

ABSTRACTIONS OF THE DATA PLANE

**DIMACS Working Group on
Abstractions for Network
Services, Architecture, and Implementation**

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WHAT WE ATE FOR LUNCH

CHINESE FOOD, OF COURSE, AND LEARNED THAT . . .

*“We find comfort among those who agree with us—
growth among those who don’t.”*

THE PREVAILING ABSTRACTION OF THE DATA PLANE

APPLICATION LAYER

*applications and
mnemonic names*

TRANSPORT LAYER

*reliable (or
unreliable) transport*

NETWORK LAYER

*best-effort global
packet delivery*

LINK LAYER

*best-effort local
packet delivery*

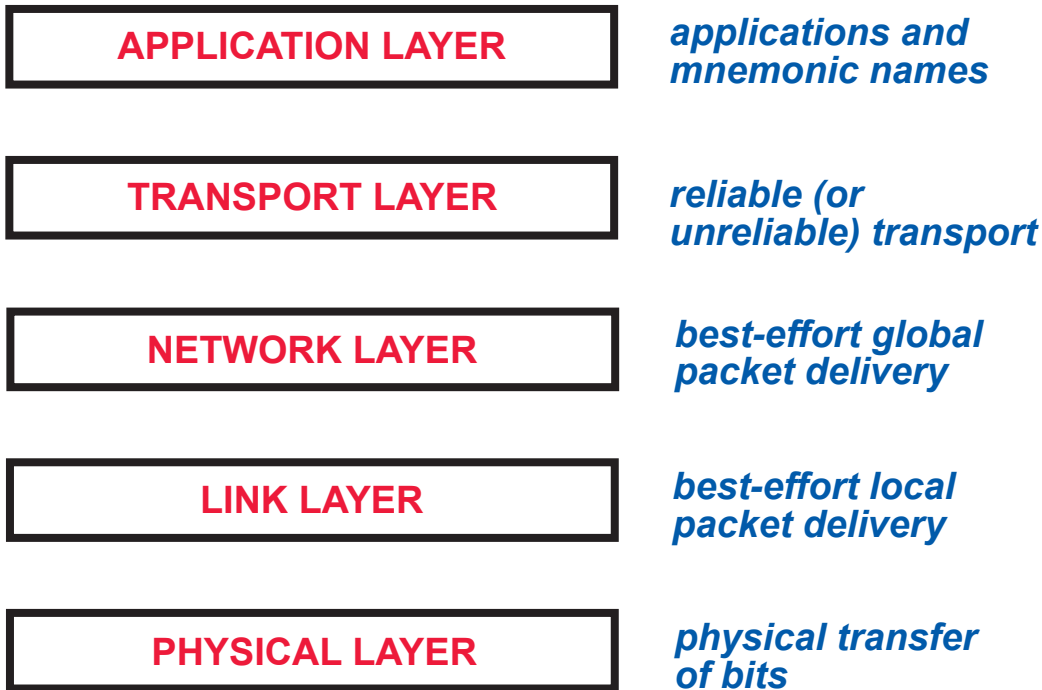
PHYSICAL LAYER

*physical transfer
of bits*

abstractions from “The future of networking, and the past of protocols”

[Shenker 2011]

WHY SHOULD WE QUESTION THIS?



1

Because there are many serious problems with the current Internet, and we must look at all possible solutions.

2

Because the purpose of the control plane is to manage the data plane, so a well-structured data plane may be the key to a well-structured control plane.

For example, “An axiomatic basis for communication” is intended to formalize what routers do . . .

. . . but much of the space is devoted to a careful discussion of the behavior of the data plane.

[Karsten, Keshav, Prasad & Beg 2007]

WHY SHOULD WE QUESTION THIS?

3 Because it is not realistic.

headers in a typical AT&T packet

Cloud Service
HTTP
TCP
IP
IPsec
IP
GTP (QoS, billing)
UDP
IP
MPLS
MPLS
Ethernet

15+ load balancing / routing algorithms are involved in getting this packet to its destination . . .

. . . most with different goals in mind;

. . . most have been analyzed / designed in some state of isolation;

. . . all are getting more dynamic every day

from “Cloud computing and my worries about the network that enables it”

[Spatscheck 2010]

A BETTER ABSTRACTION OF THE DATA PLANE?

this is more realistic, . . .

*. . . but consensus would
be difficult to achieve . . .*

. . . and not long-lasting

APPLICATION LAYER

MIDDLEWARE LAYER

TRANSPORT LAYER

LISP LAYER

NETWORK LAYER

MPLS LAYER

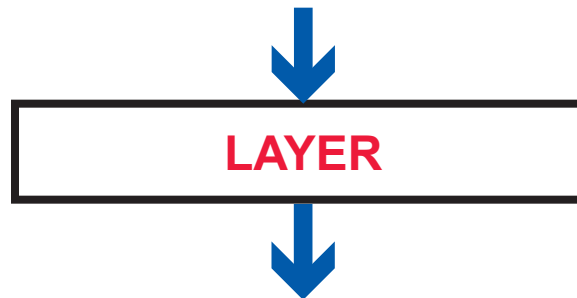
LINK LAYER

PHYSICAL LAYER

A DIFFERENT VIEW OF THE DATA PLANE

Each layer is a distributed system with the same abstract functionality and the same abstract state.

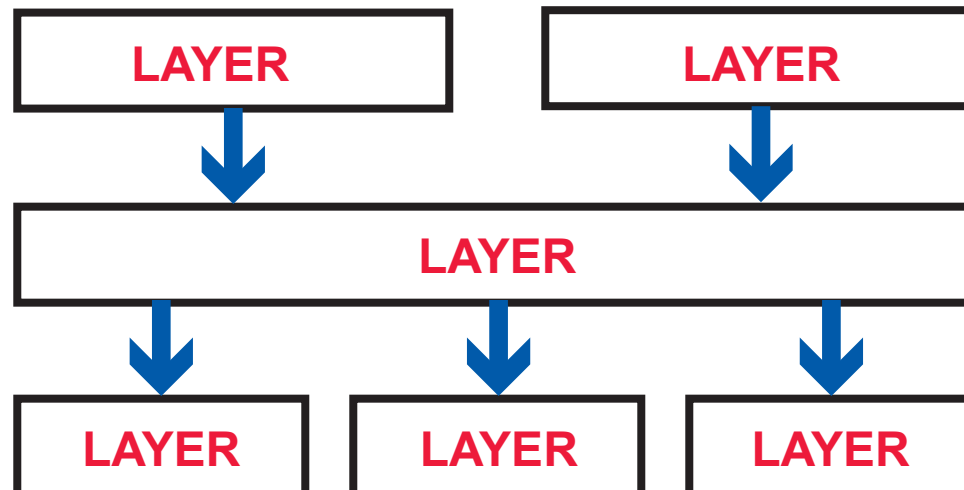
upper interface is a specification of communication services (provided)



includes transport, routing, and forwarding

lower interface is a specification of communication services (used)

This pattern is instantiated many times in a network architecture, for many purposes, at many levels, and with many different scopes.

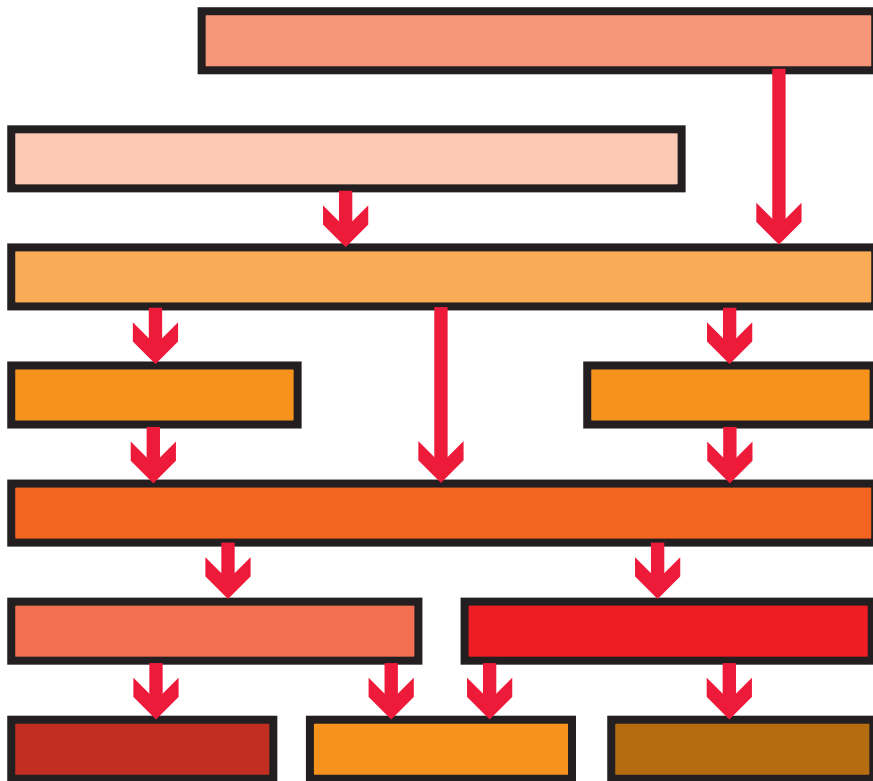


this hypothesis comes from *Patterns in Network Architecture*

[Day 2008]

WE CALL THIS THE “GEOMORPHIC VIEW” OF NETWORKS . . .

... BECAUSE THE ARRANGEMENT OF
LAYERS RESEMBLES THE EARTH’S CRUST



it is inspired by Day's ideas, with many
changes in terminology and (we hope)
improvements

OUTLINE

1 Basic information about layers

2 Frequently-asked questions

3 Examples

4 Summary and conclusions

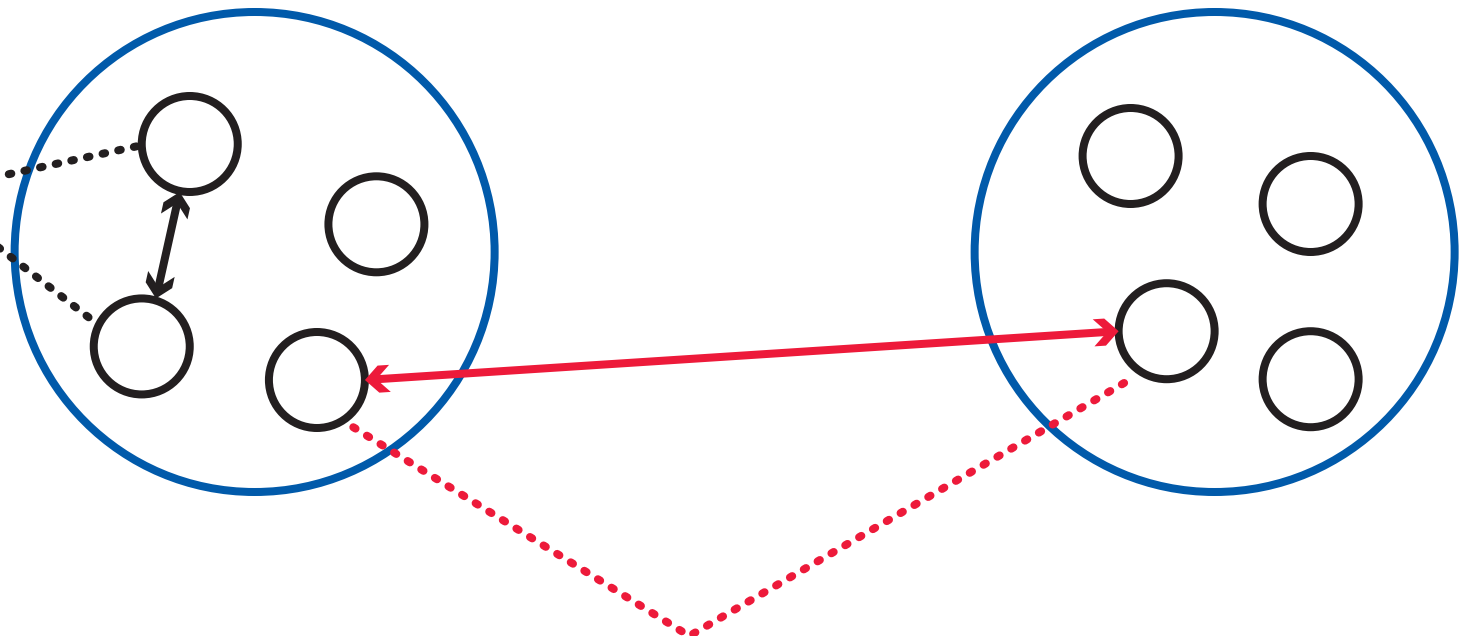
LAYERS: MACHINES AND PROCESSES

machine

machine

the OPERATING SYSTEM creates these processes and enables them to communicate quickly and reliably with each other

this communication is assumed as a building block



these processes can only communicate through a NETWORK, with all of the challenges we know well

we can choose to regard a virtual machine as a machine . . .

. . . and to regard communication through the hypervisor and softswitch of a physical machine as networked communication, and an object of study

LAYERS: MEMBERS, NAMES, AND ROUTING

a member is a process that represents its machine in that layer

each layer has its own name space

a member has a name that is unique and permanent (although re-usable)

OVERLAY
(higher layer)

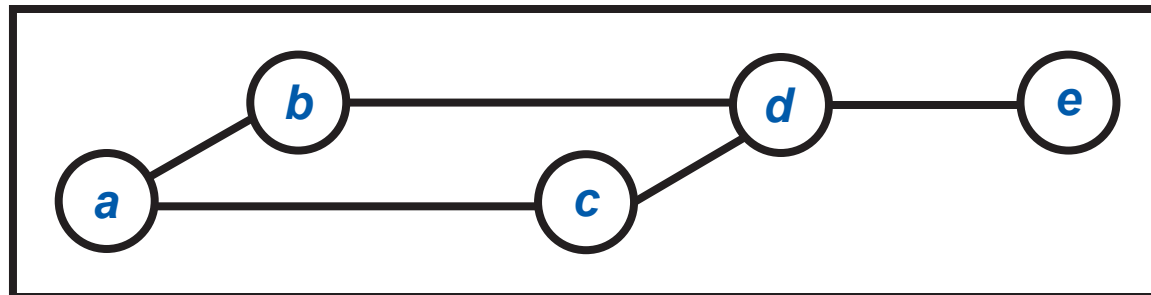


member process

link

a link is an instance or usage of a communication service

UNDERLAY
(lower layer)



members are connected to each other by links

because there is usually not a link between each pair of members, routes tell members how to reach each other

routing protocol maintains routes as links change

LAYERS: REGISTRATIONS

a registration maps an overlay process to an underlay process

both processes are on the same machine

the underlay process is a process in the lower layer that represents the overlay process to the network

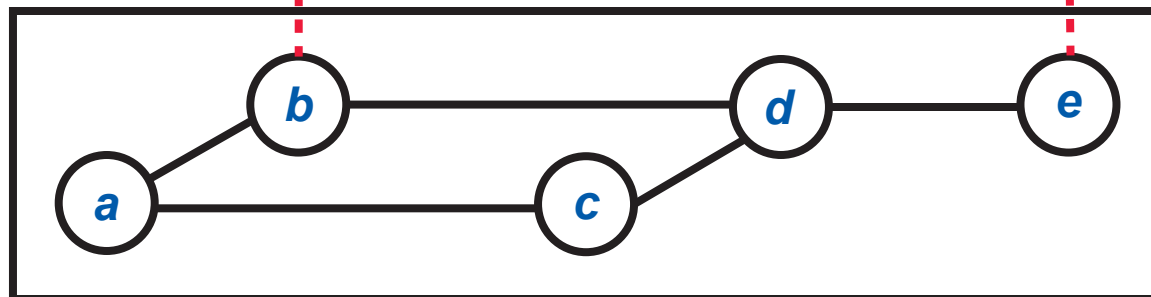
OVERLAY
(higher layer)



here these registrations are *attachments*

.....registrations.....

UNDERLAY
(lower layer)



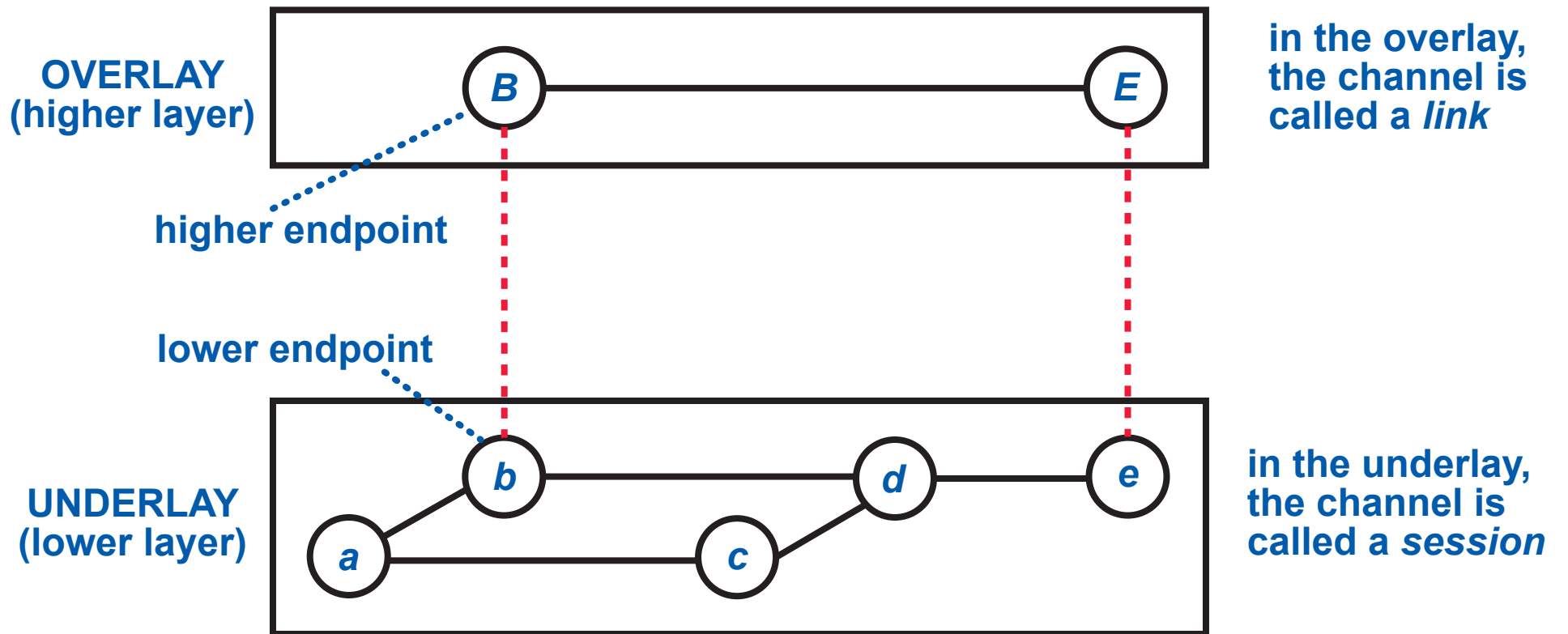
here these registrations are *locations*

registrations can be created or destroyed by either layer

LAYERS: CHANNELS

a channel is an instance or usage of a communication service

a channel can be implemented as a service by an underlay for an overlay



in the overlay, the channel is called a *link*

in the underlay, the channel is called a *session*

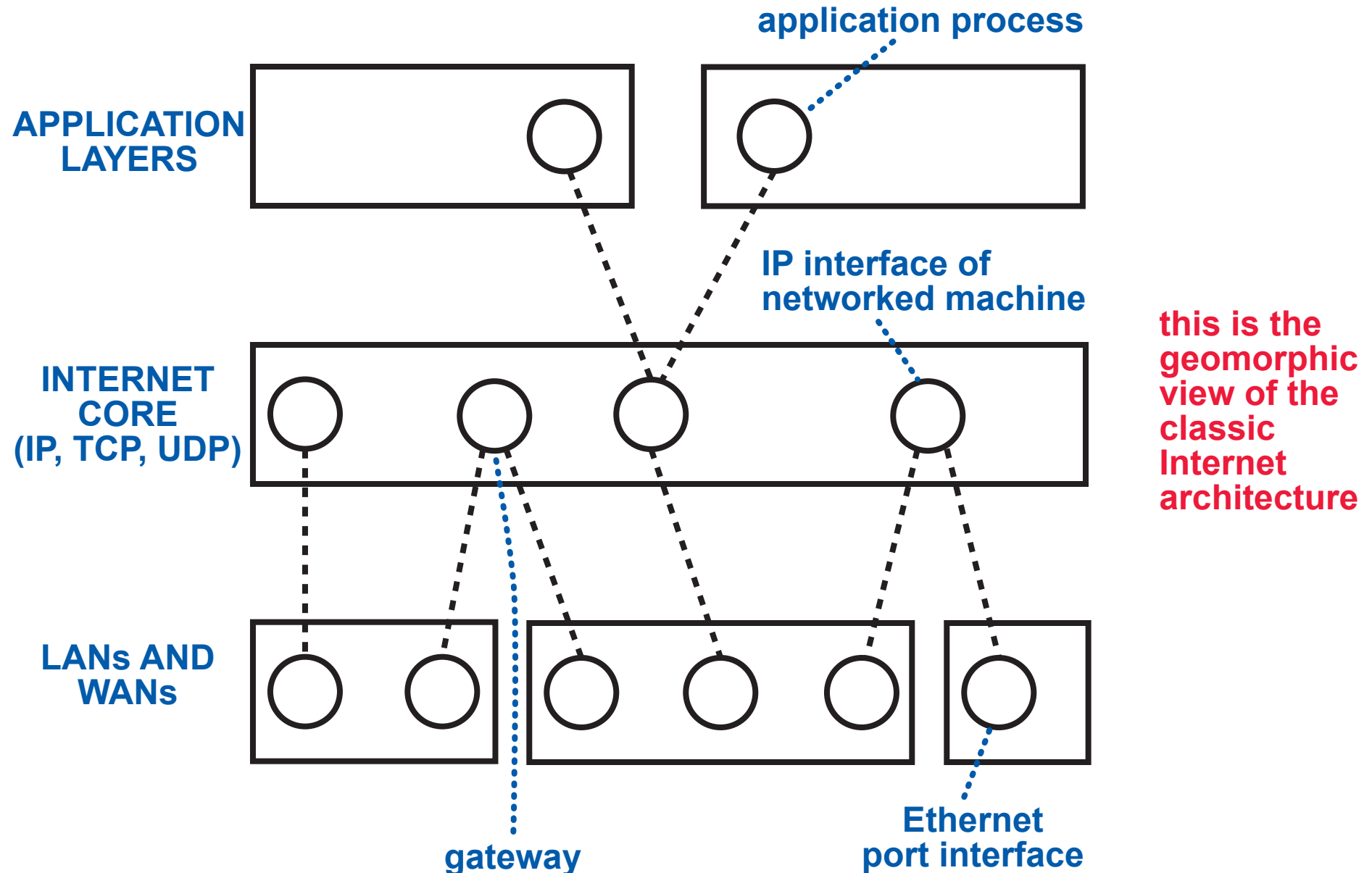
when *b* receives a channel request from *B* for *E*, it uses *locations* to find that *E* is located at *e*

underlay includes a transport protocol that enforces the service specification

LAYERS: SCOPE AND LEVEL

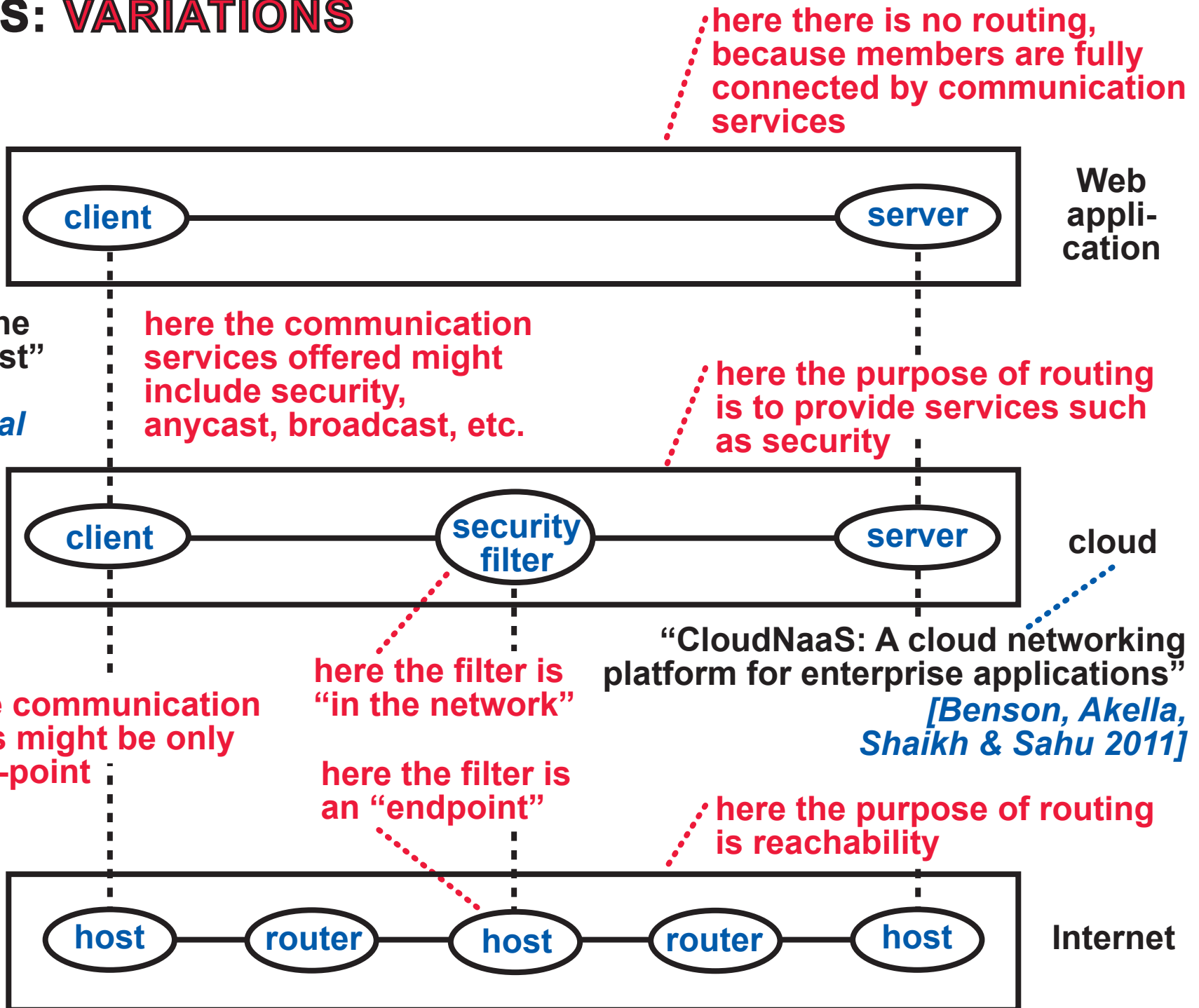
layers are arranged in a usage hierarchy, which defines levels

the scope of a layer is the set or class of processes that could be members



LAYERS: VARIATIONS

“The end-to-end argument and application design: The role of trust”
[Clark & Blumenthal 2011]



LAYERS: SOFTWARE STATE OF A LAYER

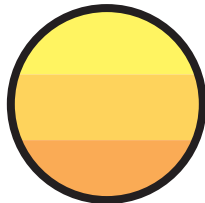
a layer is a distributed software system

this is a snapshot of its distributed, dynamic state

overlay



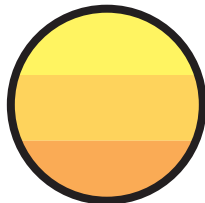
related to serving overlays



locations: set Registration

sessions: set Channel

strictly internal



members: set Process

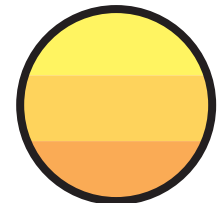
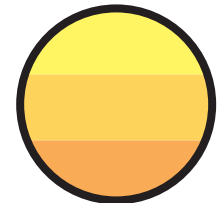
forwarding: set Route

depends on

related to using underlays

attachments: set Registration

links: set Channel



underlay



underlay

OUTLINE

1 Basic information about layers

2 Frequently-asked questions

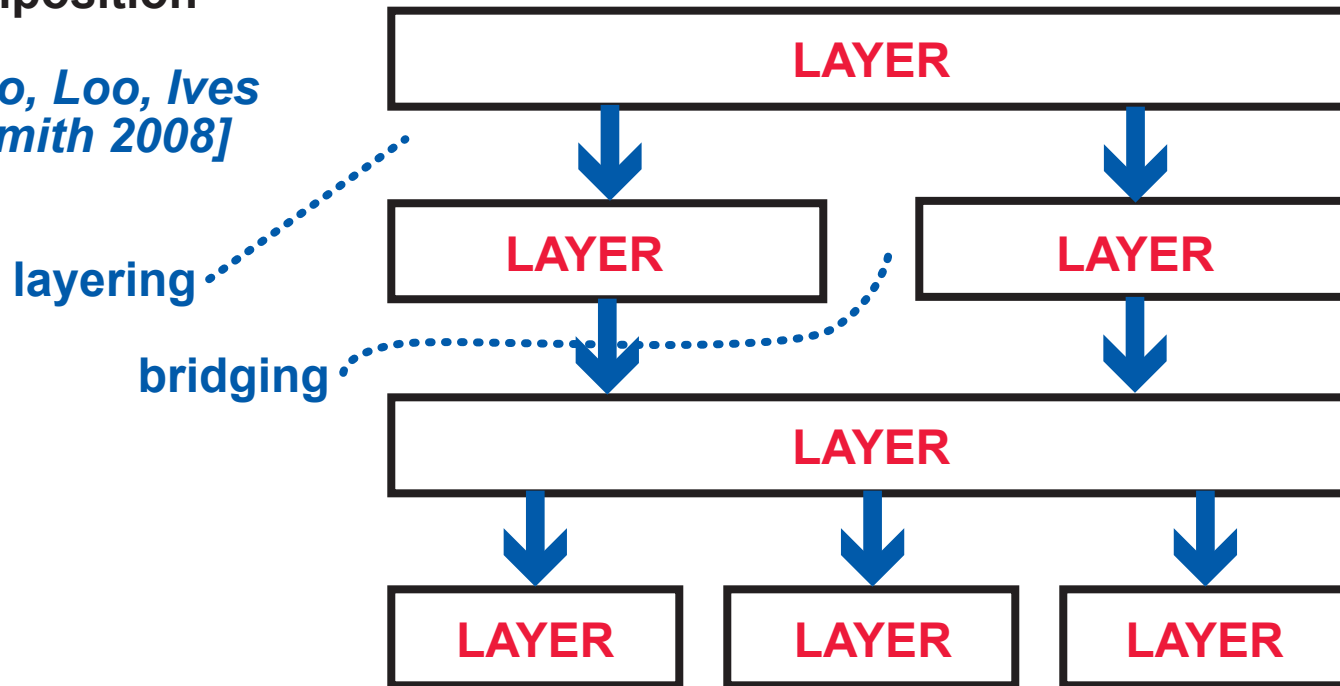
3 Examples

4 Summary and conclusions

FAQ: HOW IS THE GEOMORPHIC VIEW DIFFERENT FROM OVERLAYS?

“MOSAIC: Unified declarative platform for dynamic overlay composition”

[Mao, Loo, Ives & Smith 2008]



geomorphic view attempts to explain what is in each layer, as well as how they compose

geomorphic view has no unique reference point, so there is nothing for an “overlay” to be “over”

FAQ: IS THE GEOMORPHIC VIEW DESCRIPTIVE OR PRESCRIPTIVE?

FUNCTIONALLY, IT IS DESCRIPTIVE

there should be no major function or design that cannot be described

HOWEVER, THERE ARE FEWER MECHANISMS THAN ARE FOUND "IN THE WILD"

no arguing about names vs. identifiers vs. locators vs. addresses—each layer has one name space, designed and used for the purposes of the layer

no tunneling used as an intra-layer exception to the routing system—just inter-layer interfaces

FEWER MECHANISMS COULD MEAN:

- each design has exactly one correct description
- designs can be compared easily
- it is possible to map out structured spaces of design trade-offs
- it is possible to get implementations by code generation and re-use

GOAL IS TO CHOOSE THE MECHANISMS THAT ARE THE BEST BECAUSE THEY FACILITATE . . .

- . . . composition—of layers, mechanisms within a layer, or reasoning methods
- . . . separation of concerns, so that diverse goals can be met without interfering with each other

OUTLINE

1 Basic information about layers

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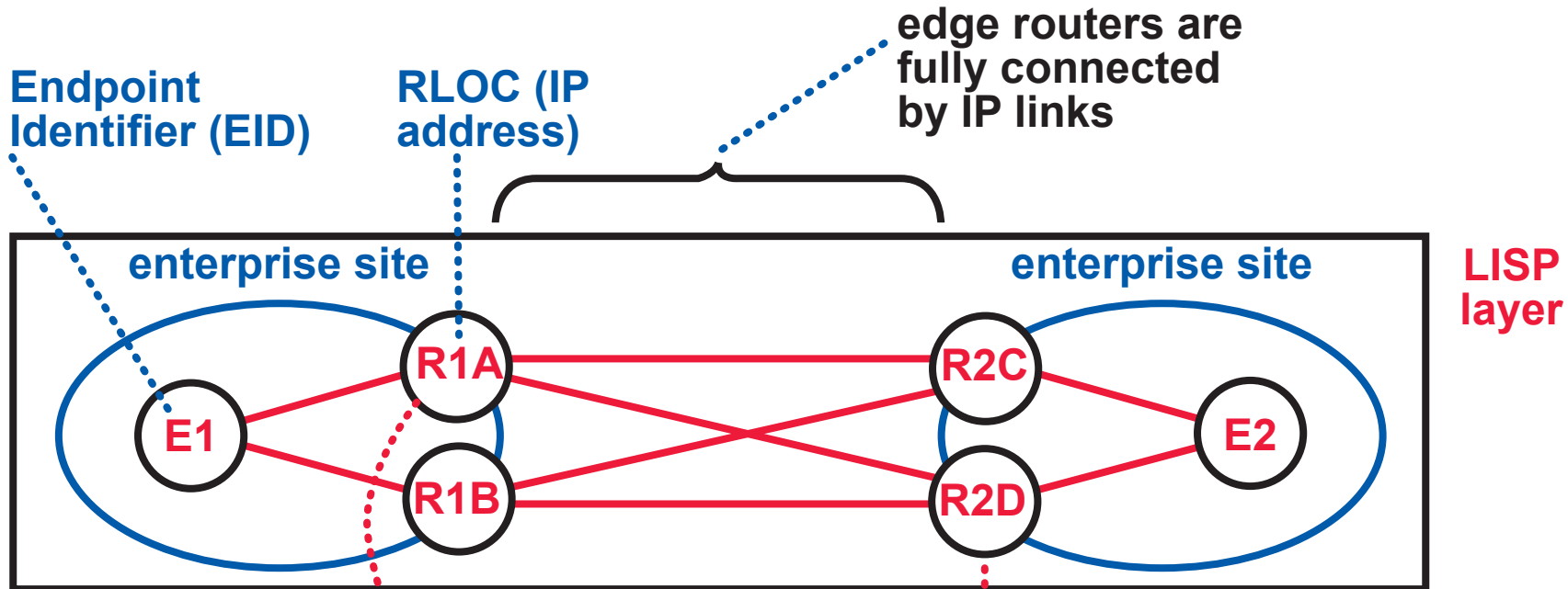
3 **Examples**

4 Summary and conclusions

EXAMPLE: COMPARING RESEARCH RESULTS

WHAT DO LISP AND SEATTLE HAVE IN COMMON?

Locator/Identifier Separation Protocol



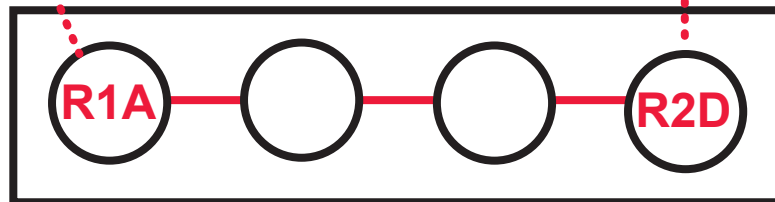
Endpoint Identifier (EID)

RLOC (IP address)

edge routers are fully connected by IP links

LISP layer

IP layer

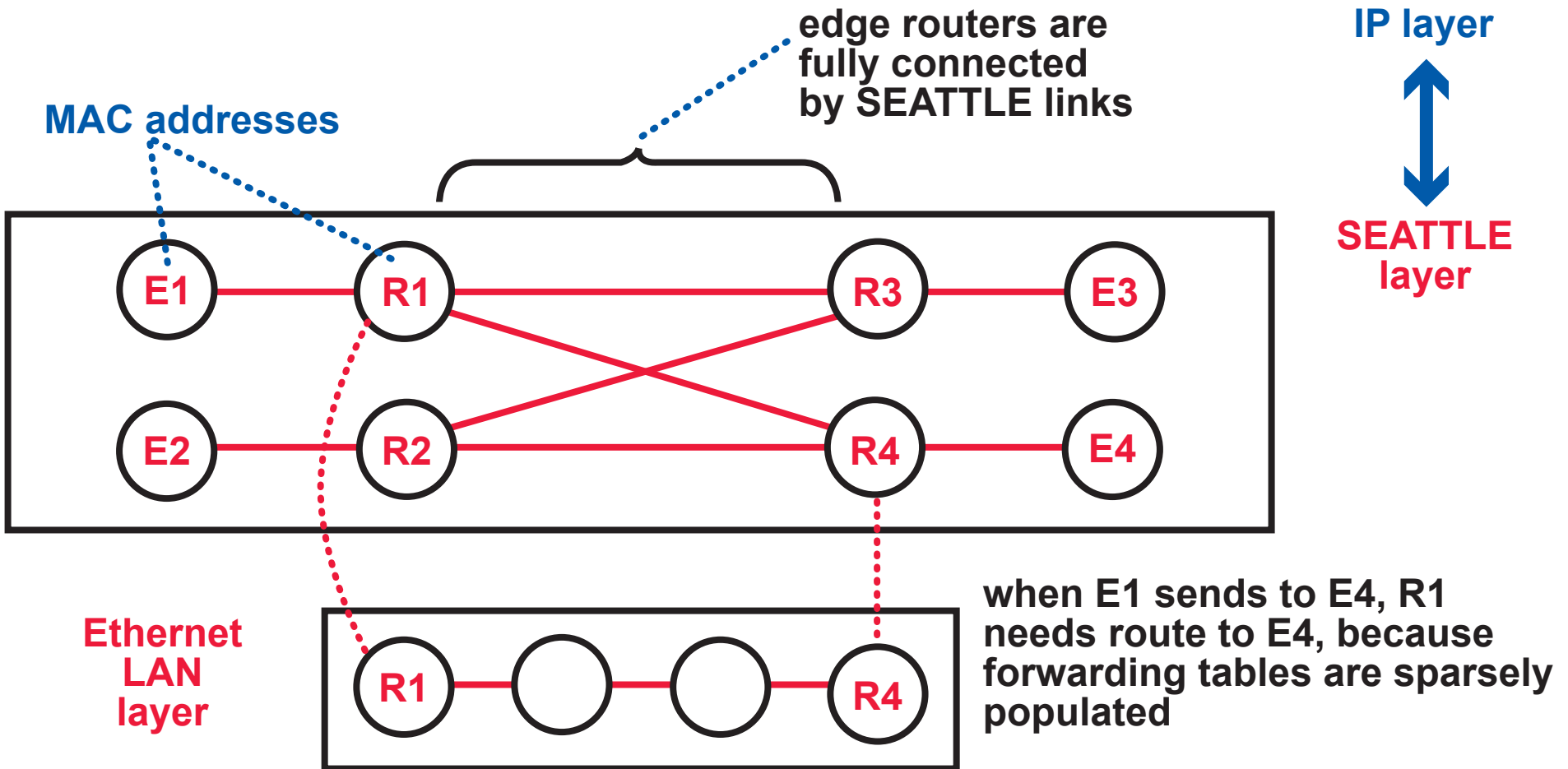


when E1 sends to E2, R1A or R1B needs route to E2, because forwarding tables are sparsely populated

route is (routes are) the same from every router, so the route can be obtained by directory lookup

WHAT DO LISP AND SEATTLE HAVE IN COMMON?

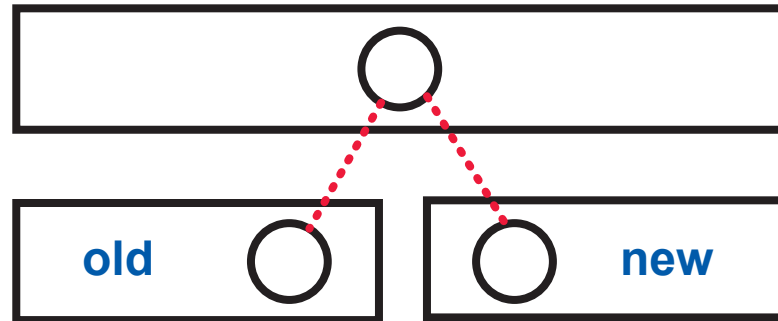
“Floodless in SEATTLE: A scalable Ethernet architecture for large enterprises”
[Kim, Caesar & Rexford 2008]



Nicira networks also have this structure; comparison focuses attention on the difference, which is how directories are implemented

EXAMPLE: COMPOSITION OF MOBILITY MECHANISMS

AS A PROBLEM, NETWORK MOBILITY IS A CHANGE IN REGISTRATION ...



... of a process, while it is participating in a channel

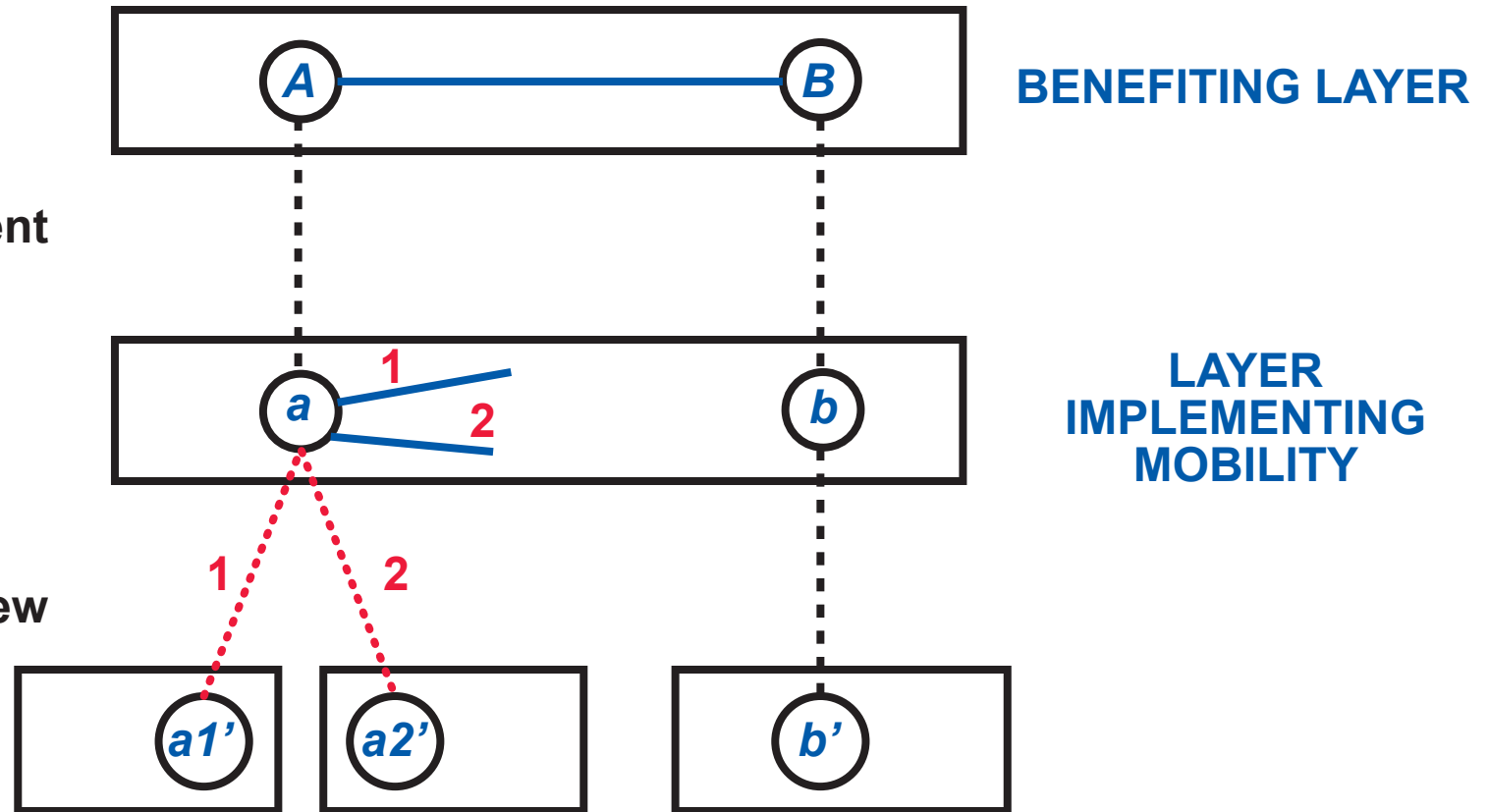
AS A SOLUTION, NETWORK MOBILITY MAINTAINS CHANNELS, DESPITE THE MOBILITY OF THE PROCESSES PARTICIPATING IN THEM



MOBILITY IMPLEMENTATION: ATTACHMENT MOBILITY

1
a becomes disconnected because attachment to **a1'** and links implemented through **a1'** fail

2
a forms a new attachment and new links through **a2'**



the hard work of implementing attachment mobility is re-routing so that other processes can reach **a** through new links

examples:

LANs/VLANs

Mobile IP

MSM-IP (uses IP multicast)

MOBILITY IMPLEMENTATION: LOCATION MOBILITY

BENEFITING LAYER



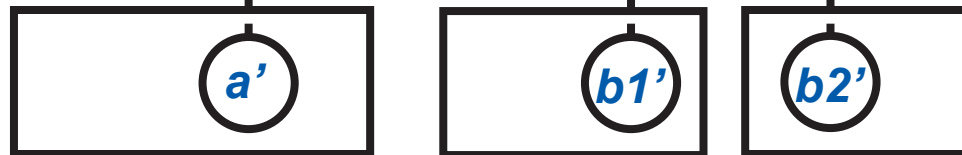
1 *B* or *b1* destroys their registration

2 *B* or *b2* creates a new registration, session state transferred from *b1* to *b2*

LAYER IMPLEMENTING MOBILITY



unless this is a case of process migration, *b1* and *b2* are on same machine



the hard work of implementing location mobility is updating locations and *a*'s session state so that *B* can be found at *b2*

disconnection may come from this layer, but there is no mobility involving this layer

examples:

TCP Migrate

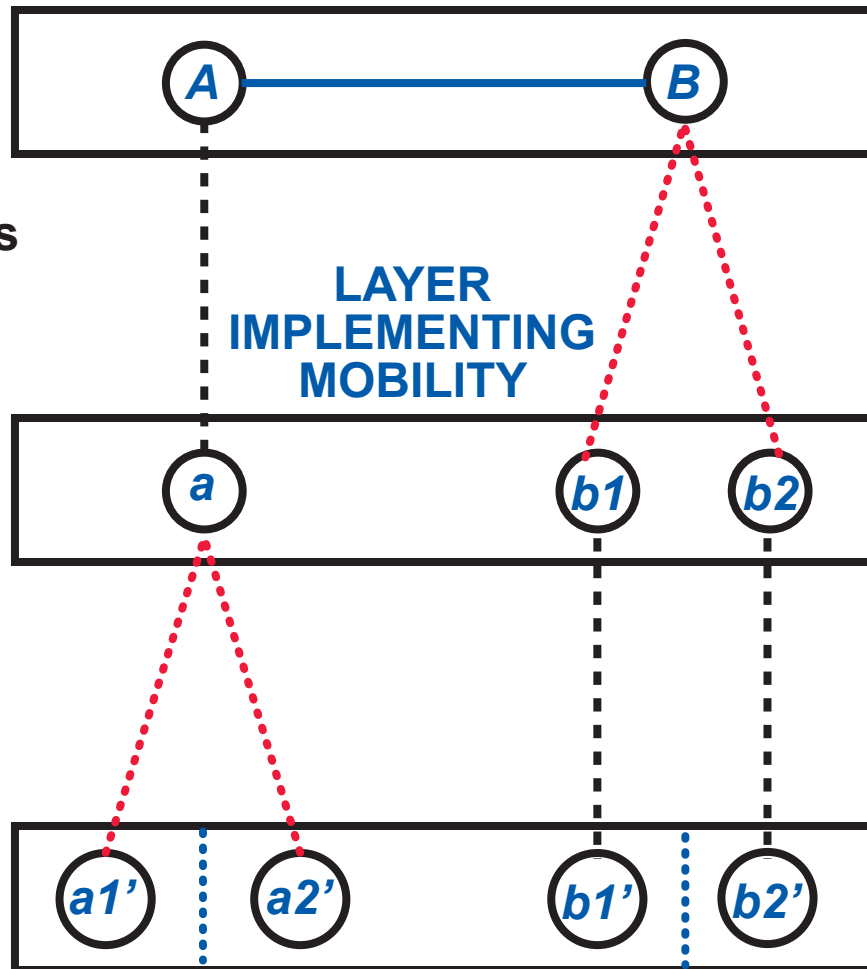
“Serval: An end-host stack for service-centric networking” [Nordstrom, Shue, Gopalan, Kiefer, Arye, Ko, Rexford & Freedman 2012]

ARE THEY REALLY DIFFERENT?

**ATTACHMENT
MOBILITY**

**LOCATION
MOBILITY**

BENEFITING LAYER



registration change
at implementing layer's
lower interface

registration change
at implementing layer's
upper interface

**LAYER
IMPLEMENTING
MOBILITY**

distributed layer
state affected:
attachments
links
forwarding

distributed layer
state affected:
locations
sessions

EXAMPLE: COMPOSITION OF MOBILITY MECHANISMS

WITHIN A LAYER:

- mobility at one end of a session is independent of mobility at the other end—either one can be attachment or location mobility, even simultaneously
- at one end of a session, location mobility can take over if attachment mobility is failing

established by verification of a formal model of the session protocol, reasoning about the layer state

“Compositional network mobility”

[Zave & Rexford 2012]

ACROSS LAYER BOUNDARIES:

- mobility mechanisms in adjacent layers are logically independent, can co-exist and even operate simultaneously

established by reasoning about actions at the layer interface

affected by location mobility

affected by attachment mobility

locations

sessions

members

routes

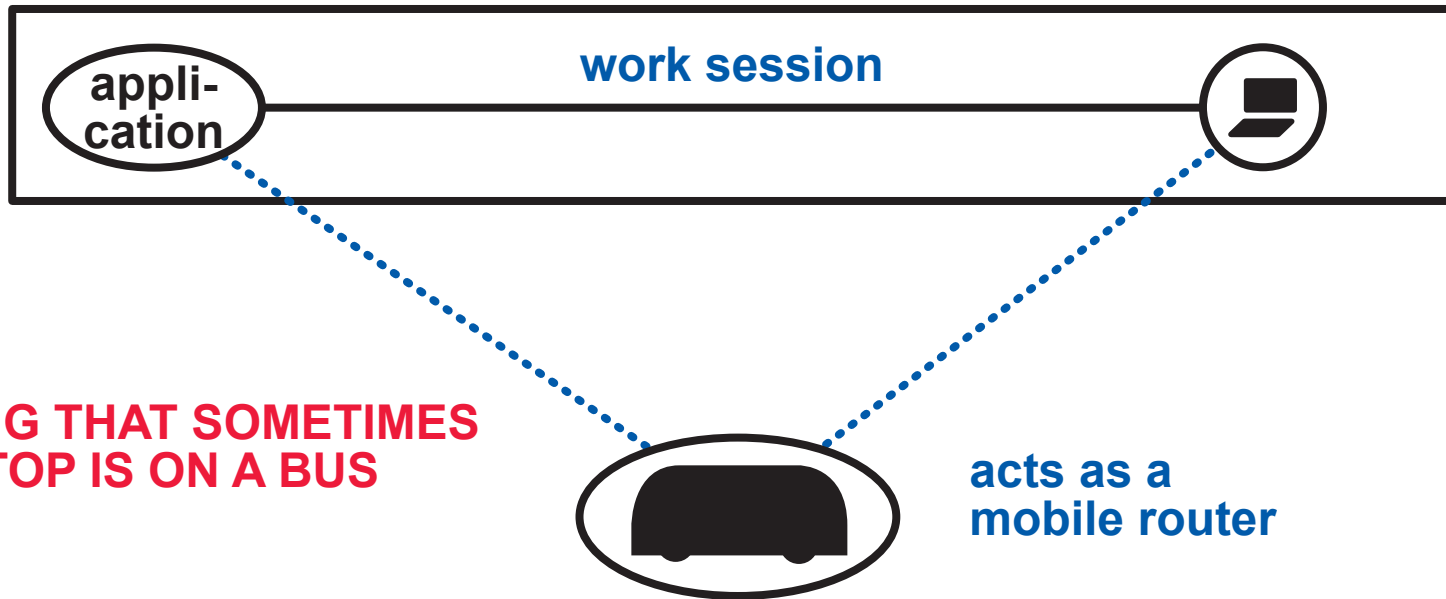
attachments

links

there is no dependency between the two state partitions

EXAMPLE: NEW DESIGNS FROM THE MOBILITY SPACE

THE GOAL IS TO PROVIDE MOBILITY FOR THIS LAPTOP . . .

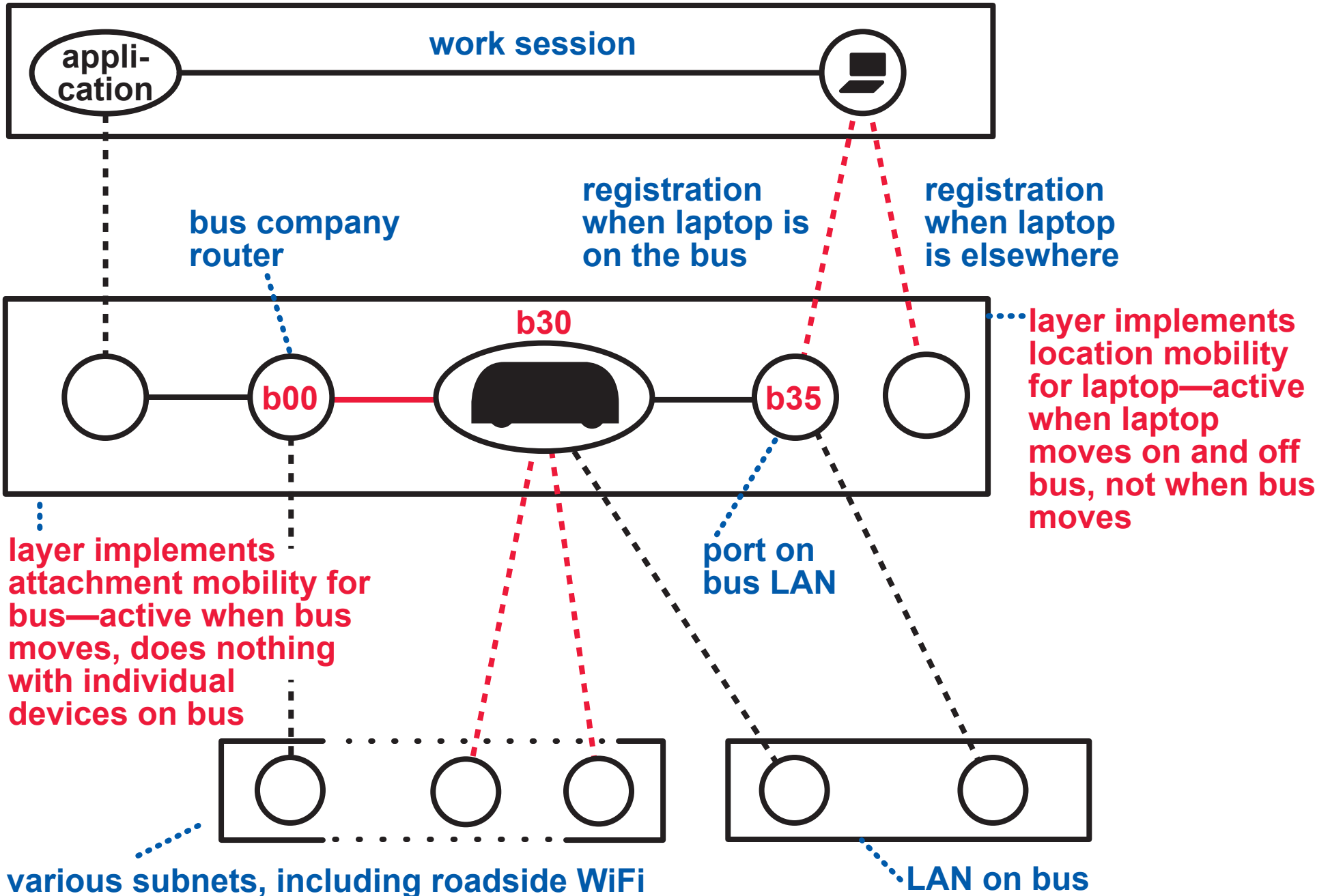


. . . NOTING THAT SOMETIMES THE LAPTOP IS ON A BUS

we want to avoid, e.g., . . .

- . . . solutions that require updates for every passenger when the bus moves
- . . . solutions that require an update for the bus when a passenger gets on or off

EXAMPLE: NEW DESIGNS FROM THE MOBILITY SPACE



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SUMMARY

SOME PURPOSES OF ABSTRACTIONS

to provide applications with richer communication services and cleaner interfaces to them

to manage complexity with separation of concerns and use of general-purpose theories

to compare and relate research results

to map out structured spaces of design trade-offs

to compose successful solutions to diverse problems

to implement layers by means of code generation and code re-use

SOME CHARACTERISTICS (OR, MORE ACCURATELY, GOALS) OF THE GEOMORPHIC VIEW

basic structures such as members, channels, and routing can look very different at different levels of the stack

the same basic pattern or template is instantiated many times, for many different purposes, at different levels and scopes

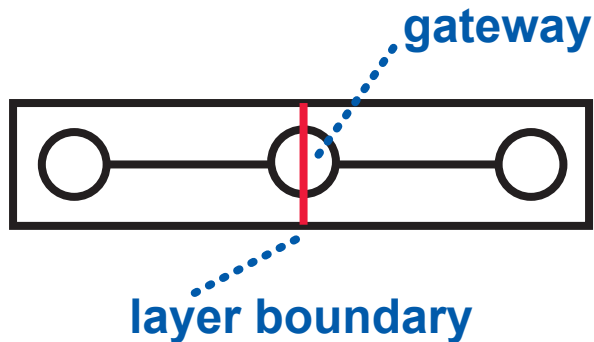
each design has exactly one correct description

by partitioning the control state of a layer, facilitates composition of mechanisms within a layer

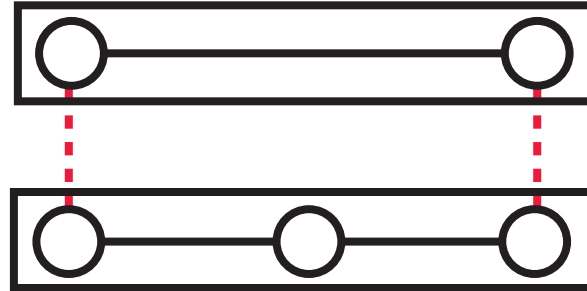
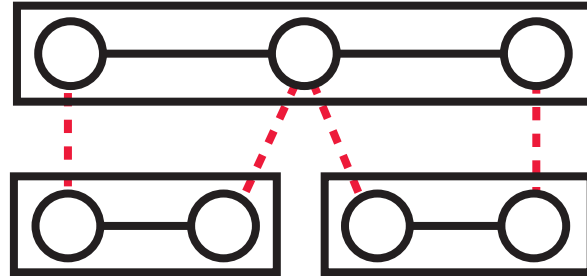
by solidifying layer interfaces, facilitates composition of layers in a network stack

A BIG TRADE-OFF

STRUCTURES YOU OFTEN SEE



THE GEOMORPHIC VIEW OF THESE STRUCTURES



WHAT ABOUT THE REDUNDANCY AND OVERHEAD?

How much?
Routers already have forwarding rules from multiple layers.

Redundancy and overhead can be removed by optimization—at the cost of less resilience to change.

- each mechanism is *ad hoc*, and its interactions with other mechanisms are unpredictable

- interactions among mechanisms can be studied, will become well-known

- each layer can achieve multiple purposes

THE GEOMORPHIC VIEW IS NEW AND IMMATURE

IN PARTICULAR, THE VIEW OF LAYER STATE (partitions, dependencies) IS VERY SIMPLISTIC . . .

. . . because the only issue we have studied in enough detail is mobility

OTHER ISSUES TO BE INVESTIGATED

- anycast, multicast, broadcast, etc.
- multihoming
- middleboxes
- enrollment
- authentication
- access control
- privacy
- failure recovery
- resource management

ON THE OTHER HAND, THESE TWO IDEAS NOW SEEM INTUITIVELY OBVIOUS:

- To understand the control plane, you must first understand the data plane.
- The geomorphic view of the data plane—consisting of “complete” layers that can be instantiated freely with different levels, scopes, and purposes—is a better abstraction of the data plane than the classic Internet architecture.

