

On Approximate Horn Minimization

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- ▶ 1. Preliminaries
- ▶ 2. Different versions of the problem
- ▶ 3. Approximability
- ▶ 4. Inapproximability

Preliminaries

- ▶ **Horn clause**: at most one unnegated variable, e.g.
 - ▶ $a, b \rightarrow c$ or $\bar{a} \vee \bar{b} \vee c$
 - ▶ **definite** Horn clause: one unnegated variable
- ▶ **(definite) Horn formula**: conjunction of (definite) Horn clauses
 - ▶ $(a, b \rightarrow c) \wedge (a, c \rightarrow d) \wedge (d, e \rightarrow f)$
- ▶ **size**: number of clauses or variables
- ▶ **implication** (can be decided efficiently):
 - ▶ $(a, b \rightarrow c) \wedge (a, c \rightarrow d) \models (a, b \rightarrow d)$
- ▶ **equivalence** of φ and ψ :
 - ▶ $\forall a \varphi(a) \Leftrightarrow \psi(a)$
 - ▶ $\forall C \varphi \models C \Leftrightarrow \psi \models C$

Horn minimization

- ▶ given a Horn formula φ and a number k , is there a Horn formula with at most k clauses equivalent to φ ?
- ▶ $x_i \rightarrow y_j, \quad y_1, \dots, y_n \rightarrow x_i$ $n^2 + n$ clauses
- ▶ $x_i \rightarrow x_{i+1}, \quad x_n \rightarrow y_i, \quad y_1, \dots, y_n \rightarrow x_1$ $2n$ clauses
- ▶ **restricted version**: given a Horn formula φ and a number k , does φ have a **subformula** with at most k clauses equivalent to φ ?

Previous work

- ▶ Boros, Čepek (1994): NP-complete
- ▶ Hammer, Kogan (1993):
 - ▶ NP-complete if the number of literals is to be minimized; in P for quasi-acyclic formulas
 - ▶ (restricted) Horn minimization (number of clauses) has an $(n - 1)$ -approximation algorithm (n : number of different variables) $o(n)$?
- ▶ Maier (1983), Ausiello, D'Atri, Saccà (1986), Guigues, Duquenne (1986), Angluin, Frazier, Pitt (1992): minimization of the number of bodies can be done efficiently
- ▶ Boros, Čepek, Kogan (1997): iterative decomposition algorithm

Transitive reduction

- ▶ given a directed graph G and a number k , is there a directed graph with at most k edges having the same transitive closure? - **in P**
- ▶ given a directed graph G and a number k , does G have a **subgraph** with at most k edges having the same transitive closure? - **NP-complete**
- ▶ Berman, DasGupta, Karpinski (2007): 1.5-approximation algorithm
- ▶ definite Horn clauses of size 2

\mathcal{R} -equivalence

- ▶ Flögel, Kleine Büning, Lettmann (1993): introduce new variables
- ▶ $x_1/y_1, \dots, x_n/y_n \rightarrow u$ 2^n clauses
- ▶ $x_i \rightarrow z_i, y_i \rightarrow z_i, z_1, \dots, z_n \rightarrow u$ $2n + 1$ clauses
- ▶ same set of consequences over the original variables
- ▶ $\varphi \sim_{\mathcal{R}} \psi$, where \mathcal{R} is a set of variables: same set of consequences over \mathcal{R}
- ▶ *co* – *NP*-complete even in rather restricted cases

Narrow extension

- ▶ introduce new variables in a **restricted** way
- ▶ $x_i \rightarrow y_j, i = 1, \dots, n, j = 1, \dots, m$ **nm clauses**
- ▶ $x_i \rightarrow z, z \rightarrow y_j, i = 1, \dots, n, j = 1, \dots, m$ **$n + m$ clauses**
- ▶ **new variables can be heads, or singleton bodies with old heads**
- ▶ $\varphi \sim_{\mathcal{R}} \psi$, where variables new are introduced by narrow extension: can be decided efficiently

A $o(n)$ approximation algorithm

Theorem

There is an efficient $O(\frac{n}{\log n})$ -approximation algorithm for narrow extended Horn minimization.

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- ▶ find an equivalent Horn formula with the minimal number of bodies
- ▶ find a decomposition of the bipartite graph between heads and bodies

Hardness of approximation

Theorem

If $NP \not\subseteq DTIME(n^{\text{polylog}(n)})$ then for every $0 < \delta < 1$ restricted Horn minimization is not

$$2^{\log^\delta \text{size}(\varphi)}$$

approximable.

- ▶ approximation preserving reduction from the MINREP problem

Decomposition of graphs

- ▶ Chung, Erdős, Spencer (1983), Bublitz (1986): every n -vertex graph can be partitioned into complete bipartite graphs with $O(n^2 / \log n)$ vertices altogether
- ▶ Tuza (1984): every (a, b) -bipartite graph can be partitioned into complete bipartite graphs with $O(ab / \log(a + b))$ vertices altogether

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- ▶ Kővári, Sós, Turán (1954): every (n, n) -bipartite graph with $c_q n^{2-1/q}$ edges contains a complete (q, q) -bipartite graph (Zarankiewicz problem)
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- ▶ Kirchner (2008): efficient algorithm to find a complete $(\sqrt{\log n}, \sqrt{\log n})$ bipartite graph in a dense n -vertex graph

Efficient decomposition

Theorem

There are efficient algorithms for finding

- ▶ *a complete bipartite subgraph with parts about*

$$\frac{\log n}{\log(n^2/m)}$$

in n -vertex graphs with m edges ($m > n^{3/2}$)

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in n -vertex graphs with m edges ($m > n^{3/2}$)

- ▶ *a decomposition of n -vertex graphs into complete bipartite graphs with $O(n^2/\log n)$ vertices altogether*
- ▶ *a decomposition of (a, b) -bipartite graphs into complete bipartite graphs with $O(ab/\log(a+b))$ vertices altogether*