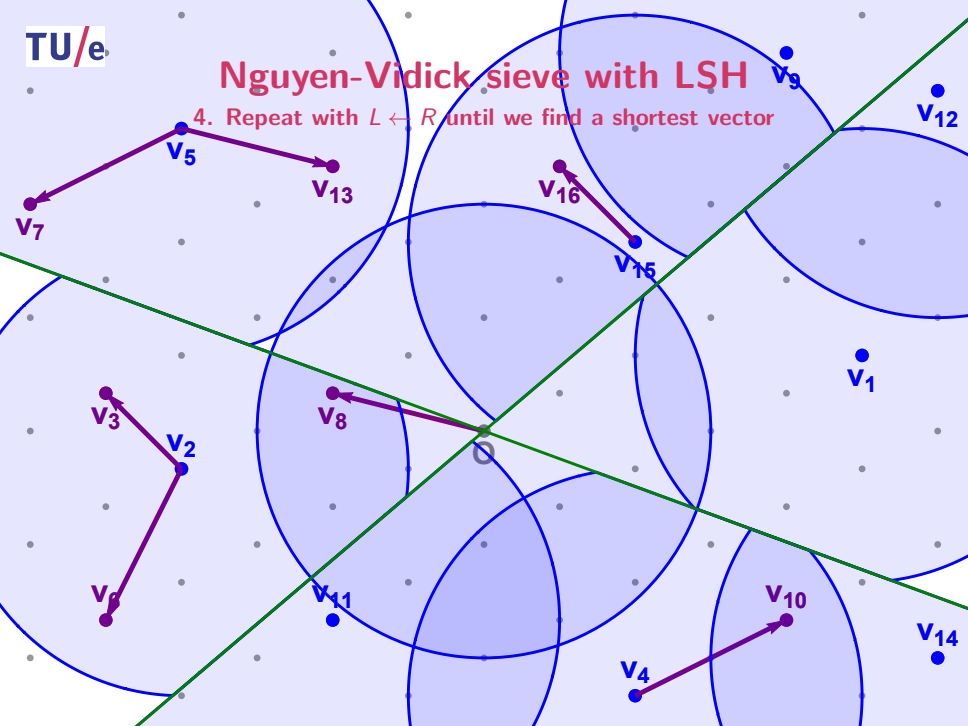


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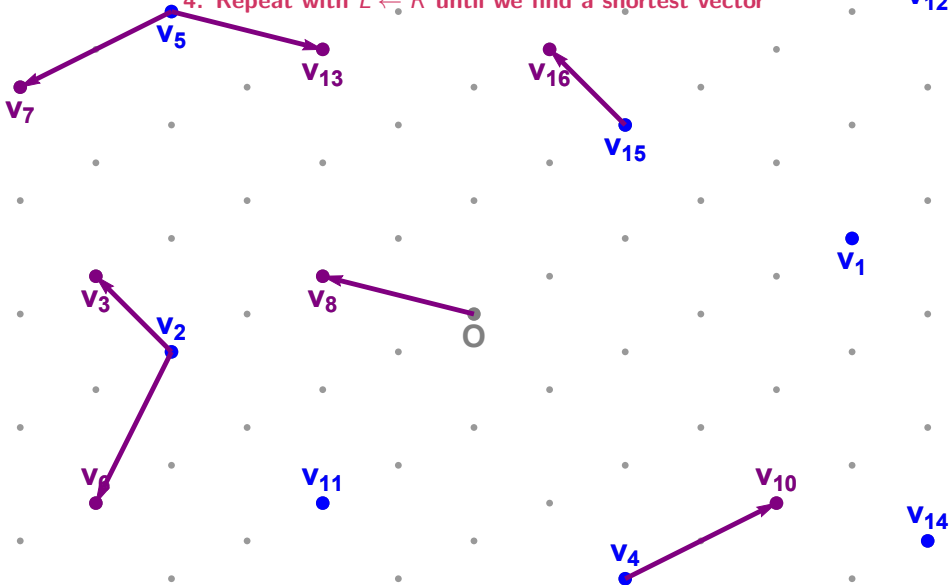


## Nguyen-Vidick sieve with LSH

4. Repeat with  $L \leftarrow R$  until we find a shortest vector

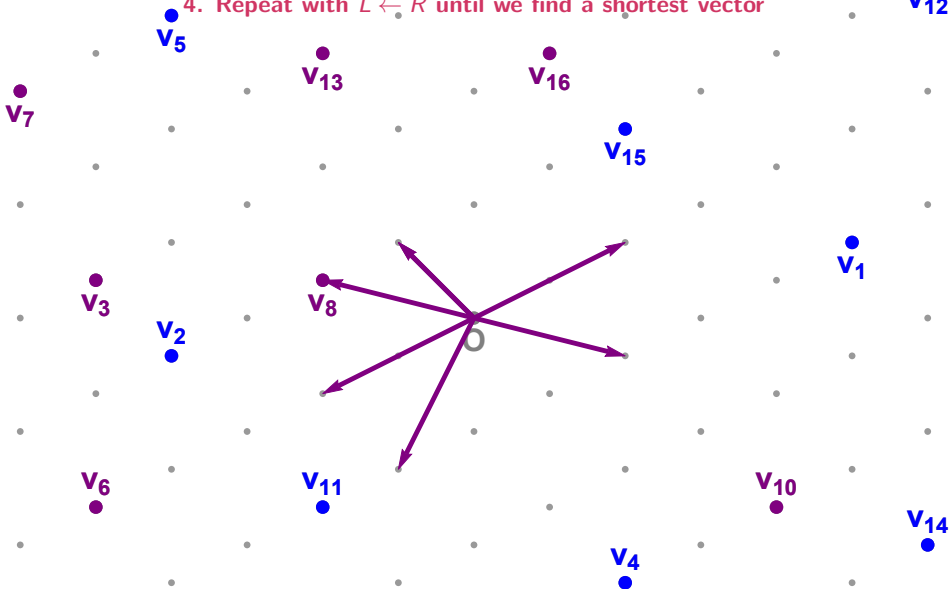
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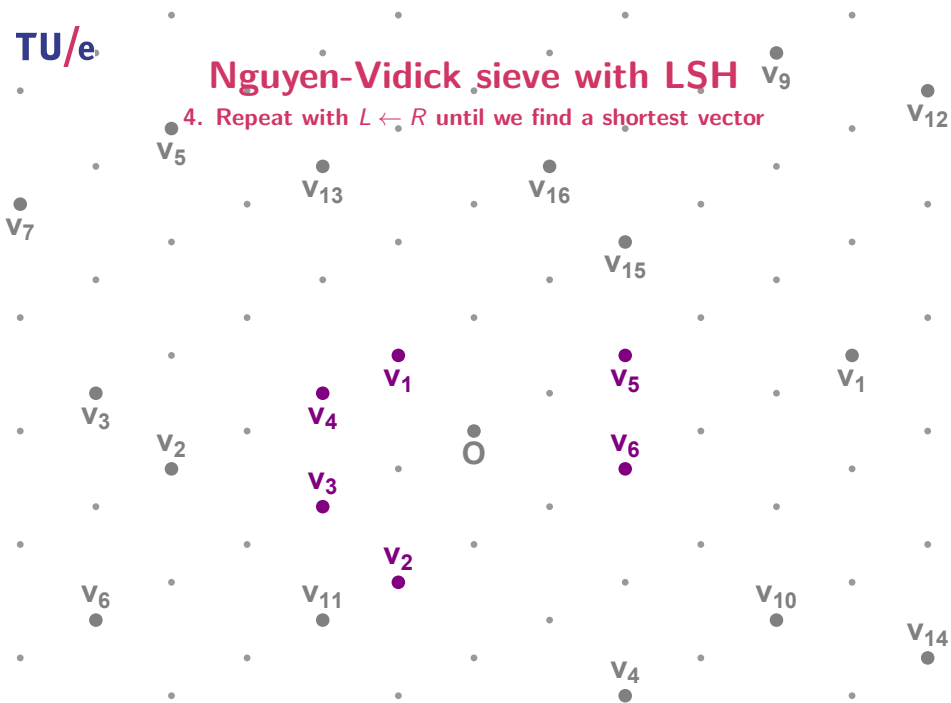
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## Nguyen-Vidick sieve with LSH

Overview





# Nguyen-Vidick sieve with LSH

## Overview

- Two parameters to tune
  - ▶  $k = O(n)$ : Number of hyperplanes, leading to  $2^k$  regions
  - ▶  $t = 2^{O(n)}$ : Number of different, independent “hash tables”

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  - ▶ Store  $2^{0.13n}$  hash tables, each containing all  $2^{0.21n}$  vectors

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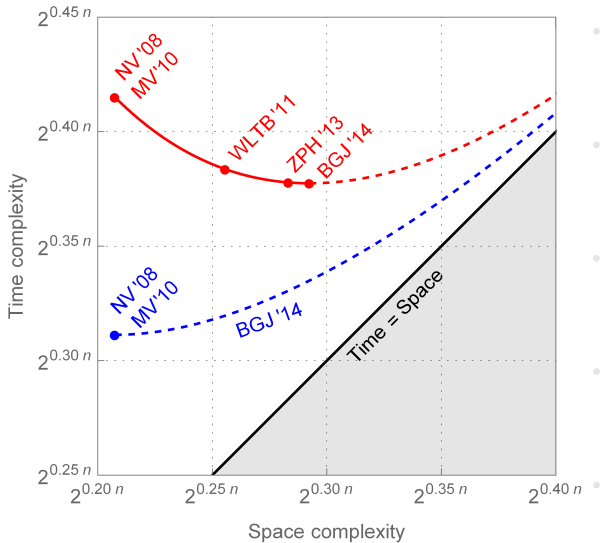
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## Heuristic

Sieving with LSH runs in time  $2^{0.34n+o(n)}$  and space  $2^{0.34n+o(n)}$ .

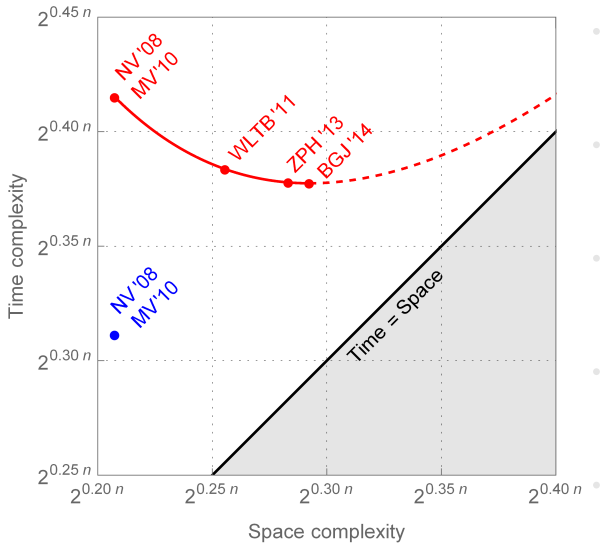
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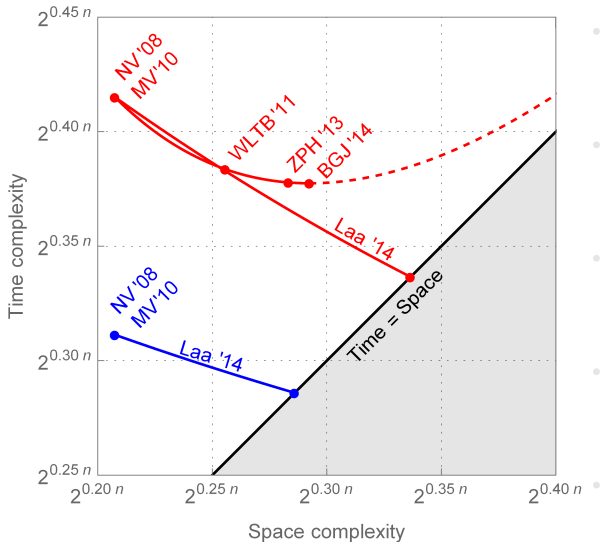
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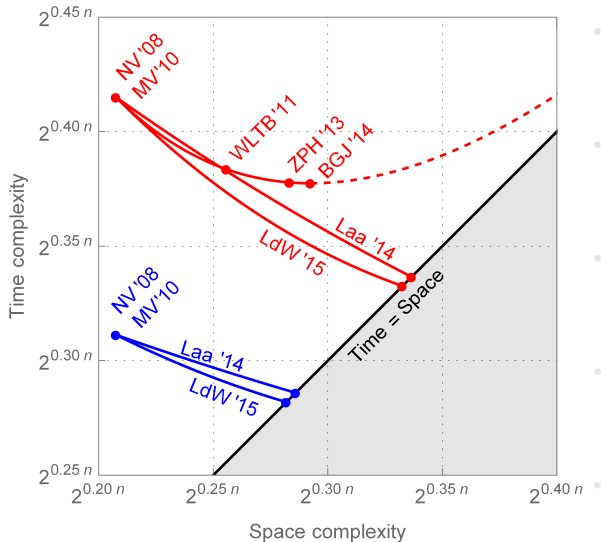
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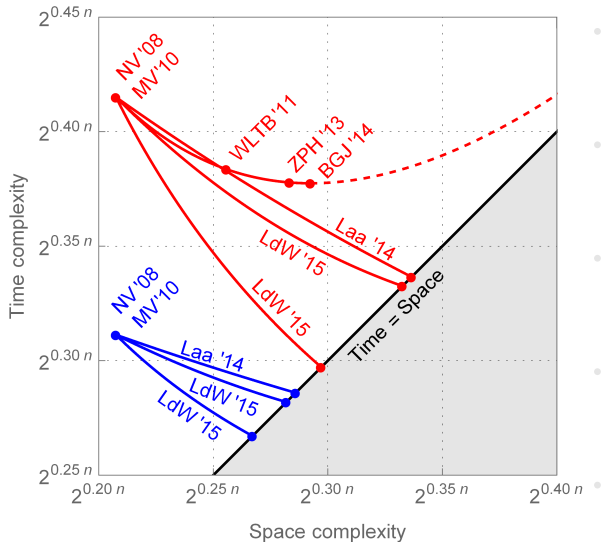
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# Sieving with LSH

Space/time trade-off



# Questions

