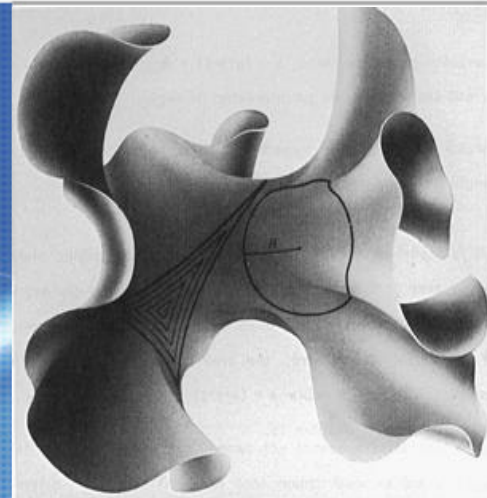


# *Network curvature, and its implications for network management\**



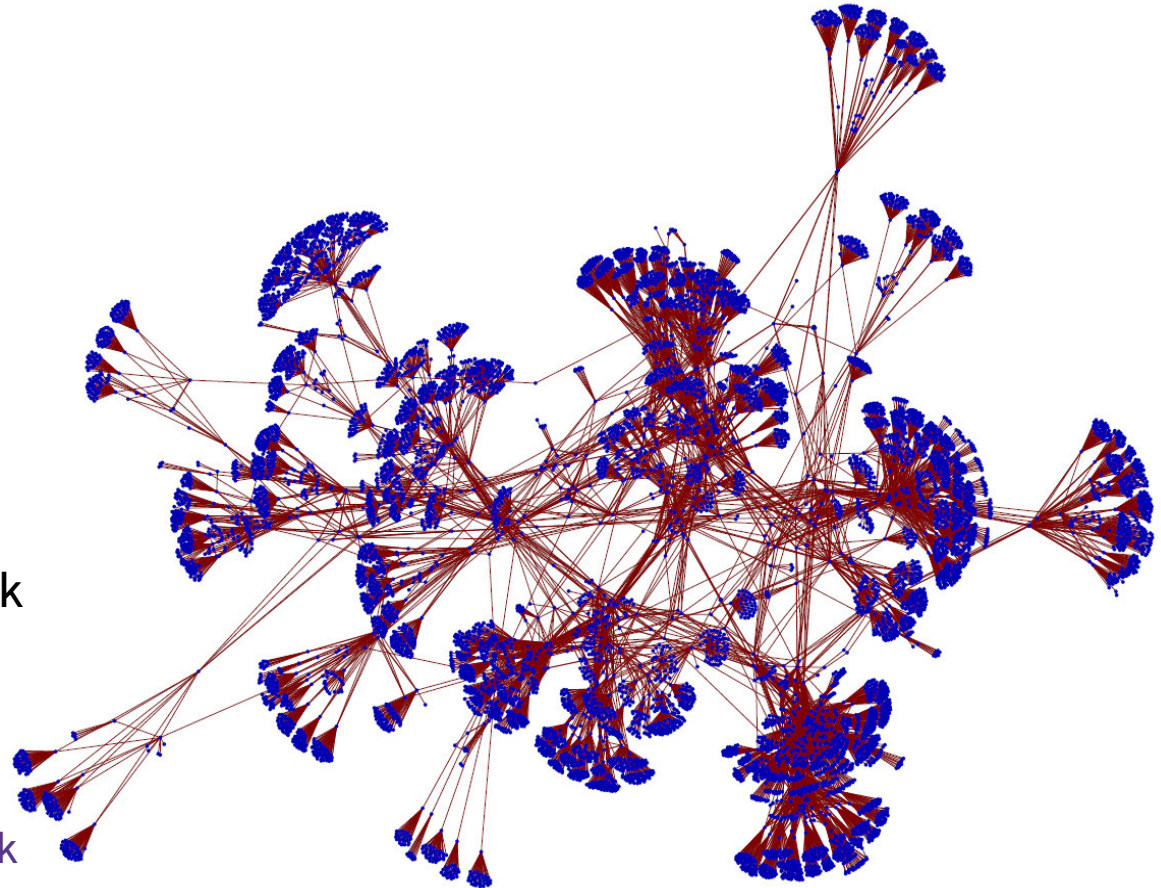
**Iraj Saniee**, *Math of Networks & Comm. Dept.*, Bell Labs, Alcatel-Lucent, Murray Hill, NJ  
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DIMACS Workshop, November 12-13, 2009

# Understanding Large-Scale Networks

- Hard to visualize due to scale
- Unclear what is essential and what is not for overall *performance, reliability* and *security*
- Need better ways to “summarize” critical network information
- How?

A promising direction is to look at two main geometric characteristics of objects: their *dimension* & *curvature*



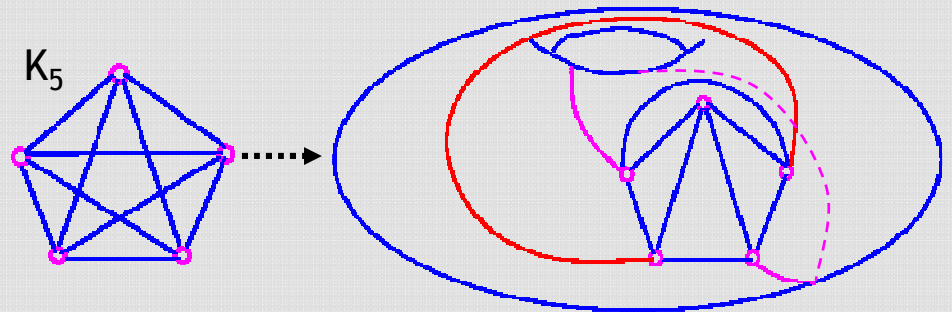
Rocketfuel dataset 7018 (AT&T)  
10152 nodes, 28638 links, diameter 12

# Local Dimension and Graph Genus

How “planar” is a given network? Turns out all graphs are *locally planar*.

**Ringel-Youngs.** (“All graphs with  $N \geq 3$  nodes are locally 2-dimensional.”) For  $N \geq 3$ , any  $G=(N,L)$  can be embedded in  $T^g$ , a torus with  $g$  holes, where

$$g \leq \lceil (N-3)(N-4)/12 \rceil$$



**Issue.** Even if computation of  $g$  were easy, R-Y wouldn't help much because the advantage of local 2-D structure is obscured by impact of holes on the large-scale properties of graph. *How so?*

## Our Data Source -- Rocketfuel

Look at scaling of the **average shortest path length  $\langle h \rangle$** . In 2-dim grid,  $\langle h \rangle \sim \sqrt{N}$  (or  $\sim N^{1/\Delta}$  in  $\Delta$ -dim grid) but many real networks have slower scaling of  $\langle h \rangle$  than  $\sqrt{N}$  -- e.g., think of the “small world” E-R random graphs.

- Look at “Rocketfuel” data, [Washington University researchers’ detailed connectivity data from various ISPs 2002-2003]
- Clearly,  $\langle h \rangle$ , average SPs with respect to the hop metric, do not scale like  $\sqrt{N}$
- Need to consider something else!

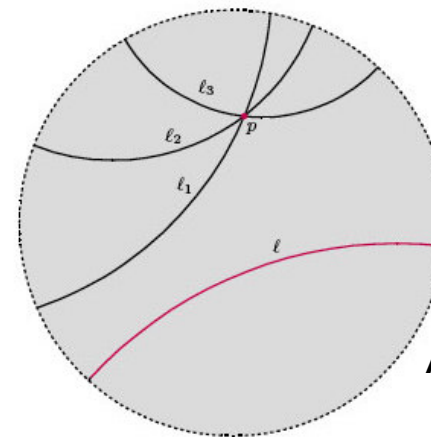
Network ID	Network Name	Size #node - #links	Average Shortest Path Length
1221	Telstra (Australia)	2998 - 7612	5.53
1239	Sprintlink (US)	8341 - 28050	5.18
1755	EBONE (US)	605 - 2070	6.0
2914	Verio (US)	3045 - 10726	6.0
3257	Tiscali (EU)	855 - 2346	5.3
3356	Level 3 (US)	3447 - 18780	5.0
3967	Exodus (US)	895 - 4140	5.9
4755	VSNL (India)	121 - 456	3.2
6461	Abovenet (US)	2720 - 7648	5.7
7018	AT&T (US)	10152 - 28638	6.9

## Other Locally 2-Dimensional Embeddings: The Poincaré Disk $H^2$

Consider the unit disk  $\{x \in R^2; |x| < 1\}$  with length metric given by

$$ds^2 = \frac{dx^2 + dy^2}{(1 - x^2 - y^2)^2}$$

the *hyperbolic* metric.



A few geodesics

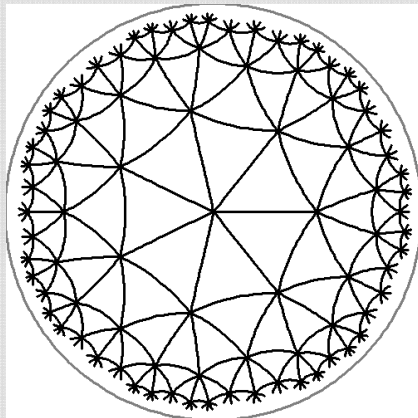
**Advantages.** In *the small scale* it is 2-dimensional, but has *slower* scaling of geodesics (shortest paths) than  $\sqrt{N}$ .

**Relationship to graphs?** The Poincaré disk comes with numerous natural “scaffoldings” or tilings.

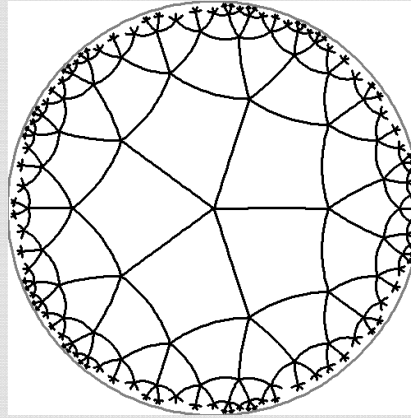
# Scaffoldings of $H^2$ : Hyperbolic Regular Graphs

Consider  $X_{p,q}$ , tilings (isometries) of  $H^2$ , that at each vertex consist of  $q$  regular  $p$ -gons for integers  $p$  &  $q$  with  $(p-2)(q-2) > 4$  (flat with equality)

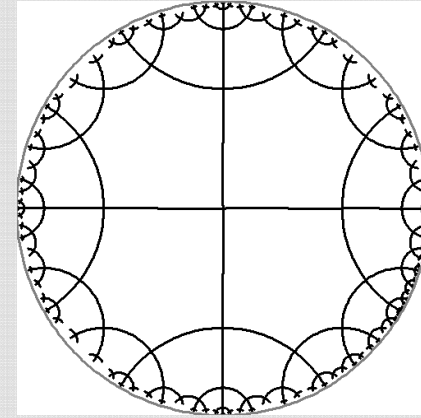
Examples:



$X_{3,7}$



$X_{4,5}$



$X_{6,4}$

**Note.** Since networks of interest to us are typically finite, we'll consider truncations of  $X_{p,q}$ , the part within a (large enough) radius  $r$  from the center. Call this  $TX_{p,q}$ .

# Some Key Properties of $TX_{p,q}$

1. **Negative local curvature.** The local angular deviation from  $2\pi$  (Gauss-Bonnet curvature) of  $X_{p,q}$  at each node

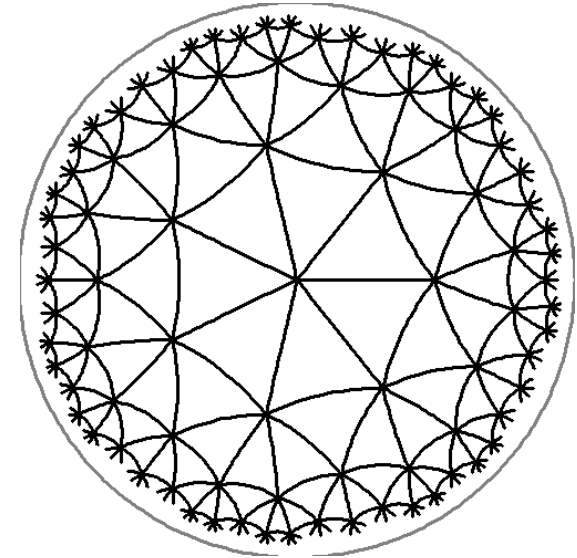
$$2\pi\{4 - (p-2)(q-2)\} < 0$$

gives the (combinatorial curvature) at node  $v$

$$\kappa_v = \frac{1}{p} + \frac{1}{q} - \frac{1}{2} < 0$$

- 2 . **Exponential growth.** Number of nodes within a ball of radius  $r$  is proportional to  $\lambda^r$  for some  $\lambda \equiv \lambda(p,q) > 1$  (e.g., for  $X_{3,7}$ ,  $\lambda = \phi$ , the golden ratio) or equivalently

- 2'. **Logarithmic scaling of geodesics.** For a (finite truncation of)  $X_{p,q}$  with  $N$  nodes, the average geodesic (shortest path length) scales like  $O(\log(N))$



$X_{3,7}$

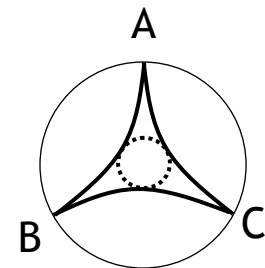
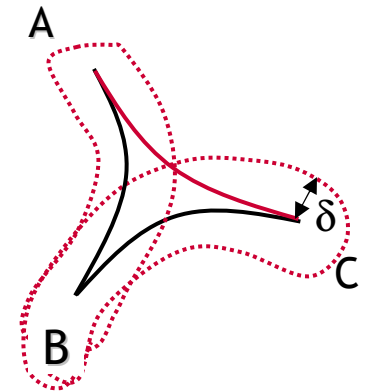
# Local Curvature versus Curvature in the Large

Computation of *total* curvature of non-flat networks with varying nodal degrees via  $\sum_{v \in G} \kappa_v$  does not appear to be possible/easy to provide information about *the large-scale* properties of the network

A more direct definition of (*negative*) *curvature in the large* is the thin-triangle condition for a geodesic metric space (or a  $CAT(-\kappa)$  space):

(M. Gromov's Thin Triangle Condition for a hyperbolic geodesic metric space) There is a (minimal) value  $\delta \geq 0$  such that for *any* three nodes of the graph connected to each other by geodesics, each geodesic is within the  $\delta$ -neighborhood of the union of the other two.

**Example.** For  $H^2$ ,  $\delta = \ln(\sqrt{2} + 1)$ . [Sketch. Largest inscribed circle must be in largest area triangle,  $\text{Area}_H(ABC) = \pi - (\alpha + \beta + \chi)$ , maximized to  $\pi$  when  $\alpha, \beta, \chi = 0$  or when A, B, & C are on the boundary.]





# What About Communication Networks?

Communication networks are geodesic metric spaces via reasonable link metrics (e.g., the hop metric)

Is there evidence for negative curvature in *real* communication networks?

We consider 10 Rocketfuel networks and some prototypical flat and curved *famous* networks to test this hypothesis

Network ID	Network Name	Size #node - #links	Diameter
1221	Telstra (Australia)	2998 - 7612	12
1239	Sprintlink (US)	8341 - 28050	13
1755	EBONE (US)	605 - 2070	13
2914	Verio (US)	3045 - 10726	13
3257	Tiscali (EU)	855 - 2346	14
3356	Level 3 (US)	3447 - 18780	11
3967	Exodus (US)	895 - 4140	13
4755	VSNL (India)	121 - 456	6
6461	Abovenet (US)	2720 - 7648	12
7018	AT&T (US)	10152 - 28638	12
TX(3,7)	Synthetic	4,264-15,022	14
Power-law (Albert-Barabasi)	Synthetic	10,000 - 39,994	9
Strogatz-Watts (2D)	Synthetic (p=0.2)	80x80 - 26,578	20
Truncated (3,6) Grid	Synthetic	469 - 2,520	24
Truncated (4,4) Grid	Synthetic	80x80 - 25,280	158

# Experiments and Methodology

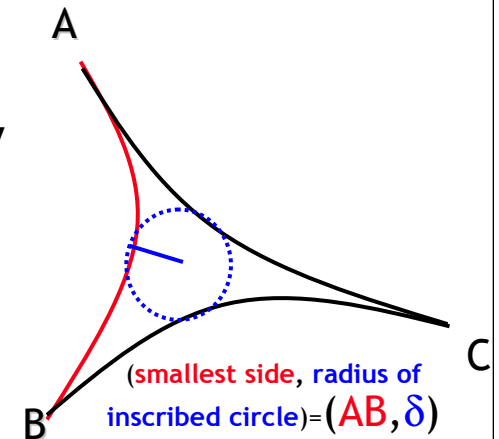
We ran experiments on all Rocketfuel networks plus a few prototypical flat/curved networks to test our key hypothesis:

## 1. Dimension. “Growth test” - Polynomial or exponential?

- Consider the volume  $V(r)$  as a function of radius  $r$  for arbitrary centers
- [In flat graphs volume growth is typically polynomial in radius  $r$ ]

## 2. Curvature. “Triangle test” - Are triangles are universally $\delta$ -thin

- Randomly selected 32M, 16M, 1.6M triangles for networks with more than 1K nodes and exhaustively for the remainder
- For each triangle noted shortest side  $L$  and computed the  $\delta$
- Counted number of such triangles, indexed by  $\delta$  and  $L$
- [In flat graphs  $\delta$  grows without bound as the size of the smallest side increases]



We conduct “growth” and “triangle” tests

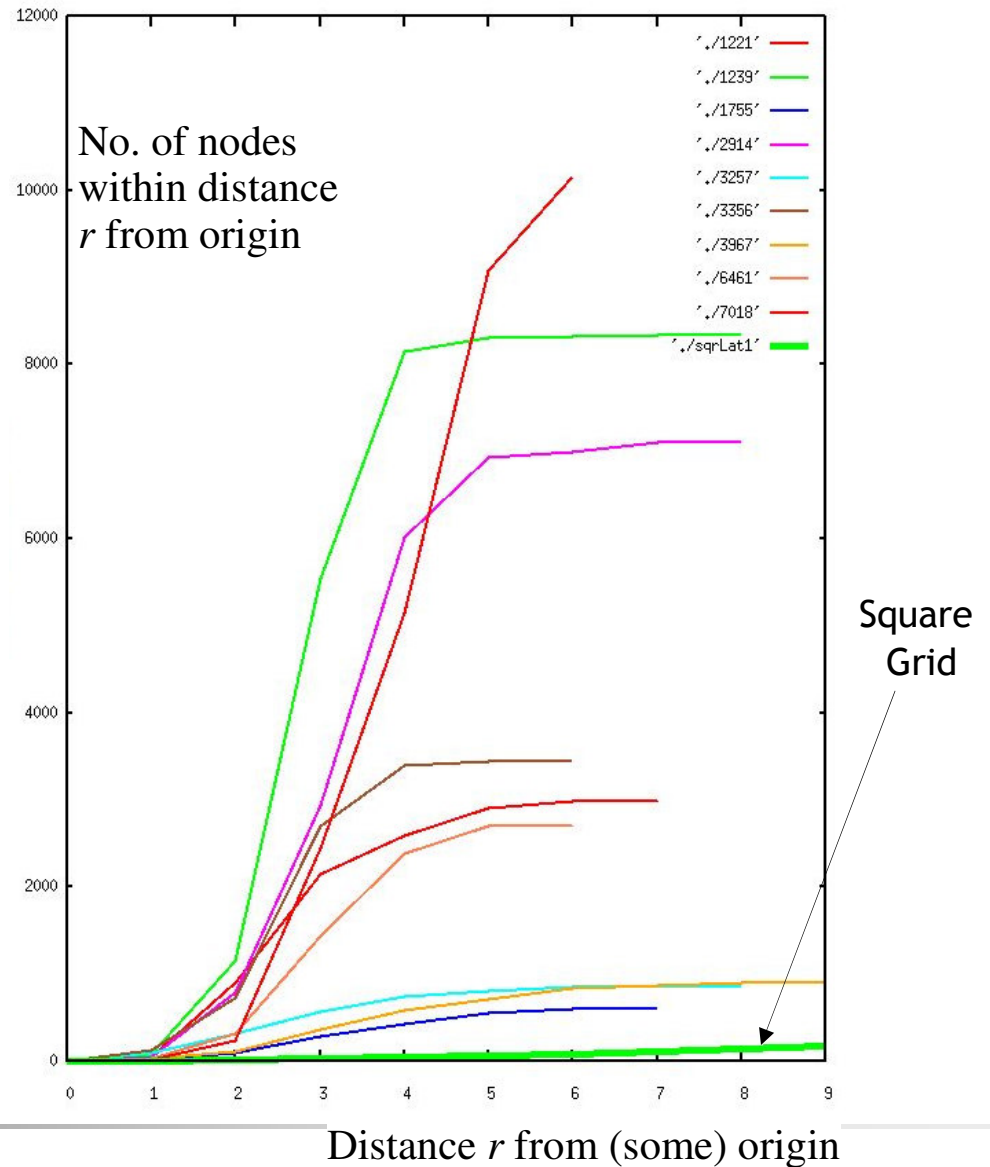
# 1. Growth Charts

Recall that:

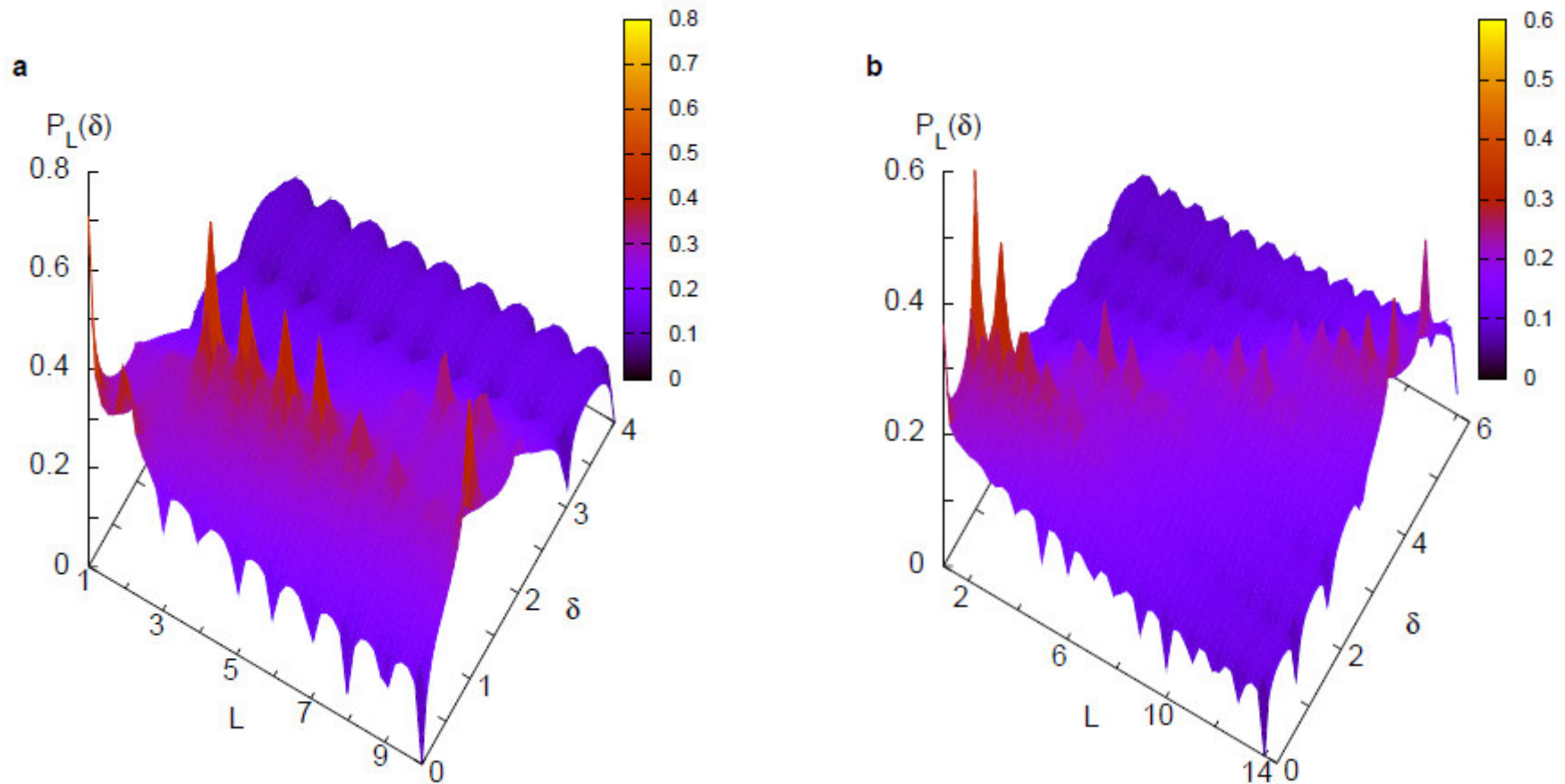
Euclidean growth  $V(r) \approx r^D$

Exponential growth  $V(r) \approx \theta^r$

*Volume* (number of points with distance  $r$ ) as a function of *radius*  $r$  from a “center” of the graph. Flattening of curves for larger  $r$  is due to approach to boundary / finite size of network.

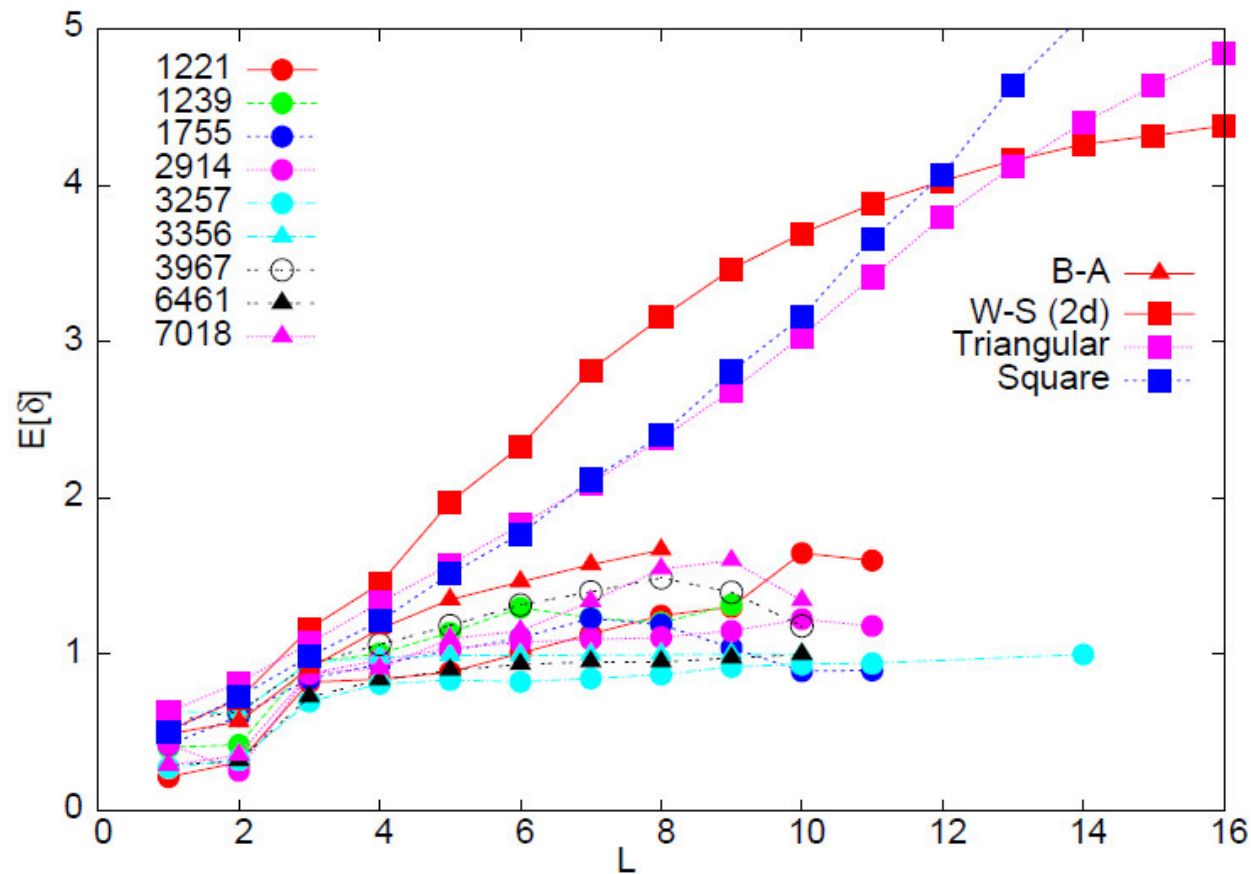


## 2. Triangle Test - Rocketfuel 7018 (AT&T) & Triangular Grid



- (a) Probability  $P_L(\delta)$  for randomly chosen triangles whose shortest side is  $L$  to have a given  $\delta$  for the network 7018(AT&T network) which has 10152 nodes and 14319 bi-directional links and diameter 12. The quantities  $\delta$  and  $L$  are restricted to integers, and the smooth plot is by interpolation.
- (b) Similar to (a), for a (flat) triangular lattice with 469 nodes and 1260 links. (The smaller number of nodes is sufficient for comparing with (a) since the range for  $L$  is large due to the absence of the small world effect.)

# Summary of Triangle Tests for Rocketfuel Networks



The average  $\delta$  as a function of  $L$ ,  $E[\delta](L)$ , for the 10 IP-layer networks studied here, and for the Barabasi-Albert model with  $k = 2$  and  $N = 10000$  (11th curve) and the hyperbolic grid  $X_{3,7}$  (12th curve). On the other hand, a Watts- Strogatz type model on a square lattice with  $N = 6400$ , open boundary conditions and 5% extra random connections (13<sup>th</sup> curve) and two flat grids (the triangular lattice with diameter 29 and the square lattice with diameter 154) are also shown.

## Where to go from here?

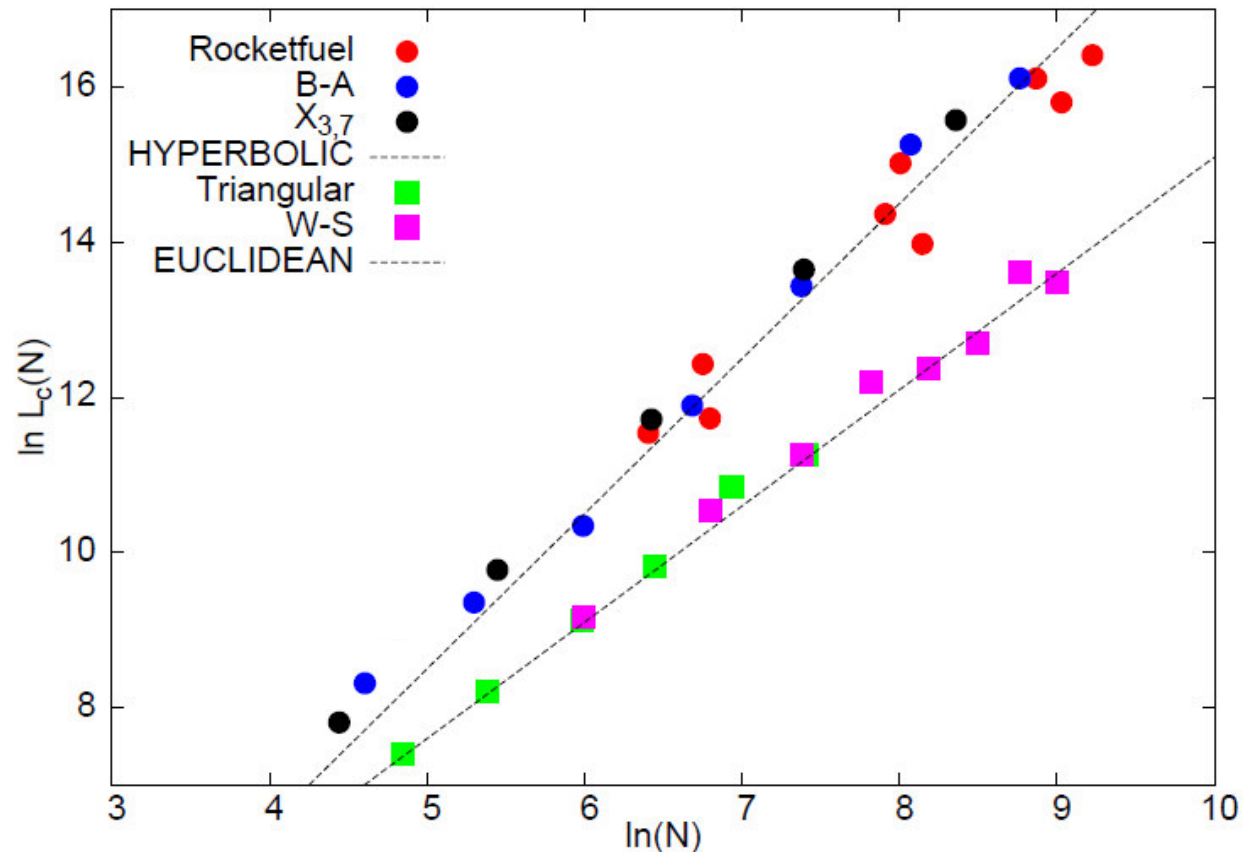
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- OK, these ten RF datasets and some “well-known” large-scale networks exhibit
  - Exponential growth / logarithmic scaling of shortest paths
  - Negatively curvature in the large

### *So what?*

- Turns out negatively-curved structures exhibit specific features that affect critical properties of networks
  - Existence of a “core”

# The Downside of Hyperbolicity: Quadratic Scaling of Load (“Betweenness Centrality” and Existence of “Core”)



Plot of the maximum load  $L_c(N)$  -- maximal number of geodesics intersecting at a node -- for each network in the Rocketfuel database as a function of the number of nodes  $N$  in the network. Also shown are the maximum load for the hyperbolic grid  $X_{3,7}$ , the Barabasi-Albert model with  $k = 2$ , the Watts-Strogatz model and a triangular lattice, for various  $N$ . The dashed lines have slopes of 2.0 and 1.5, corresponding to the hyperbolic and Euclidean cases respectively.

## Key Claims:

# Network Curvature -> Congestion, Reliability and Security

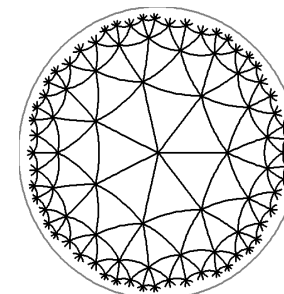
Congestion is not necessarily a manifestation of the heavy-tailed property of the distribution of the node degrees, but rather of a more subtle feature, given the (negative) curvature of the network.

Numerical studies show that congestion does not necessarily occur at vertices of high degree, nor at the so-called highly connected nodes but rather at the points relative to which the network has minimum inertia (the “core”).

### Geodesics (shortest paths)

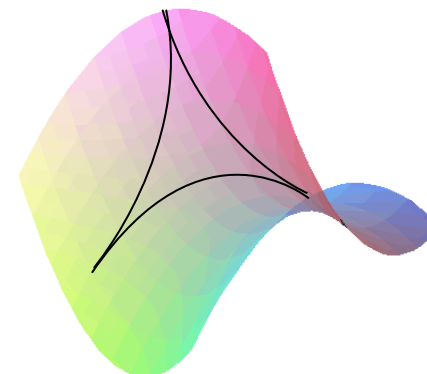
- (Upside) Are very effective, as diameter is small compared to  $N$ , e.g., TTL of  $\sim 20$  good enough for all of the Internet!
- (Downside) Lead to
  - congestion
  - non-random failure can be very severe
  - certain nodes exhibit significant security threats

### Need for non-geodesic routing to avoid the downside



$X_{3,7}$

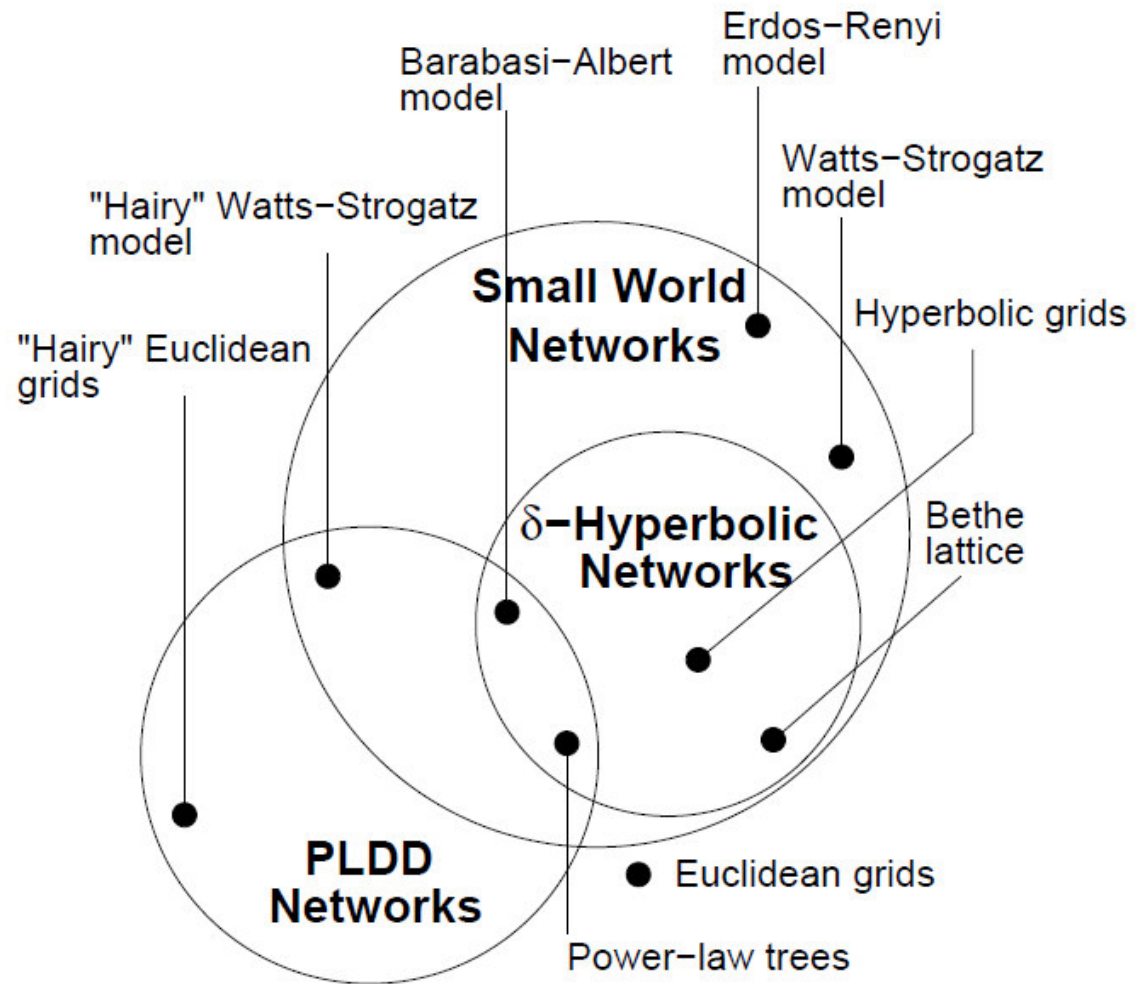
Nodal *loads* need not be related to nodal *degrees*



Geodesic triangles on negatively curved spaces: thin regardless of the edge lengths



# Taxonomy for Large-Scale Networks



Taxonomy of key characteristics of networks and their overlaps in a schematic diagram.

## To Summarize

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- Networks, in addition to the *small world* property -- that has been widely verified -- and the power law degree distribution -- that is often but not always observed -- have another signature: *curvature*
- Rocketfuel ISP networks exhibit negative curvature *and* small world behavior and (within numerical accuracy) power law degree distribution
- Synthetic network models can show negative curvature (B-A) or not (S-W)
- Analytical prediction on continuum and hyperbolic lattices shows  $N^2$  scaling load at their “core”. This agrees with nodal power-law and Rocketfuel, but not with S-W.
- Thus hyperbolicity seems sufficient AND necessary for an important network performance measure. Real networks show this behavior, so it is relevant!
- Need for routing protocols that don't fall into *geodesic grooves*

# References

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- [1] O. Narayan, I. Saniee, The Large Scale Curvature of Networks, arXiv:0907.1478 (July 2009)
- [2] O. Narayan, I. Saniee, Scaling of load in communications networks, arXiv:0906.4138 (June 2009)