Profile vectors in the poset of subspaces

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Motivation

For a family \mathcal{F} of subsets of $[n] = \{1, 2, ..., n\}$ P.L. Erdős, P. Frankl, G.O.H. Katona (1985) defined

$$p(\mathcal{F}) \in \mathbf{R}^{n+1}, p(\mathcal{F})_i = |\{F \in \mathcal{F} : |F| = i\}|.$$

For a set A of families they started to investigate $\langle \mu(\mathbf{A}) \rangle$, the convex hull of the profiles of the families in A.

In the same way for a family \mathcal{U} of subspaces of an *n*-dimensional vectorspace *V* over GF(q) we define

 $p(\mathcal{U}) \in \mathbf{R}^{n+1}, p(\mathcal{U})_i = |\{U \in \mathcal{U} : \dim U = i\}|.$

The reduction method

The method provides a superset of the profile polytope if we know the profile polytope of the families reduced to some simpler structure.

In the Boolean poset, reduction to

- maximal chain ({ $C_0 \subseteq C_1 \subseteq ... \subseteq C_n$ } with $|C_i| = i$),
- pair of complement chains (union of two maximal chains C₁ and C₂ such that C ∈ C₁ iff [n] \ C ∈ C₂),
- circle (the intervals (sets of consecutive elements) for a fixed cyclic permutation of [n]).

In the poset of subspaces, reduction to

• chain,

• pair of orthogonal chains.

Easy consequences I

Using the reduction method one can determine the profile polytope of the following sets of families (among others):

- *k*-Sperner families,
- orthogonal complement-free families,
- families without pairs of subspaces U_1, U_2 with $U_1 \cap U_2 = \{\underline{0}\}, \langle U_1 \cup U_2 \rangle = V$ (property *).

Theorem. The extreme points of the profile polytope of k-Sperner families are the following:

for all $0 \le z \le k$ and $\beta = \{\beta_1, ..., \beta_z\}$ with $0 \le \beta_1 < ... < \beta_z \le n$ the vectors v_β

$$(v_{eta})_i = \left\{ egin{array}{cc} [n] & ext{if} & i \in eta \ 0 & ext{otherwise.} \end{array}
ight.$$

(1)

Easy consequences II

Theorem. The extreme points of the convex hull of the profile vectors of orthogonal complement-free families are the vectors corresponding to the families consisting of

a, a set I of levels with the property that the *i*th and the n - ith levels cannot be both in I, if n is odd,

b, a set I of levels with the property that the *i*th and the n - ith levels cannot be both in I and possibly half of the subspaces with dimension n/2 one from each pair of orthogonal complementary subspaces, if n is even.

Theorem. The extreme points of the profile polytope of families with property * are the vectors corresponding to the families consisting of

a, a set I of levels with the property that the *i*th and the n - ith levels cannot be both in I, if n is odd,

b, a set I of levels with the property that the *i*th and the n - ith levels cannot be both in I and possibly a maximal star of subspaces (subspaces containing a fixed vector) with dimension n/2, if n is even.

The method of inequalities

To do:

- finding (linear) inequalities valid for the profiles of all families in the set examined
- determining the extreme points of the polytope determined by the above inequalities
- finding a family for any extreme point that has this vector as profile

Intersecting families

Theorem. (Erdős - Ko - Rado for vectorspaces) If $\mathcal{U} \subseteq \begin{bmatrix} V \\ k \end{bmatrix}$ is a *t*-intersecting family and $n \ge 2k - t$, then

$$|\mathcal{U}| \le \max\{ {n-t \brack k-t}, {2k-t \brack k} \}.$$

- Hsieh '77 for large enough ns (including $t = 1, n \ge 2k + 1$)
- Greene Kleitman '78 for t = 1, k|n (so for n = 2k as well)
- Frankl Wilson '86 for all cases

A generalization of Hsieh's theorem

Theorem. The following generalization of Hsieh's theorem holds:

a, if $2k \le n \le 2k+2$ and d=0 or d=n-k

or

b, if $n \ge 2k+3$ and $k+d \le n$

then for any intersecting family \mathcal{U} of k-dimensional subspaces of an n-dimensional vectorspace V with all members disjoint from a fixed d-dimensional subspace W of V we have

$$\mathcal{U}| \le {\binom{n-1}{k-1}}^{*(d)},$$

where $\binom{n}{k}^{*(d)}$ denotes the number of *k*-dimensional subspaces of an *n*-dimensional vectorspace *V* that are disjoint from a fixed *d*-dimensional subspace *W* of *V*.

'Sketch' of the proof

Proof.

- calculating the eigenvalues of the apropriate incidency matrix (as Frankl and Wilson did) is not straightforward
- if k|d|n or k|n, d = 0 (so in particular when n = 2k, d = 0) the trick of Greene and Kleitman (partitioning $V \setminus W$ into isomorphic copies of $V_k \setminus \{\underline{0}\}$) works
- for the remaining cases the 'lengthy computations' of Hsieh go through, since he used 'mostly' the quick growth of the *q*-nomial coefficients, and the ratio of growth needed there is valid for the growth of

$$\frac{{\binom{n}{k}}^{*(d)}}{{\binom{n-1}{k-1}}^{*(d)}}$$

(for the q = 2 case we needed some further computations)

The inequalities

Lemma. For the profile vector f of any family \mathcal{U} of intersecting subspaces of an n-dimensional vectorspace, and for any k < n/2 and $n/2 < d \le n - k$, the following holds

$$c_{k,d}f_k + f_d \leq \begin{bmatrix} n \\ d \end{bmatrix},$$

where $c_{k,d} = q^d \frac{\binom{n-k}{d}}{\binom{n-d-1}{k-1}}$, and eqaulity holds if $f_k = 0, f_d = \binom{n}{d}$ or if $f_k = \binom{n-1}{k-1}, f_d = \binom{n-1}{d-1}$.

Proof. doublecount the disjoint pairs formed from \mathcal{U}_k and $\mathcal{U}'_d = \begin{bmatrix} V \\ d \end{bmatrix} \setminus \mathcal{U}_d$ giving

$$\binom{n}{d}^{*(k)} f_k \le \left(\binom{n}{d} - f_d \right) \binom{n-1}{k-1}^{*(d)}$$

The profile polytope

Theorem. The essential extreme points of the profile polytope of the set of intersecting families are the vectors $v_i \ 1 \le i \le n/2$ for even n and there is an additional essential extreme point v^+ for odd n, where

$$(v_i)_j = \begin{cases} 0 & \text{if } 0 \le j < i\\ \binom{n-1}{j-1} & \text{if } i \le j \le n-i\\ \binom{n}{j} & \text{if } j > n-i. \end{cases}$$
(2)

(3)

and

$$(v^+)_j = \begin{cases} 0 & \text{if } 0 \le j < n/2\\ \begin{bmatrix} n\\j \end{bmatrix} & \text{if } j > n/2. \end{cases}$$

Remark. This is the 'same' as in the Boolean case, just using q-nomial coefficients instead of binomial coefficients.

t-intersecting families

In the Boolean case the complete intersection theorem of Ahlswede and Khatchatrian (1997) settled the problem of the largest *t*-intersecting families when *n* is small compared to *k* and *t*. The optimal families vary 'widely' when *n* ranges from 2k - t to infinity.

By the theorem of Frankl and Wilson we know that in the case of subspaces there are only three possibilities for the optimal family.

'Conjecture'. Determining the profile polytope of t-intersecting families is much easyer in the poset of subspaces than in the Boolean case.