

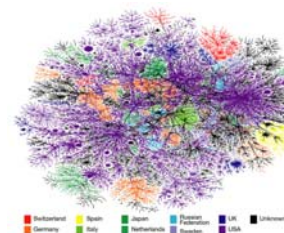
## Large Complex Networks: Deterministic Models (Recursive Clique-Trees)



<http://www.caida.org/tools/visualization/plankton/>

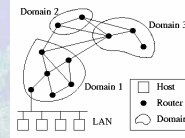
**Francesc Comellas**

Departament de Matemàtica Aplicada IV,  
Universitat Politècnica de Catalunya, Barcelona  
comellas@ma4.upc.edu



WWW

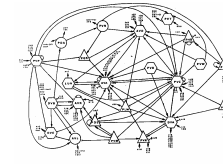
Internet



Air routes



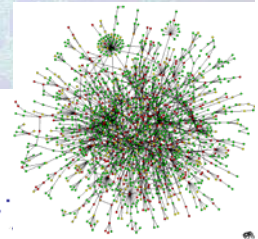
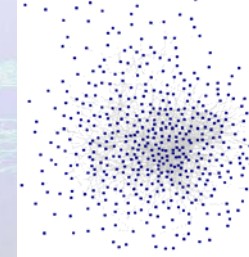
Power grid



C. Elegans



Erdős number



Proteins

### Complex systems

Different elements (nodes)  
Interaction among elements (links)

### Complex networks

Mathematical model: **Graphs**

### Real networks very often are

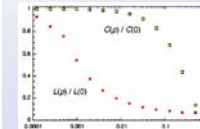
**Large**  
**Small-world**  
small diameter  $\log(|V|)$ , large clustering  
**Scale-free**  
power law degree distribution ("hubs")  
**Self-similar / fractal**

### Deterministic models

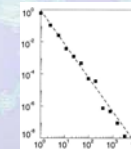
Based on cliques  
(hierarchical graphs, recursive  
clique-trees, Apollonian graphs)

### Most "real" networks are

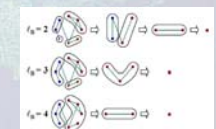
**small-world**    **scale-free**    **self-similar**



Small diameter (logarithmic)  
Milgram 1967  
High clustering  
Watts &  
Strogatz 1998



Power law  
(degrees)  
Barabási &  
Albert 1999



Fractal  
Song, Havlin &  
Makse 2005, 2006

## Main parameters (invariants)

Diameter - average distance

### Degree

$\Delta$  degree.  
 $P(k)$ : Degree distribution.

### Clustering

Are neighbours of a vertex also neighbours among them?

## Small-world networks

small diameter (or average dist.)  
high clustering

6 degrees of separation !    Stanley Milgram (1967)  
160 letters    Omaha -Nebraska- -> Boston



Small world phenomenon in social networks  
What a **small-world** !

### Structured graph

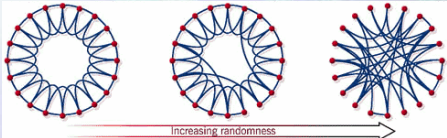
- high clustering
- large diameter
- regular

### Small-world graph

- high clustering
- small diameter
- almost regular

### Random graph

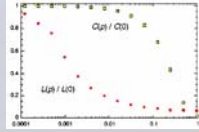
- small clustering
- small diameter



$|V|=1000$   $\Delta=10$   
 $D=100$   $d=49.51$   
 $C=0.67$

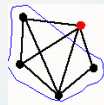
$|V|=1000$   $\Delta=8-13$   
 $D=14$   $d=11.1$   
 $C=0.63$

$|V|=1000$   $\Delta=5-18$   
 $D=5$   $d=4.46$   
 $C=0.01$



Watts & Strogatz,  
 Collective dynamics of "small-world" networks,  
 Nature 393, 440-442 (1998)

### Network characteristics



**Clustering**  $C(v) = \frac{\# \text{ of links among neighbors}}{n(n-1)/2}$

**Diameter or Average distance**

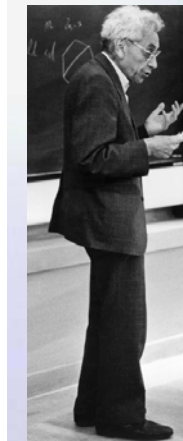
Maximum communication delay

**Degree distribution**

Resilience

Real life networks are clustered, large  $C(p)$ , but have small average distance  $L(p)$ . Very often they are also scale-free

	L	L <sub>rand</sub>	C	C <sub>rand</sub>	N
WWW	3.1	3.35	0.11	0.00023	153127
Actors	3.65	2.99	0.79	0.00027	225226
Power Grid	18.7	12.4	0.080	0.005	4914
C. Elegans	2.65	2.25	0.28	0.05	282



### Erdős number

<http://www.oakland.edu/enp/>

1- 509  
 2- 7494

N= 268,000 Jul 2004  
 (connected component)

D=23 R=12  $D_{avg} = 7.64$   
 $\delta=1$   $\Delta=509$   $\Delta_{avg} = 5.37$   
 $C = 0.14$

- Erdős number 0 --- 1 person
- Erdős number 1 --- 504 people
- Erdős number 2 --- 6593 people
- Erdős number 3 --- 33605 people
- Erdős number 4 --- 83642 people
- Erdős number 5 --- 87760 people
- Erdős number 6 --- 40014 people
- Erdős number 7 --- 11591 people
- Erdős number 8 --- 3146 people
- Erdős number 9 --- 819 people
- Erdős number 10 --- 244 people
- Erdős number 11 --- 68 people
- Erdős number 12 --- 23 people
- Erdős number 13 --- 5 people

(MathSciNet Jul 2004)

### Notable Erdős coauthors:

- Frank Harary (257 coauthors)
- Noga Alon (143 coauthors)
- Saharon Shelah (136)
- Ronald Graham (120)
- Charles Colbourn (119)
- Daniel Kleitman (115)
- A. Odlyzko (104)

Erdős did not write a joint paper with his PhD advisor, Leopold Fejér

### Some other Erdős coauthors

articles with Erdős

- András Sárközy 57
- András Hajnal 54
- Ralph Faudree 45
- Richard Schelp 38
- Vera Sós 34
- Alfréd Rényi 32
- Cecil C. Rousseau 32
- Pál Turán 30
- Endre Szemerédi 29
- Ronald Graham 27
- Stephan A. Burr 27
- Joel Spencer 23
- Carl Pomerance 21
- Miklos Simonovits 21
- Ernst Straus 20
- Melvyn Nathanson 19

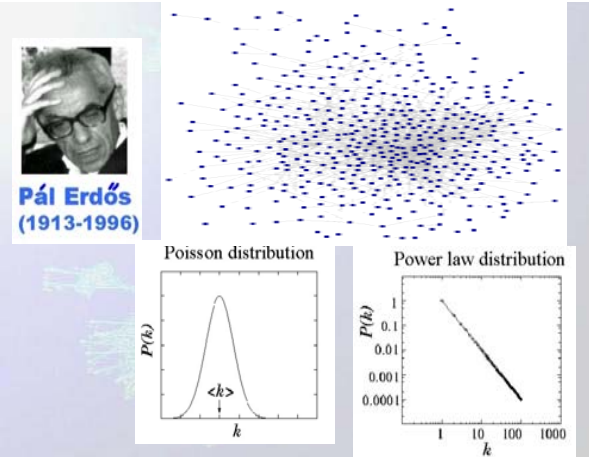
- Richard Rado 18
- Jean Louis Nicolas 17
- Janos Pach 16
- Béla Bollobás 15
- Eric Milner 15
- John L. Selfridge 13
- Harold Davenport 7
- Nicolaas G. de Bruijn 6
- Ivan Niven 7
- Mark Kac 5
- Noga Alon 4
- Saharon Shelah 3
- Arthur H. Stone 3
- Gabor Szegő 2
- Alfred Tarski 2
- Frank Harary 2
- Irving Kaplansky 2
- Lee A. Rubel 2

### Fields medals

Alain Connes	1982	France	3
William Thurston	1982	USA	3
Shing-Tung Yau	1982	China	2
Simon Donaldson	1986	Great Britain	4
Gerd Faltings	1986	Germany	4
Michael Freedman	1986	USA	3
Lars Ahlfors	1936	Finland	4
Jesse Douglas	1936	USA	4
Laurent Schwartz	1950	France	4
Atle Selberg	1950	Norway	2
Kunihiko Kodaira	1954	Japan	2
Jean-Pierre Serre	1954	France	3
Klaus Roth	1958	Germany	2
Rene Thom	1958	France	4
Lars Norstrand	1962	Sweden	2
John Milnor	1962	USA	3
Michael Atiyah	1966	Great Britain	4
Paul Cohen	1966	USA	5
Alexander Grothendieck	1966	Germany	5
Stephen Smale	1966	USA	4
Alan Baker	1970	Great Britain	2
Heisuke Hironaka	1970	Japan	4
Serge Novikov	1970	USSR	3
John G. Thompson	1970	USA	3
Enrico Bombieri	1974	Italy	2
David Mumford	1974	Great Britain	2
Pierre Deligne	1978	Belgium	3
Charles Fefferman	1978	USA	2
Gregori Margulis	1978	USSR	3
Daniel Quillen	1978	USA	4
Richard Borcherds	1998	S Afr/Gc Brtn	2
William T. Gowers	1998	Great Britain	4
Maxim L. Kontsevich	1998	Russia	4
Curtis McMullen	1998	USA	3
Vladimir Voevodsky	2002	Russia	4
Laurent Lafforgue	2002	France	inf
Andrei Okounkov	2006	USA	3
Terence Tao	2006	USA	3
Wendelin Werner	2006	France	3

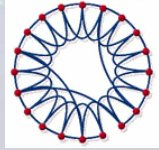
Winning bid: US \$1,831.80  
 Ended: Apr-30-04 09:58:51 PDT  
 Start time: Apr-20-04 09:58:51 PDT  
 History: 33 bids (US \$0.09 starting bid)  
 Winning bidder: madd\_greg (1)  
 Item location: Ann Arbor, MI United States (Detroit)  
 Ships to: Worldwide

Seller information: jaggedy (1,009 ★) m20  
 Feedback Score: 1609  
 Positive Feedback: 99.8%  
 Member since Jul-31-99 in United States



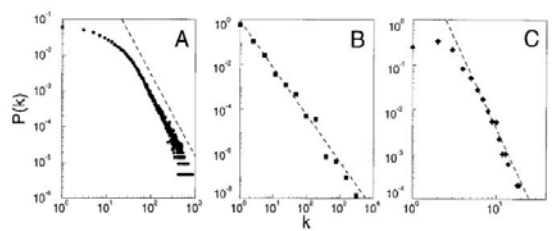
## Scale-free networks

## Scalability vs Fractality

1	1227	7	0	<b>SWCirculant</b>
2	1656	8	5	
3	1060	9	93	V =1000 Δ=8-13
4	401	10	806	
5	252	11	90	D = 14 d = 11.1
6	137	12	5	<b>Small World</b>
7	84	13	1	
8	46	14	0	C = 0.63
9	27			
10	26			
11	11			
12	5			
13	5			
14	3			
15	0			
16	0			
17	0			
18	1			
19	1			

**Power grid**  
|V|=4491 δ=1 Δ=19  
D = 46 d = 34.54  
Small World  
C = 0.08

A-L. Barabási i R. Albert,  
Emergence of scaling in random networks.  
Science 286, 509-510 (1999)



$$P(k) = k^{-\gamma}$$

A: actors	N=212.250	k=28.78	$\gamma=2.3$
B: WWW	N=325.729	k=5.46	$\gamma=2.67$
C: power grid	N= 4.94	k=2.67	$\gamma=4$

## Real networks for which we know the topology:

$$P(k) \sim k^{-\gamma}$$

<b>NON BIOLOGICAL</b>	$\gamma > 2$
www (in)	$\gamma = 2.1$
actors	$\gamma = 2.3$
citations	$\gamma = 3$
power grid	$\gamma = 4$
<b>BIOLOGICAL</b>	$\gamma < 2$
yeast protein-protein net	$\gamma = 1.5, 1.6, 1.7, 2.5$
E. Coli metabolic net	$\gamma = 1.7, 2.2$
yeast gene expression net	$\gamma = 1.4-1.7$
gene functional interaction	$\gamma = 1.6$

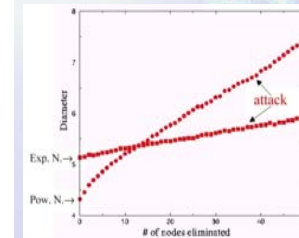
But, usual random models give:  $P(k) \sim e^{-k}$

## Interest on scale-free nets:

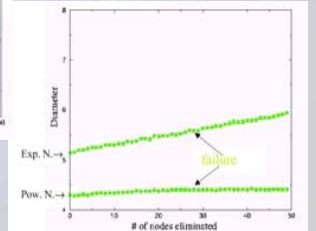
### Resilience / Survival of the WWW

Albert, Jeong, Barabási  
Nature 406, 378 (2000)

What happens when nodes fail randomly?



And when there are intentionate attacks to the best connected nodes?

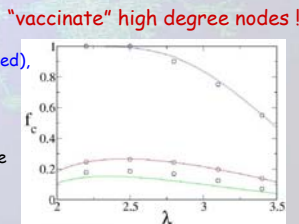


## Epidemics spreading / "vaccination" WWW, social networks

R. Cohen, D. ben-Avraham, S. Havlin;  
Efficient immunization of populations and computers  
Phys. Rev. Lett. 91, 247901 (2003)

$f_c$  threshold  
 $\lambda$  power law exponent  
upper- totally random  
lower- acquaintance immunisation (red),  
double acq. imm. (green)

method:  
\* select a node at random  
\* ask it to select a high degree node and immunize it



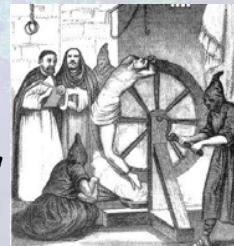
## Search in power-law networks

Adamic, Lukose, Puniyani, Huberman;  
Phys. Rev. E 64, 046135 (2001)

## Even the Inquisition knew about scale-free networks!!

From random "vaccination"

Arnau d'Amaurí 1209. Besièrs  
*Caedite eos.*  
*Robit enim Dominus qui sunt eius*  
Kill them all, God will know his own



to selection (see figure)

P. Ormerod, A.P. Roach;  
The Medieval inquisition: scale-free networks and the suppression of heresy.  
Physica A 339 (2004) 645-652

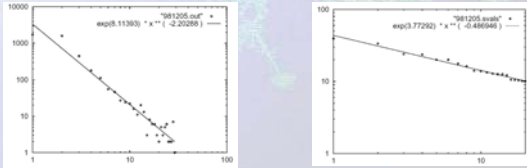
## spectral properties

- Connectivity and vulnerability (diameter, cut sets, distances between subsets)
- Scalability, expansion (Cheeger constants)
- Routings (spanning trees)
- Load balancing
- Clustering (triangles)
- Reconstruction (Ipsen & Mikhaliiov, 2001)
- Dynamical aspects (interlacing theorem)

# Experimental and simulation results

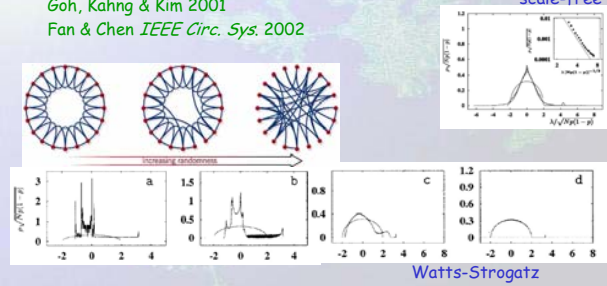
## WWW / Internet eigenvalues

Faloutsos, Faloutsos & Faloutsos, 1999



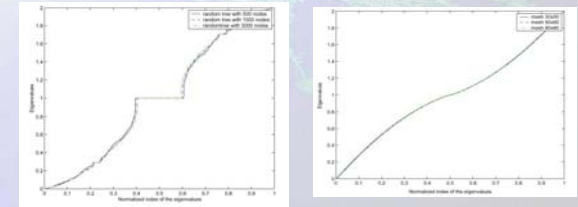
## Adjacency matrix eigenvalues (Watts & Strogatz model, Scale-free models)

Farkas, Derenyi, Barabási & Vicsek *Phys. Rev E* 2001  
 Goh, Kahng & Kim 2001  
 Fan & Chen *IEEE Circ. Sys.* 2002



## Normalized Laplacian eigenvalues (meshes, random trees)

Vukadinovic, Huang, Erlebach 2002



How to model real networks ?

Erdős-Rényi,  
 Watts-Strogatz  
 Barabási-Albert

other models ?

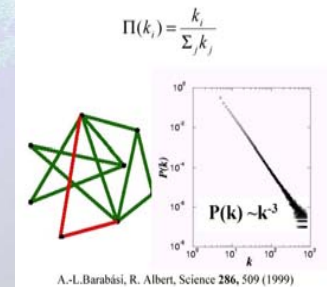
## Why appears a power law?

1. Networks grow continuously by addition of new nodes
2. Growth is NOT uniform: A new node will join, with high probability, an old well connected node

WWW: New documents point to "classic" references  
 Erdős: I would prefer to publish with a well known mathematician

"Standard" model: Barabási, Albert; *Science* 286, 509 (1999)

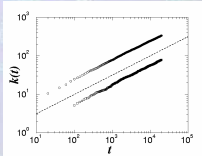
Preferential attachment : At each time unit a new node is added with m links which connect to existing nodes . The probability P to connect to a node i is proportional to its degree k<sub>i</sub>



## Mean Field Theory

$$\frac{\partial k_i}{\partial t} \propto \Pi(k_i) = A \sum_j k_j = \frac{k_i}{2t}, \text{ with initial condition } k_i(t_0) = m$$

$$k_i(t) = m \sqrt{\frac{t}{t_0}}$$



$$P(k_i(t) < k) = P(t_i > \frac{m^2 t}{k^2}) = 1 - P(t_i \leq \frac{m^2 t}{k^2}) = 1 - \frac{m^2 t}{k^2(m_0 + t)}$$

$$\therefore P(k) = \frac{\partial P(k_i(t) < k)}{\partial k} = \frac{2m^2 t}{m_0 + t} \frac{1}{k^3} \sim k^{-3} \quad \gamma = 3$$

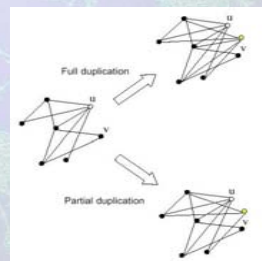
A.-L.Barabási, R. Albert and H. Jeong, *Physica A* 272, 173 (1999)

## Duplication models:

Fan Chung, Lu, Dewey, Galas; (2002)

Nodes are duplicated together with all (or part) of their edges.

can produce  $\gamma < 2$  as in biological networks  
 keep some network properties

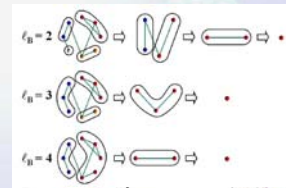
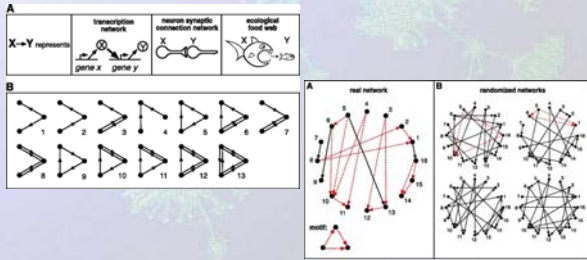


## Fractal networks

### Motifs, graphlets

Milo, Shen-Orr, Itzkoviz, Kashtan, Chkrovskii, Alon  
*Science* 298, 824-827 (2002)

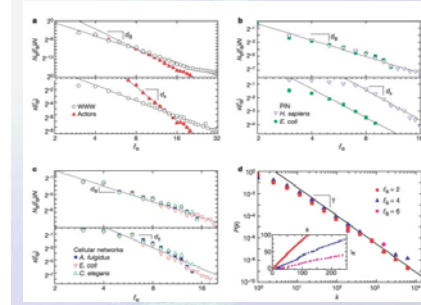
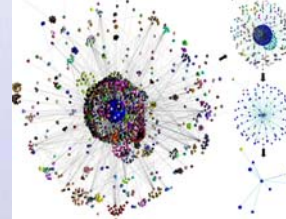
Pržulj, Corneil, Jurisica  
*Bioinformatics* 20, 3508-3515 (2004)



Song, Havlin, Makse  
*Nature* 433, 392-396 (2005)  
*Nature Physics* 2, 275-281 (2006)

Self-similarity of complex networks

Origins of fractality in the growth of complex networks



$$N_B \approx \ell_B^{-d_B}$$

$$k \rightarrow k' = s(\ell_B)k$$

$$s(\ell_B) \approx \ell_B^{-d_k}$$

$$\gamma = 1 + d_B/d_k$$

WWW, protein interaction networks are fractal

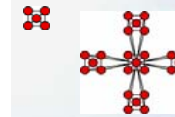
Internet (AS) is not fractal

Barabási-Albert is not fractal

Real complex networks: self-organized criticality (SOC) by some optimization process !!

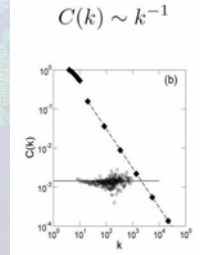
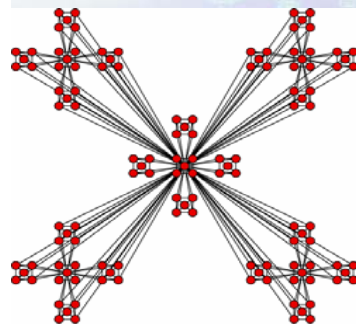
Cliques-trees, as deterministic models for real networks.

Hierarchical graphs  
Recursive clique-trees  
Apollonian graphs



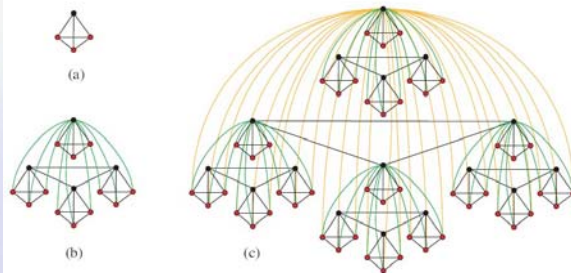
### Hierarchical graphs

Ravasz, Barabasi, Hierarchical organization in complex networks  
*Phys. Rev. E* (2003).



$$\gamma = 1 + (\ln 5 / \ln 4)$$

Real-life networks are fractal (Song, Havlin, Makse) but some fractal-looking graphs are not !



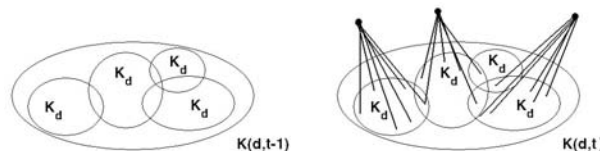
Barrière, Comellas, Dalfó (2006)

$$\gamma \approx 1 + \frac{\ln(d+1)}{\ln d}$$

### Recursive clique-trees

SN Dorogovtsev, AV Gotsev, JFF Mendes., *Phys. Rev. E* (2002)

F. Comellas, Guillaume Fertin, André Raspaud, *Phys. Rev. E* (2004)



### Recursive clique-trees

F. Comellas, G. Fertin, A. Raspaud, *Phys. Rev. E*

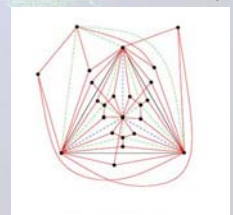
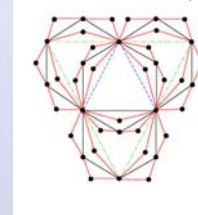
Initial graph:  $K_q$  - the complete graph with  $q$  vertices-

Operation:  $t \rightarrow t+1$ , obtain  $K(q,t+1)$  from  $K(q,t)$  by adding for every clique  $K_q$  of  $K(q,t)$ :

a: A new vertex  $u$

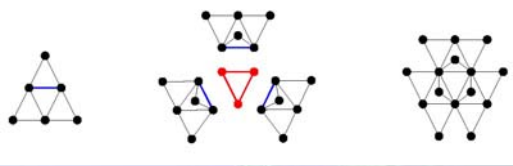
b:  $q$  edges joining  $u$  with the vertices of this clique

example  $q=2$



example  $q=3$

### Recursive equivalent operation



•Order, size

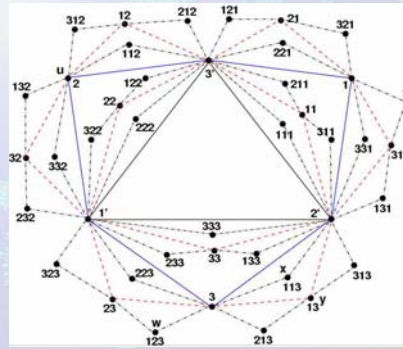
•Degree distribution  $\gamma \approx 1 + \frac{\ln(d+1)}{\ln d}$   $2 < \gamma < 2.58496$

•Clustering  $0.8 \leq C \leq 1$

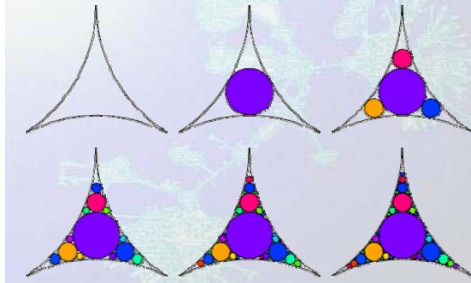
•Diameter logarithmic

### Distance-labeling and routing (example q=2)

F.Comellas, G. Fertin, A. Raspaud, *Sirocco 2003*



### Apollonian graphs

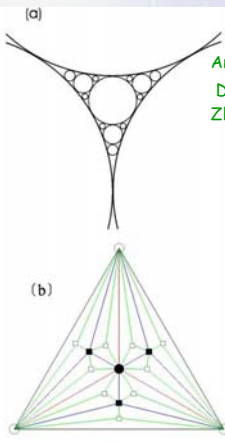


Apollonian packings

And let us not confine our cares  
To simple circles, planes and spheres,  
But rise to hyper flats and bends  
Where kissing multiple appears,  
In n-ic space the kissing pairs  
Are hyperspheres, and Truth declares -  
As n+2 such osculate  
Each with an n+1 fold mate  
The square of the sum of all the bends  
Is n times the sum of their squares.

Thorold Gosset, The Kiss Precise, Nature 139 (1937) 62.

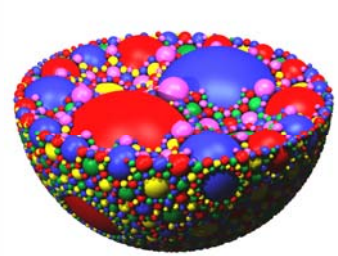
### Apollonian graphs



Andrade et al. *Phys. Rev. Lett.* (2005)

Doye, Massen *Phys. Rev. E* (2005)

Zhang, Rong, Comellas, Fertin *J. Phys.A* (2006)



<http://graphics.ethz.ch/~peikert/personal/packing/images/apoll3d.png>

d=2

d=3

Step (t)	New edges	Number of $K_{d+1}$
0	$\frac{d(d+1)}{2}$	1
1	$d+1$	$d+1$
2	$(d+1)^2$	$(d+1)^2$
3	$(d+1)^3$	$(d+1)^3$
...	...	...
i	$(d+1)^i$	$(d+1)^i$
i+1	$(d+1)^{i+1}$	$(d+1)^{i+1}$
...	...	...

$$N_t = (d+1) + \sum_{j=0}^{t-1} (d+1)^j = \frac{(d+1)^t - 1}{d} + d + 1$$

$$|E|_t = \frac{d(d+1)}{2} + \sum_{j=1}^t (d+1)^j = \frac{d(d+1)}{2} + \frac{(d+1)^{t+1} - d - 1}{d} \quad (1)$$

TABLE II: Distribution of vertices and degrees for  $A(d,t)$  at each step  $t$ .

Step(t)	Num. vert.	Degree
0	$d+1$	$d$
1	$d+1$	$d+1$
2	$d+1$	$d+1$
3	$d+1$	$d+1+d$
4	$d+1$	$(d+1) + (d+1)$
5	$d+1$	$d+1$
6	$d+1$	$d^2 + d + 1 + d$
7	$d+1$	$(d+1)d + (d+1) + (d+1)$
8	$d+1$	$(d+1) + (d+1)$
9	$(d+1)^2$	$d+1$
10	$(d+1)^2$	$d+1$
11	$(d+1)^2$	$d+1$
12	$(d+1)^2$	$d+1$
13	$(d+1)^2$	$d+1$
14	$(d+1)^2$	$d+1$
15	$(d+1)^2$	$d+1$
16	$(d+1)^2$	$d+1$
17	$(d+1)^2$	$d+1$
18	$(d+1)^2$	$d+1$
19	$(d+1)^2$	$d+1$
20	$(d+1)^2$	$d+1$
21	$(d+1)^2$	$d+1$
22	$(d+1)^2$	$d+1$
23	$(d+1)^2$	$d+1$
24	$(d+1)^2$	$d+1$
25	$(d+1)^2$	$d+1$
26	$(d+1)^2$	$d+1$
27	$(d+1)^2$	$d+1$
28	$(d+1)^2$	$d+1$
29	$(d+1)^2$	$d+1$
30	$(d+1)^2$	$d+1$
31	$(d+1)^2$	$d+1$
32	$(d+1)^2$	$d+1$
33	$(d+1)^2$	$d+1$
34	$(d+1)^2$	$d+1$
35	$(d+1)^2$	$d+1$
36	$(d+1)^2$	$d+1$
37	$(d+1)^2$	$d+1$
38	$(d+1)^2$	$d+1$
39	$(d+1)^2$	$d+1$
40	$(d+1)^2$	$d+1$
41	$(d+1)^2$	$d+1$
42	$(d+1)^2$	$d+1$
43	$(d+1)^2$	$d+1$
44	$(d+1)^2$	$d+1$
45	$(d+1)^2$	$d+1$
46	$(d+1)^2$	$d+1$
47	$(d+1)^2$	$d+1$
48	$(d+1)^2$	$d+1$
49	$(d+1)^2$	$d+1$
50	$(d+1)^2$	$d+1$
51	$(d+1)^2$	$d+1$
52	$(d+1)^2$	$d+1$
53	$(d+1)^2$	$d+1$
54	$(d+1)^2$	$d+1$
55	$(d+1)^2$	$d+1$
56	$(d+1)^2$	$d+1$
57	$(d+1)^2$	$d+1$
58	$(d+1)^2$	$d+1$
59	$(d+1)^2$	$d+1$
60	$(d+1)^2$	$d+1$
61	$(d+1)^2$	$d+1$
62	$(d+1)^2$	$d+1$
63	$(d+1)^2$	$d+1$
64	$(d+1)^2$	$d+1$
65	$(d+1)^2$	$d+1$
66	$(d+1)^2$	$d+1$
67	$(d+1)^2$	$d+1$
68	$(d+1)^2$	$d+1$
69	$(d+1)^2$	$d+1$
70	$(d+1)^2$	$d+1$
71	$(d+1)^2$	$d+1$
72	$(d+1)^2$	$d+1$
73	$(d+1)^2$	$d+1$
74	$(d+1)^2$	$d+1$
75	$(d+1)^2$	$d+1$
76	$(d+1)^2$	$d+1$
77	$(d+1)^2$	$d+1$
78	$(d+1)^2$	$d+1$
79	$(d+1)^2$	$d+1$
80	$(d+1)^2$	$d+1$
81	$(d+1)^2$	$d+1$
82	$(d+1)^2$	$d+1$
83	$(d+1)^2$	$d+1$
84	$(d+1)^2$	$d+1$
85	$(d+1)^2$	$d+1$
86	$(d+1)^2$	$d+1$
87	$(d+1)^2$	$d+1$
88	$(d+1)^2$	$d+1$
89	$(d+1)^2$	$d+1$
90	$(d+1)^2$	$d+1$
91	$(d+1)^2$	$d+1$
92	$(d+1)^2$	$d+1$
93	$(d+1)^2$	$d+1$
94	$(d+1)^2$	$d+1$
95	$(d+1)^2$	$d+1$
96	$(d+1)^2$	$d+1$
97	$(d+1)^2$	$d+1$
98	$(d+1)^2$	$d+1$
99	$(d+1)^2$	$d+1$
100	$(d+1)^2$	$d+1$

Discrete degree spectrum (with larger and larger jumps)

$$P_{cum}(k) \equiv \sum_{k' \geq k} N(k', t) / N_t \sim k^{1-\gamma}$$

$$\gamma \approx 1 + \frac{\ln(d+1)}{\ln d}$$

$$2 < \gamma < 2.58496$$

### Random Apollonian graphs

Instead of adding simultaneously a new vertex to each clique (never used before), we add a unique vertex to a random clique.

Initially  $A(d,0)$  is  $K_{d+2}$

Step  $t$  choose clique  $K_{d+1}$  NEVER USED and add a node (and the corresponding edges)

Order increments by 1 at each step

$$N_t = t + d + 2$$

### Degree distribution (self-averaging)

Given a vertex, when its degree increases by 1, the number of  $K_{d+1}$  which contains it increases by  $d-1$ . Thus, when the vertex attains degree  $k$ , the number of  $K_{d+1}$  is  $(d+1) \cdot (k-d-1) \cdot (d-1) = (d-1) \cdot k_i - d^2 + d + 2$

$$\frac{\partial k_i}{\partial t} = \frac{(d-1)k_i - d^2 + d + 2}{dt + d + 2}$$

with initial condition  $k_i(t_i) = d+1$  we obtain

$$k_i(t) = \frac{d^2 - d - 2}{d-1} + \frac{d+1}{d-1} \left( \frac{dt + d + 2}{dt_i + d + 2} \right)^{\frac{d-1}{d}}$$

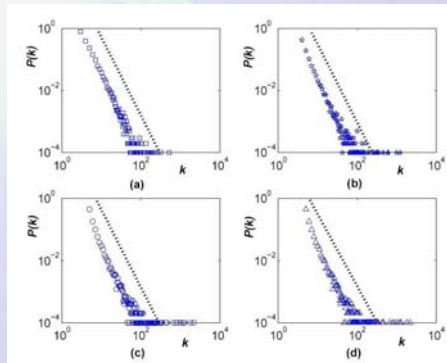
$$P(k_i(t) < k) = P\left(t_i > \frac{(dt + d + 2) \left(\frac{d+1}{d-1}\right)^{\frac{d}{d-1}} - d + 2}{d \left(k - \frac{d^2 - d - 2}{d-1}\right)^{\frac{d}{d-1}} - d + 2}\right)$$

$$P(k) = d(d+1)^{\frac{d}{d-1}} \left( (d-1)k - (d^2 - d - 2) \right)^{\frac{1-2d}{d-1}}$$

If  $k \gg d$  we have  $P(k) \sim k^{-\gamma}$  with  $\gamma(d) = \frac{2d-1}{d-1}$

$\gamma=3$  (for  $d=2$ , random seq.)

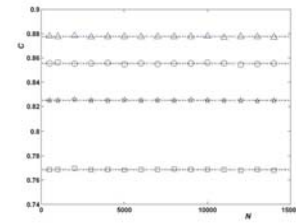
vs  $\gamma=2.58496$  ( $d=2$  parallel)



Degree distribution when  $N=10000$ ,  $d=2,3,4,5$

$$C(k) = \frac{\frac{d(d+1)}{2} + d(k-d-1)}{\frac{k(k-1)}{2}} = \frac{d(2k-d-1)}{k(k-1)} \quad \text{Clustering}$$

$$C = \int_{d+1}^{\infty} C(k)P(k)dk = \int_{d+1}^{\infty} \frac{d^2(2k-d-1)(d+1)^{\frac{d}{d-1}}}{k(k-1)} \left( (d-1)k - (d^2-d-2) \right)^{\frac{1-2d}{d-1}} dk$$



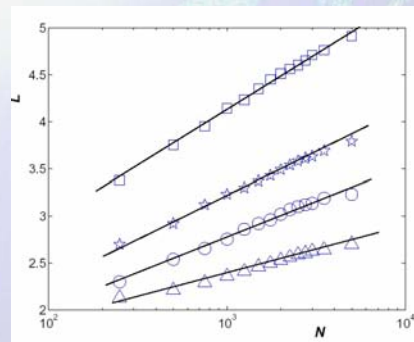
$$C = \frac{46}{3} - 36 \ln \frac{3}{2} = 0.7366$$

$$C = 18 + 36\sqrt{2} \arctan \sqrt{2} + \frac{9}{2}\pi - 18\sqrt{2}\pi = 0.8021$$

HDRAN clustering

$N=10000$ ,  $d=2,3,4,5$

### Average path length



- Zhongzhi Zhang, Lili Rong, F. Comellas, Guillaume Fertin, **High dimensional Apollonian networks** *J. Phys. A* (2006)
- Zhongzhi Zhang, Lili Rong, F. Comellas, **High dimensional random Apollonian networks** *Physica A* (2006),

### Deterministic recursive clique-trees

	Adding at the same time a vertex to each $d$ -clique with repetition	Adding at the same time a vertex to each $d$ -clique without repetition
Case $d = 2$	<i>Pseudofractal scale-free</i> Dorogotsev, Goltsev, Mendes Phys.Rev.E 65 (2002) 066122	<i>Deterministic SW network</i> Zhang, Rong, Guo Physica A cond-mat/0503637
Case $d = 3$		<i>Apollonian network</i> Andrade, Herrmann, Andrade, Silva Phys.Rev.Lett. 94 (2005) 018702 Doye, Massen Phys. Rev. E 71 (2005) 016128.
General case $d = 2 \dots \infty$ (includes cases $d=2,3$ )	<i>Recursive clique-trees</i> Comellas, Fertin, Raspaud Phys.Rev.E 69 (2004) 037104.	<i>High dimensional Apollonian network</i> Zhang, Comellas, Fertin, Rong J. Phys. A. 39 (2006) 1811 (introduced by Doye and Massen, Phys. Rev. E 71 (2005) 016128.)

### Random recursive clique-trees

	Adding a single vertex to a random clique with repetition	Adding a single vertex to a random clique without repetition
Case $d = 2$		<i>Random SW network</i> Ozik, Hunt, Ott Phys.Rev.E 69 (2004) 02618
Case $d = 3$		<i>Random Apollonian network</i> Zhou, Yan, Wang Phys.Rev.E 71 (2005) 046141
General case $d = 2 \dots \infty$ (includes cases $d=2,3$ )	<i>Random recursive clique-tree</i> see Appendix	<i>HD random Apollonian network</i> Zhang, Comellas, Rong Physica A. cond-mat/0502591

### Deterministic vs Random

Graph family	$P(k)$ or $\gamma$ -exponent	Clustering
Deterministic SW [78]	$2^{-\frac{k}{d}}$	$0.69 = \ln 2$
Random SW [77]	$\frac{3}{4} \left(\frac{2}{3}\right)^{-k}$	$0.65 (= \frac{3}{2} \ln 3 - 1)$
Apollonian [7,34]	$2.58 (= 1 + \frac{\ln 3}{\ln 2})$	0.83
Random Apollonian [82]	$\frac{3N-5}{N} \approx 3$	$0.74 (= \frac{46}{3} - 36 \ln \frac{3}{2})$
High-Dim. Apollonian [81]	$1 + \frac{\ln(d+1)}{\ln d}$ (2 to 2.58)	0.83 to 1
High-Dim. Random Apollonian [80]	$\frac{2d-1}{d-1}$ 2 to 3	0.74 to 1
Pseudo fractal scale-free [29]	$1 + \frac{\ln 3}{\ln 2} = 2.58$	$0.80 (= \frac{4}{5})$
Random pseudo fractal scale-free	$\frac{5}{2} = 2.5$	
Determ. recursive clique-trees [22]	$1 + \frac{\ln(d+1)}{\ln d}$ (2 to 2.58)	0.80 to 1
Random rec. clique-trees [see Appendix]	$\frac{2d-1}{d-1}$ (2 to 3)	0.74 to 1

### Why the random approach produces a different distribution

F. Comellas, Hernan D. Rozenfeld, Daniel ben-Avraham  
**Synchronous and asynchronous recursive random scale-free nets**  
*Phys. Rev. E* (2005),

In many simulations choosing an edge might be biased.

It is not the same to choose edge  $e$  from  $|E|$  edges than choose a node and then an adjacent node.

## Present and future work in SW-SF networks

How to construct a better WWW (new topologies -Akamai) ?

How to analyse very large graphs ?

- mean field and other statistical methods
- fractal techniques
- spectral theory
- new invariants
- this workshop

How to deal with dynamical networks ?

New communication protocols

