

# ***Securing Wireless Localization: Living with Bad Guys***

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# Talk Overview

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- Wireless Localization Background
- Attacks on Wireless Localization
  - Time of Flight
  - Signal Strength
  - Angle of Arrival
  - Region Inclusion
  - Hop Count
  - Neighbor Location
- Coping with Localization Threats
  - Multimodal Localization Strategies
  - Robust Statistics
- Conclusions and Future Directions

# ***What is Localization?***

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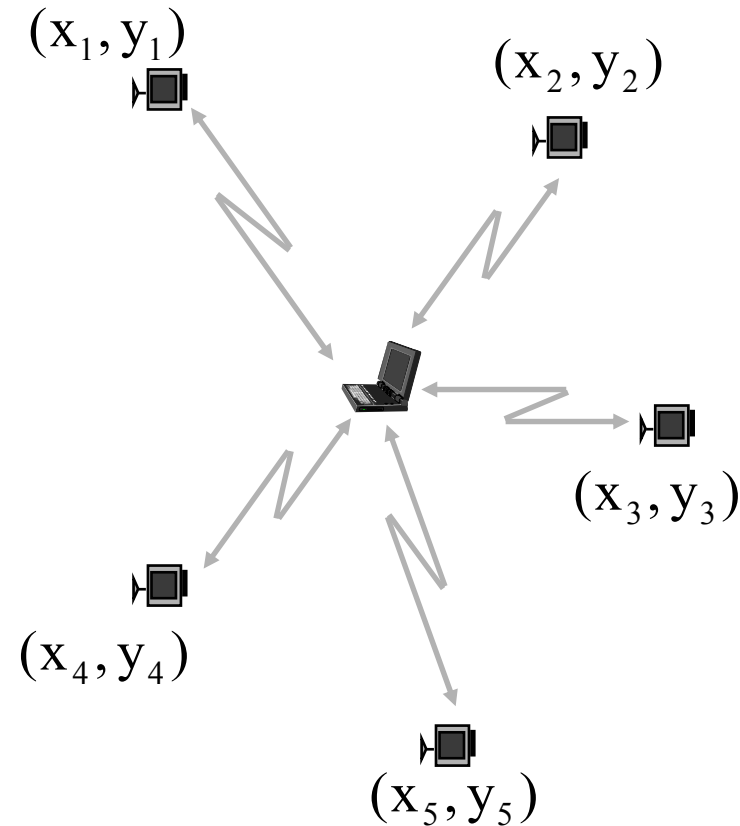
- Localization is important for facilitating location-based services
- **Goal:** Determine the location of one or more wireless devices based on some form of measurements
- Useful measurements:
  - Time of flight (TOA)
  - Time difference of flight (Tdoa)
  - Energy of flight (DoA based on Signal Strength)
  - Phase of flight (AoA = Angle of arrival from fixed stations)
  - Perspective of flight (Visual Cues)
  - Hop count to anchors: Correlated with distance
  - Neighbor Location: Find regions
- Examples...

# Use Neighbor Locations: Centroids

- Scenario:
  - A set of anchor nodes with known locations are deployed as infrastructure for localization
- Wireless devices localize by calculating the centroid of the anchor points they hear:

$$(\hat{x}, \hat{y}) = \left( \frac{x_1 + x_2 + \dots + x_n}{n}, \frac{y_1 + y_2 + \dots + y_n}{n} \right)$$

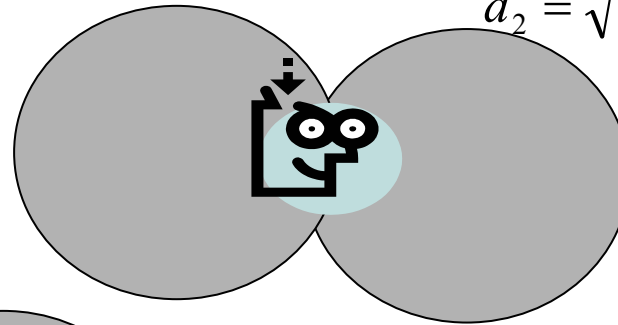
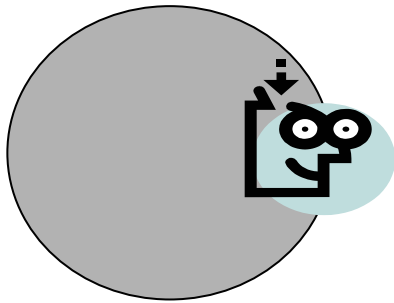
- Refine by averaging the values of the other nodes within the signal range



# Time of Flight (S=R) Localization

- Send a signal to receiver and back
- Measure RTT, know velocity of propagation
- Calculate Distance -

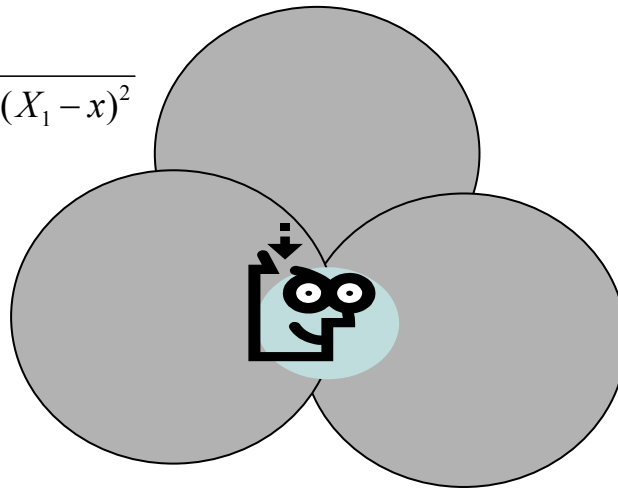
$$d_1 = \sqrt{(Y_1 - y)^2 + (X_1 - x)^2}$$
$$d_2 = \sqrt{(Y_2 - y)^2 + (X_2 - x)^2}$$



$$d_1 = c(rtt) = \sqrt{(Y_1 - y)^2 + (X_1 - x)^2}$$

Lateration

very  
common  
local  
triangulation  
solve  $[Ax=b]$



$$d_1 = \sqrt{(Y_1 - y)^2 + (X_1 - x)^2}$$

$$d_2 = \sqrt{(Y_2 - y)^2 + (X_2 - x)^2}$$

$$d_3 = \sqrt{(Y_3 - y)^2 + (X_3 - x)^2}$$

# Signal Strength

- Underlying Principle: Signal strength (RSSI) is a function of distance

- Free Space Propagation Model

$$P_r = P_t \left[ \frac{\sqrt{G_t} \lambda}{4\pi d} \right]^2$$

- Two-Path (Single Ground Reflection Model)

$$P_r = P_t \left[ \frac{\sqrt{G_t} h_t h_r}{d^2} \right]^2$$

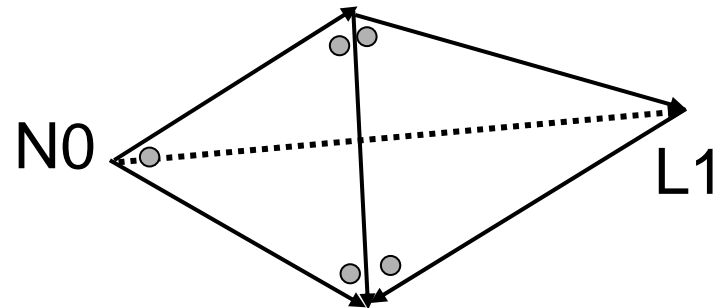
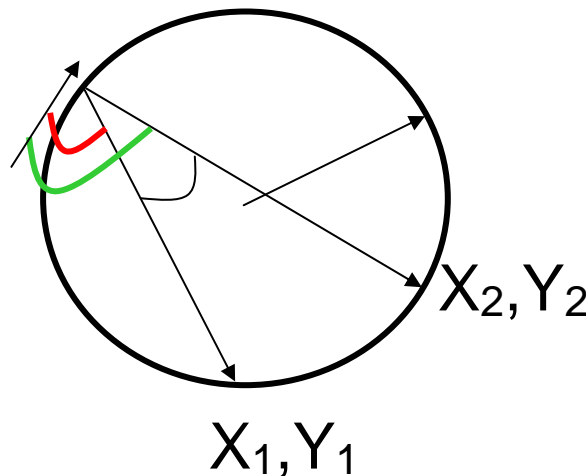
- Generalized Path Loss Model

$$P_r = P_t K \left[ \frac{d_0}{d} \right]^\gamma$$

- Use known landmark locations and RSSI-Distance relationship to setup a least squares problem

# Angle of Arrival Localization

- One can determine an orientation w.r.t a reference direction
- Angle of Arrival (AoA) from two different points and their distances
- You can locate a point on a circle. Similar AoA from another point gives you three points. Then triangulate to get a position



$$a/\sin A = b/\sin B = c/\sin C$$

# ***AoA capable nodes***

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## ■ Cricket Compass (MIT Mobicom 2000)

- Uses 5 ultra sound receivers
- 0.8 cm each
- A few centimeters across
- Uses tdoa (time difference of arrival)
- +/- 10% accuracy

## ■ Medusa sensor node (UCLA node)

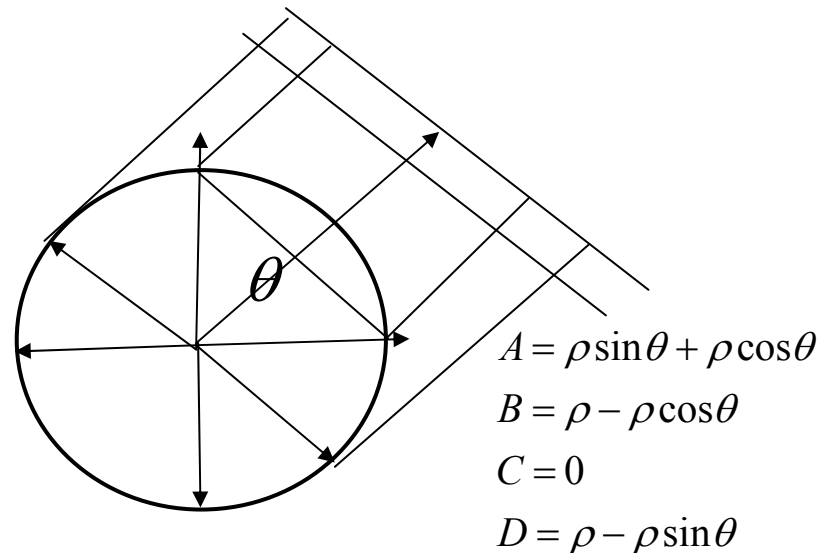
- Mani Srivatsava et.al

## ■ Antenna Arrays



# AoA Using Visual Cues

- Color cylinder
- Determine proportion of colors



Taking the ratios  $A/D$  and  $A/B$  and solving for theta

$$\sin \theta = (A + B - D) / (A + B + D)$$

$$\cos \theta = (A - B + D) / (A + B + D)$$

$$\theta = \arctan((A + B - D) / (A - B + D))$$

# Attacks on Localization

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- Most security and privacy issues for wireless networks are best addressed through cryptography and network security
- **End of Day Analysis:** Not all security issues can be addressed by cryptography!
- Non-cryptographic attacks on wireless localization:
  - Adversaries may affect the measurements used to conduct localization
  - Adversaries may physically pick up and move devices
  - Adversaries may alter the physical medium (adjust propagation speed, introduce smoke, etc.)
  - Many, many more **crazy** attacks...
- New Field: Securing Wireless Localization
  - “Secure Verification of Location Claims,” Sastry and Wagner
  - “Secure Positioning in Sensor Networks,” S. Capkun and J.P. Hubaux
  - “SeRLoc: Secure range-independent localization for wireless networks,” L. Lazos and R. Poovendran
  - “Securing Wireless Localization: Living with Bad Guys,” Z. Li, Y. Zhang, W. Trappe and B. Nath (expanded version under submission)

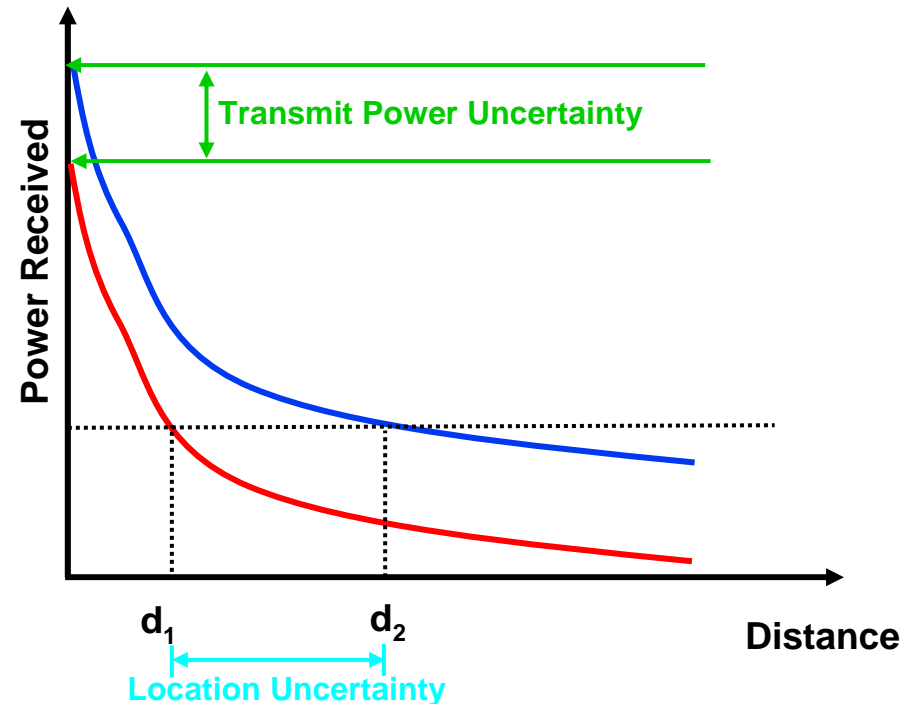
# Possible Attacks vs. Localization Algorithms

Property	Example Algorithms	Attack Threats
Time of Flight	Cricket	<ul style="list-style-type: none"><li>➤ Remove direct path and force radio transmission to employ a multipath;</li><li>➤ Delay transmission of a response message;</li><li>➤ Exploit difference in propagation speeds (speedup attack, transmission through a different medium).</li></ul>
Signal Strength	RADAR, SpotON, Nibble	<ul style="list-style-type: none"><li>➤ Remove direct path and force radio transmission to employ a multipath;</li><li>➤ Introduce different microwave or acoustic propagation loss model;</li><li>➤ Transmit at a different power than specified by protocol;</li><li>➤ Locally elevate ambient channel noise</li></ul>
Region Inclusion	APIT, SerLoc	<ul style="list-style-type: none"><li>➤ Enlarge neighborhood by wormholes;</li><li>➤ Manipulate the one-hop distance measurements;</li><li>➤ Alter neighborhood by jamming along certain directions</li></ul>

<b>Property</b>	<b>Example Algorithms</b>	<b>Attack Threats</b>
Angle of Arrival	APS	<ul style="list-style-type: none"> <li>➤ Remove direct path and force radio transmission to employ a multipath;</li> <li>➤ Change the signal arrival angel by using reflective objects, e.g., mirrors;</li> <li>➤ Alter clockwise/counter-clockwise orientation of receiver (up-down attack)</li> </ul>
Hop Count	DV-Hop	<ul style="list-style-type: none"> <li>➤ Shorten the routing path between two nodes through wormholes;</li> <li>➤ Lengthen the routing path between two nodes by jamming;</li> <li>➤ Alter the hop count by manipulating the radio range;</li> <li>➤ Vary per-hop distance by physically removing/displacing nodes</li> </ul>
Neighbor Location	Centroid, SerLoc	<ul style="list-style-type: none"> <li>➤ Shrink radio region (jamming); Enlarge radio region (transmit at higher power, wormhole);</li> <li>➤ Replay; Modify the message; Physically move locators;</li> <li>➤ Change antenna receive pattern</li> </ul>

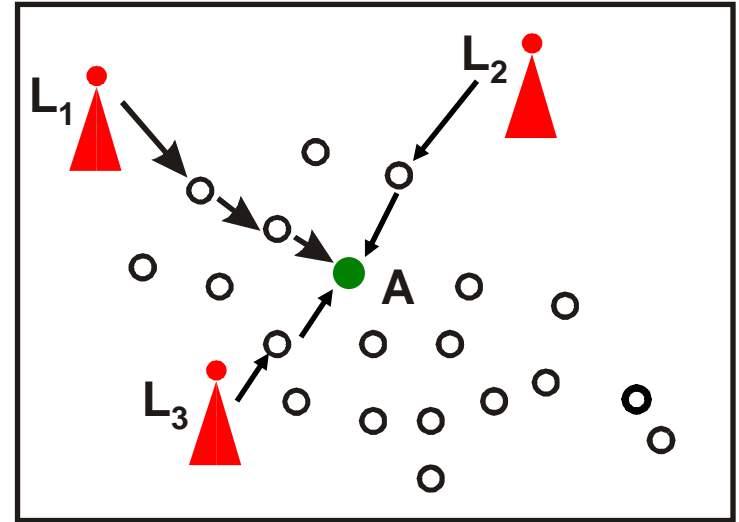
# Signal Strength Attack on Localization

- Signal strength wireless localization are susceptible to power-distance uncertainty relationships
- Adversary may:
  - Alter transmit power of nodes
  - Remove direct path by introducing obstacles
  - Introduce absorbing or attenuating material
  - Introduce ambient channel noise



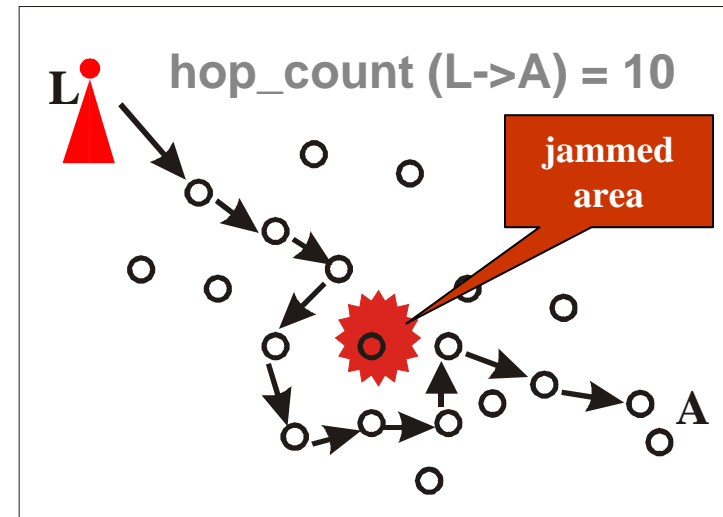
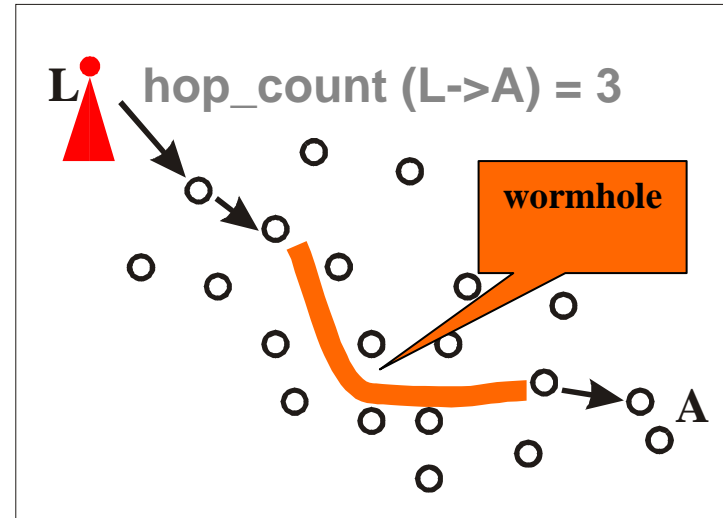
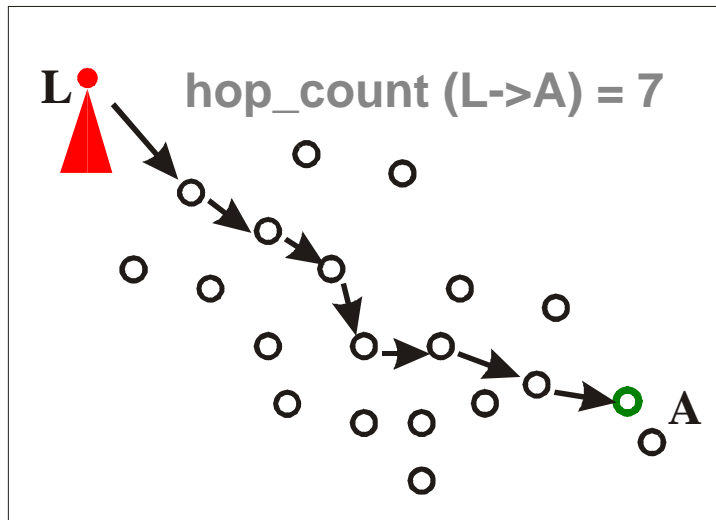
# Attacks on Hop-Count Methods

- DV-hop localization algorithm:
- Obtain the hop counts between a sensor node and several locators.
- Translate hop counts to actual distance.
- Localize using triangulation.



It is critical to obtain the correct hop counts between sensor nodes and every locator.

# Attacks on Hop-Count Methods, pg. 2



# Defenses for Wireless Localization

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## ■ Multimodal Localization:

- Most localization techniques employ a single property
- Adversary only has to attack one-dimension!!!
- **Defense Strategy:** Make the adversary have to attack several properties simultaneously
- Example: Do signal strength measurements correspond to TOF measurements?

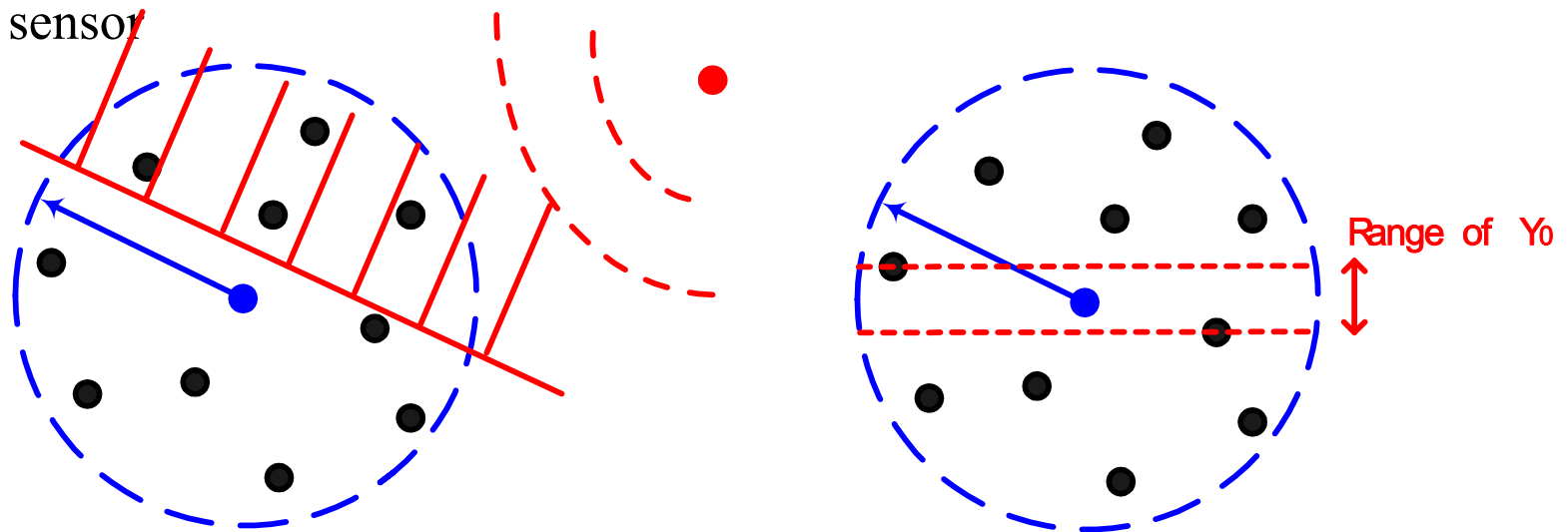
## ■ Robust Statistical Methods:

- **Defense Strategy:** Ignore the wrong values introduced by adversaries
- Develop robust statistical estimation algorithms and data cleansing methods
- Interesting behavior: Its best for the adversary not to be too aggressive!



# Multimodal Techniques

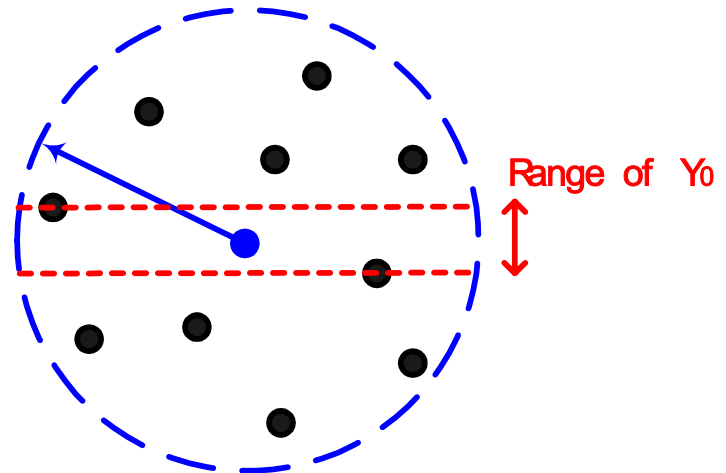
- Multimodal localization strategies: exploiting several properties simultaneously to corroborate each other and improve robustness
- Example: Centroid
  - **Attacks:** generally involve modifying neighboring list
  - **Defense:** use both neighbor location and a two-sector antenna on each sensor



$$\hat{x}_0 = \frac{1}{N} \sum_{i=1}^N x_i, \quad \hat{y}_0 = \frac{1}{N} \sum_{i=1}^N y_i$$

# Multimodal Technique

- Only the neighbors that are closest to the sensor in the x-coordinate or y-coordinate will affect the estimation
- Robust to wrong neighbor information
- Neighbor coordinates rule: the neighbors in the upper sector have larger Y coordinates than the neighbors in lower sector
  - Ensure correct orientation
  - Detect existence of attacks



# ***Robust: Localization with Anchor Nodes***

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- Anchor nodes have their positions  $\{(x, y)\}$  known
- Distances to anchor nodes  $d$  are estimated through DV-hop or signal strength or other distance estimation methods
- $\{(x, y, d)\}$  values map out a parabolic surface  $d(x, y)$  whose minimum value  $(x_0, y_0)$  is the wireless device location
- Least squares (LS) algorithm can be used to find  $(\hat{x}_0, \hat{y}_0)$

$$(\hat{x}_0, \hat{y}_0) = \arg \min_{(x_0, y_0)} \sum_{i=1}^N (\sqrt{(x_i - x_0)^2 + (y_i - y_0)^2} - d_i)^2$$

# *What if Attacks Exist?*

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- Adversary can alter the distance measurement through wormholes or jamming attacks
- One significant deviation of distance measurement may drive the location estimation far from the true value
- The fundamental reason for this vulnerability to attacks is that

**Least squares algorithm is not robust to outliers!**

- The misinformation produced by the adversary are outliers in the location estimation problem
- Redundancy within network can be exploited to combat attacks

# Robust Statistics

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- Least median squares (LMS) algorithm

$$(\hat{x}_0, \hat{y}_0) = \arg \min_{(x_0, y_0)} \text{med}(\sqrt{(x_i - x_0)^2 + (y_i - y_0)^2} - d_i)^2$$

- Proposed by Rousseeuw
- With a robust cost function, a small fraction of outliers won't affect the cost function significantly
- In the absence of noise, LMS algorithm can tolerate up to 50 percent outliers
- Exact calculation of LMS solution is computational expensive

# Least Median Squares Algorithm

- Solve random subsets of  $\{(x_i, y_i, d_i)\}$  values to get several candidate  $(\hat{x}_0, \hat{y}_0)$
- Choose the candidate with the least median residue squares
- Identify the inliers and outliers according to the least median squares subset estimate

$$s_0 = 1.4826 \left(1 + \frac{5}{N-p}\right) \sqrt{\text{med } r_i^2} \quad w_i = \begin{cases} 1, & |r_i / s_0| > \gamma \\ 0, & \text{otherwise} \end{cases}$$

- Do a reweighted least squares algorithm to get the final estimate  $(\hat{x}_0, \hat{y}_0)$

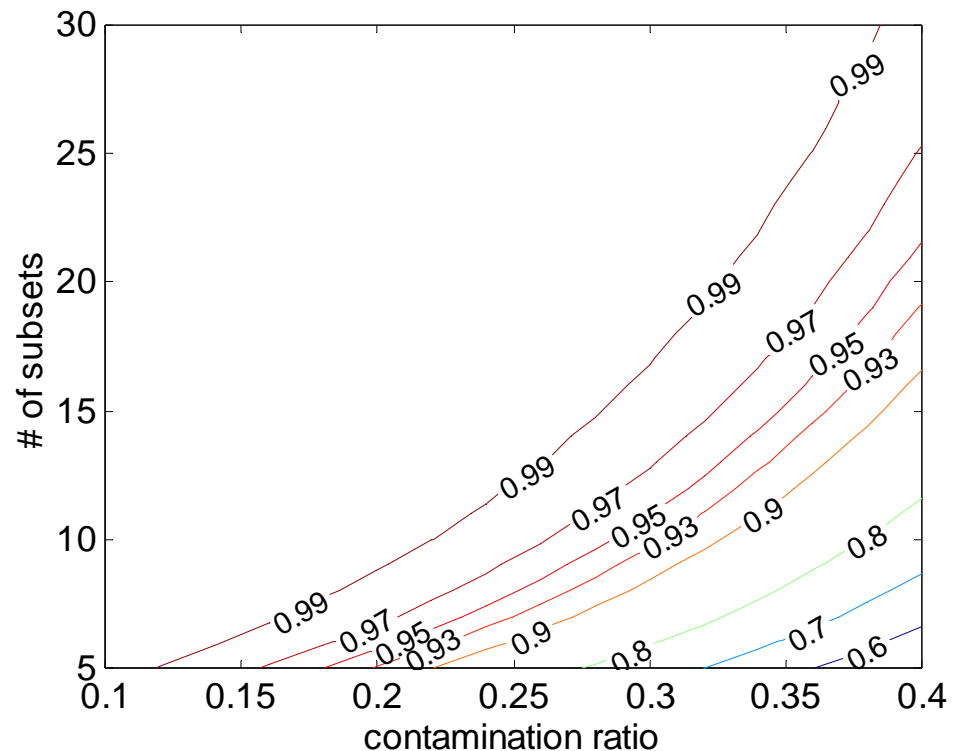
# Robust Localization with LMS

- How to choose  $M$ , the number of subsets and  $n$ , the size of a subset?
  - Hopefully, at least one subset among all subsets does not contain any contaminated sample

$$P = 1 - (1 - (1 - \varepsilon)^n)^M$$

- In our simulation:

- ✓  $n = 4$
- ✓  $M = 20$



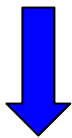
# Robust Localization with LMS (ctd.)

- How to estimate the location from the samples with reduced computation?
  - Linearization: suboptimal, but less complexity

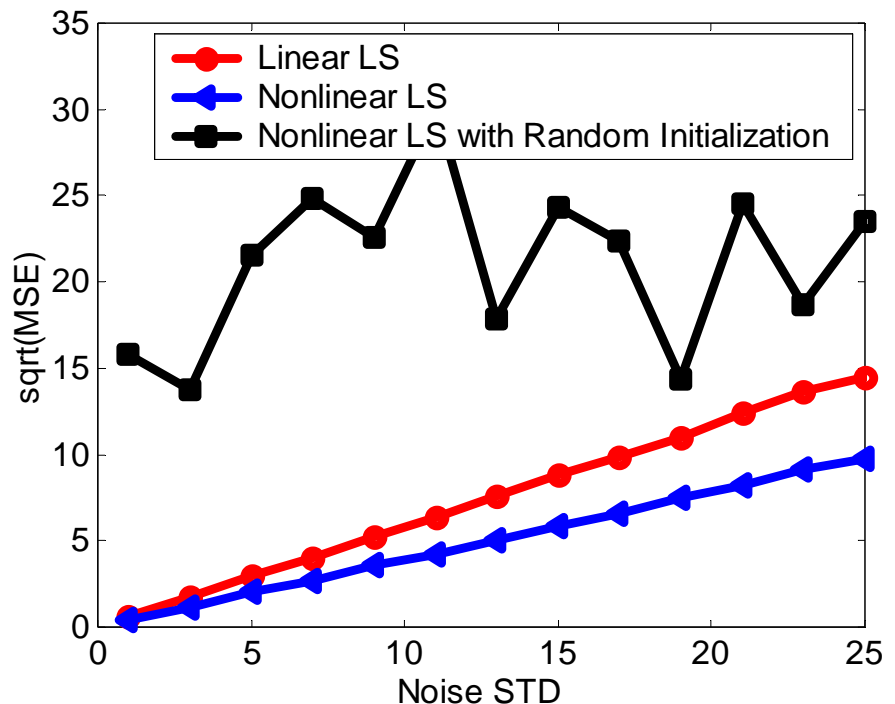
$$(x_1 - x_0)^2 + (y_1 - y_0)^2 = d_1^2$$

⋮

$$(x_N - x_0)^2 + (y_N - y_0)^2 = d_N^2$$



$$\frac{1}{N} \sum_{i=1}^N [(x_i - x_0)^2 + (y_i - y_0)^2] = \frac{1}{N} \sum_{i=1}^N d_i^2$$



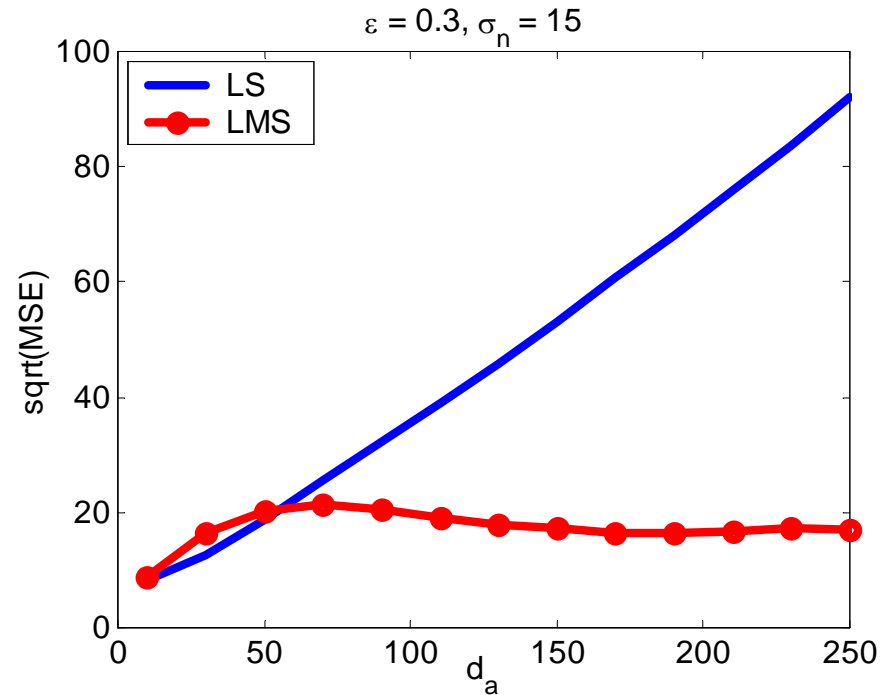
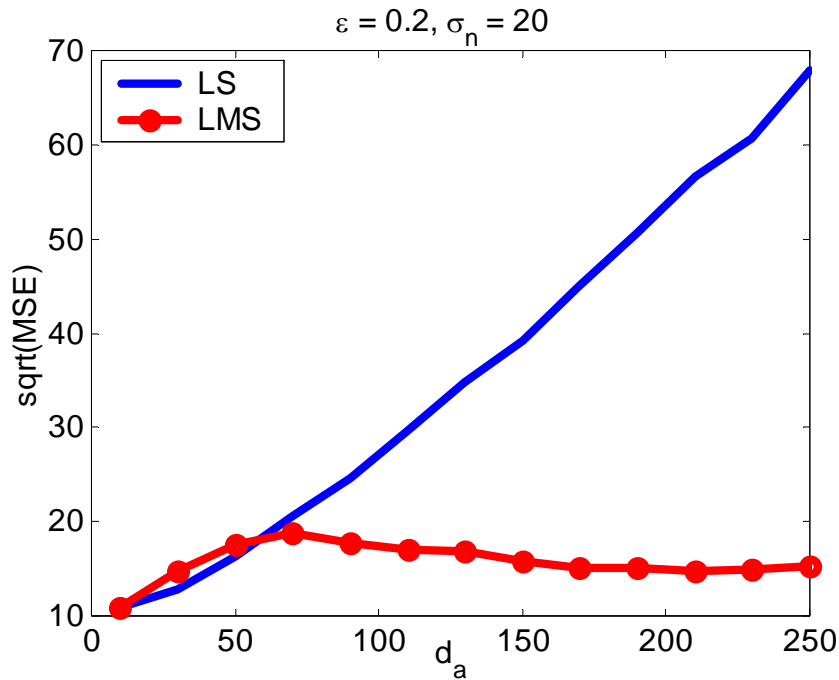


# ***Attack Model***

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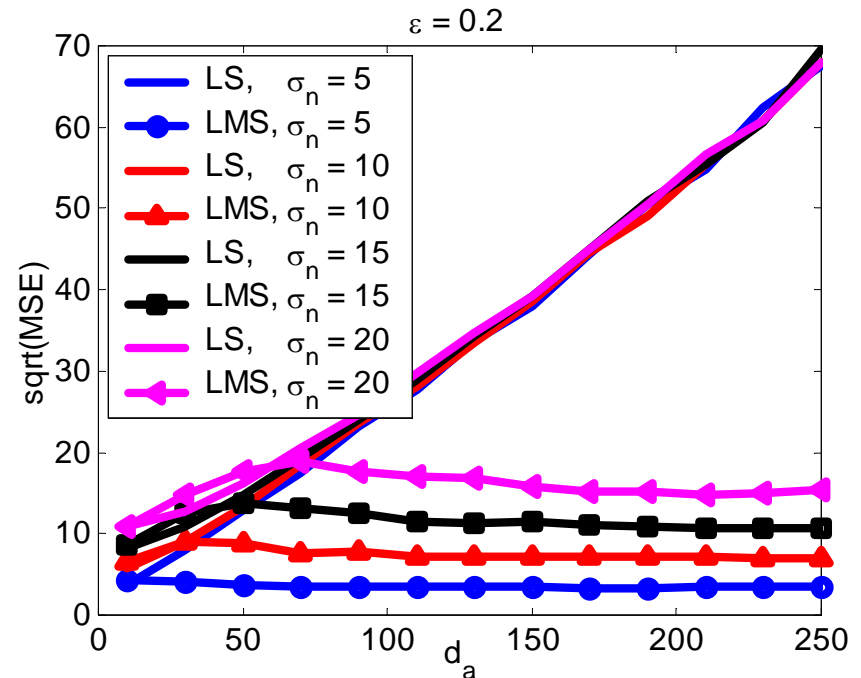
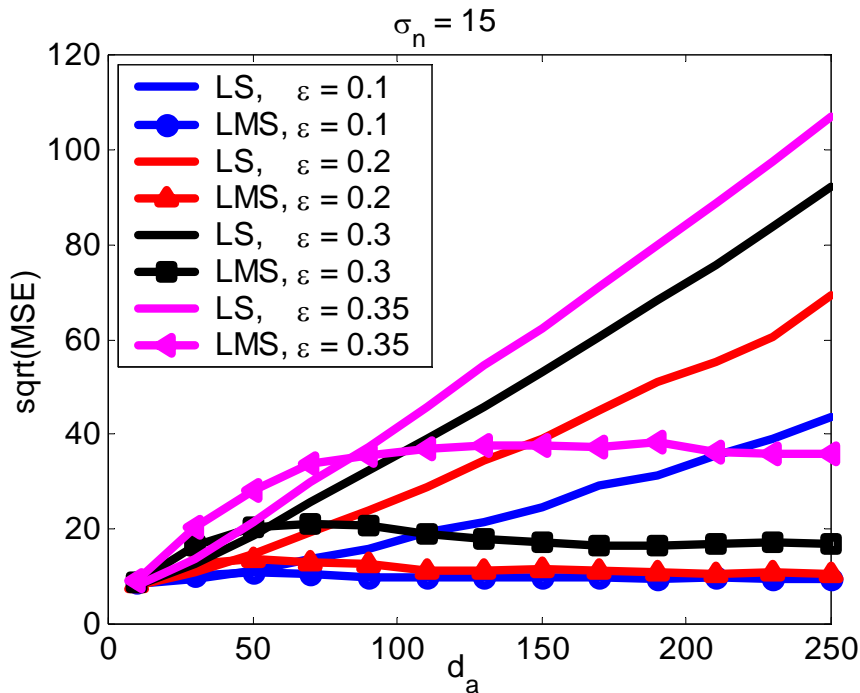
- The adversary successfully gains the power to arbitrarily modify the distance measurements to a fraction  $\varepsilon$  of the total anchor nodes
- The contamination ratio  $\varepsilon \leq 0.5$
- The adversary coordinates the tampering of measurements so that they will push the estimate toward the same wrong location  $(x_a, y_a)$
- $d_a$ , distance between  $(x_a, y_a)$  and  $(x_0, y_0)$ , is used to indicate the strength of the attack

# Performance of the LMS Algorithm



- MSE of LS algorithm increases as  $d_a$  increases
- MSE of LMS algorithm does not increase unboundedly with  $d_a$

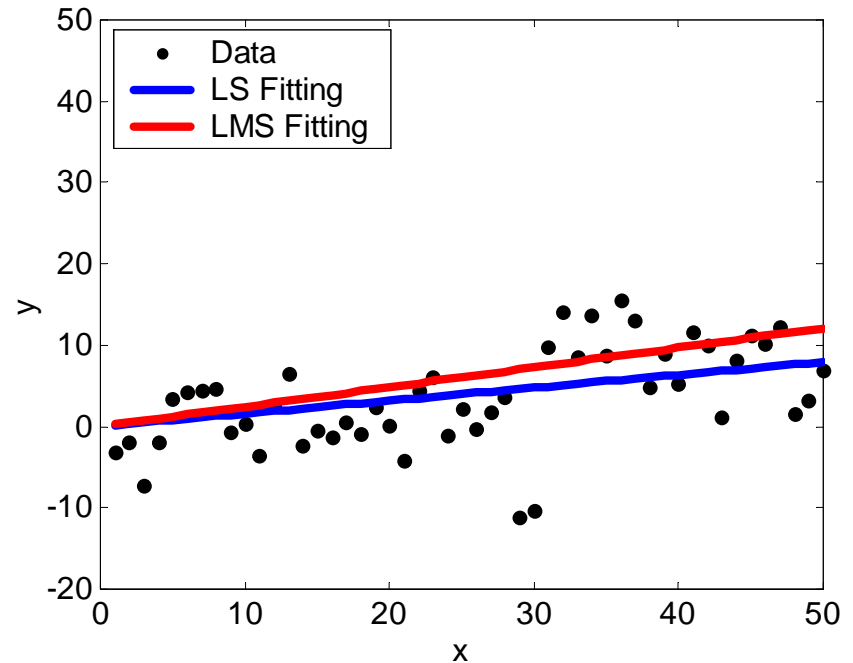
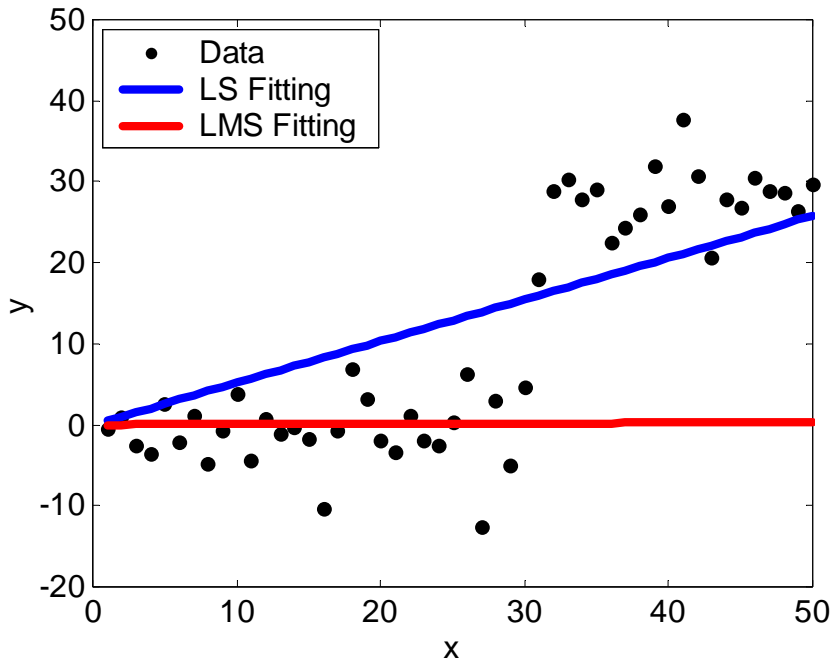
# Performance of the LMS Algorithm (ctd.)



- The larger contamination ratio, the worse the performance
- The larger the measurement noise level, the worse the performance

# When to Use LMS?

- At small  $d_a$ , LS performs better than LMS at a lower computational cost



(Conceptual Figures)

## ***When to Use LMS? (ctd.)***

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- Observation: the variance of the data with outliers is larger than that of the data without outliers
- Variance expansion indicates the attacking strength
- Estimate the variance in data using LS

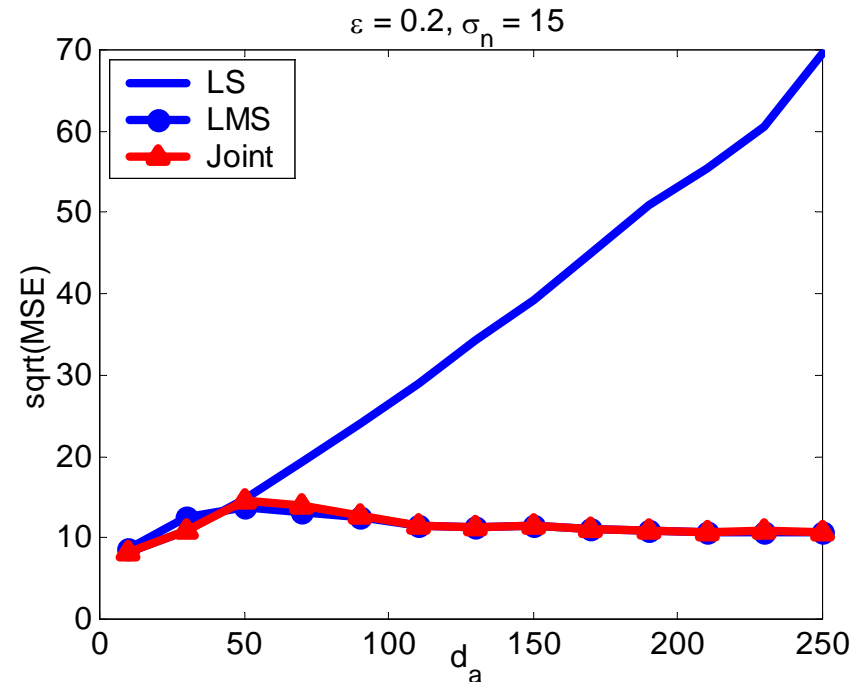
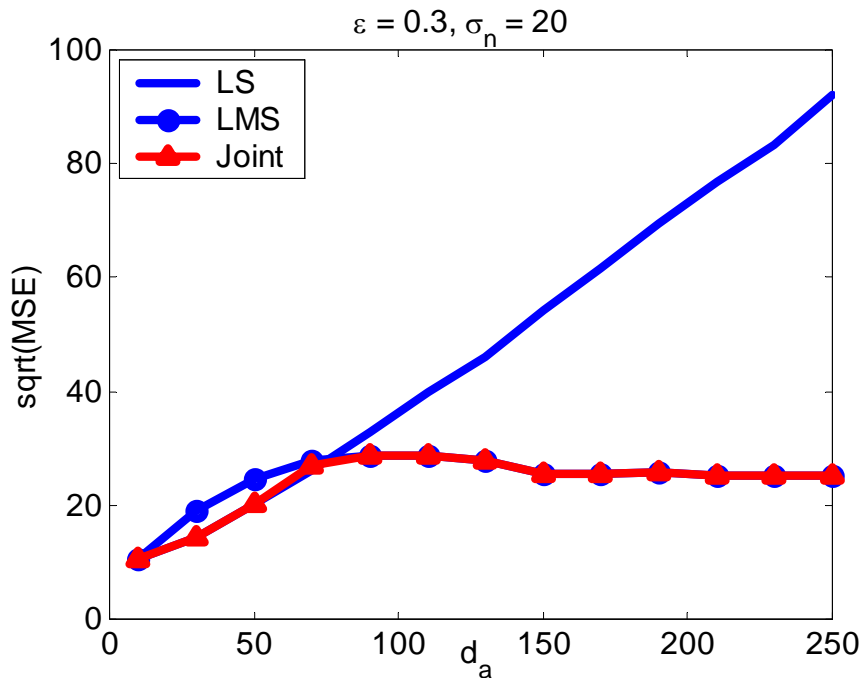
$$\hat{\sigma}_n = \sqrt{\frac{\sum r_i^2}{N - 2}}$$

- Assume the actual measurement noise level  $\sigma_n$  is known
- Use LMS only when

$$\frac{\hat{\sigma}_n}{\sigma_n} > T$$

# Performance of Joint LS and LMS Algorithm

- Empirically,  $T = 1.5$  is a good choice across all  $(\varepsilon, \sigma_n)$  pairs



**This improvement is achieved and we save computational complexity!!!**

# ***Conclusion and Remarks***

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- Wireless localization algorithms are important to future location-based services
- Several (non-cryptographic) attacks unique to wireless localization were identified
- We presented two strategies to cope with the effects of attacks on localization
  - Multimodal Localization
  - Robust Statistical Localization